## **Egg Problems**

William Gasarch-U of MD

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- ▶ If an egg breaks then you can cannot re-use it.

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- 2. How many drops needed if n floors, 1 egg. **Algorithm** Floor 1, ..., floor n 1. n 1 drops.

Get an answer AND prove its optimal.

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   Algorithm Floor 1, ..., floor 99. 99 drops.
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- How many drops needed if n floors, 1 egg.
   Algorithm Floor 1, ..., floor n − 1. n − 1 drops.
   Optimal If skip a floor then cannot know the answer.

## Two Eggs

- 1. How many drops needed if 100 floors, 2 eggs.
- 2. How many drops needed if n floors, 2 eggs. (Can ignore +O(1) terms.)

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#### YOU"VE BEEN PUNKED

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#### A Clue That This is Not Optimal

- ▶ If first drop SPLAT then takes 11 drops.
- ▶ If last drop SPLAT then takes 19 drops.
- ► There should be a way to make the worst case better and the best case worse.

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15th floor. If SPLAT, 14 left, 1+13=14 total (15+14)th Floor. If SPLAT, 13 left, 2+12=14 total

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 $(15 + \cdots + 4)$ th Floor. If SPLAT, 3 left, 12+2=14 total Sum is 101, so actually works for 101 floors.

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Find least k with  $1+2+\cdots+k \ge n$ . Note:  $k \sim 2^{1/2} \times n^{1/2}$ .

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kth floor. If SPLAT, k-1 left, 1+(k-1)=k+1 total

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### Algorithm

kth floor. If SPLAT, k-1 left, 1+(k-1)=k+1 total (k+(k-1))th Floor. If SPLAT, k-2 left, 2+(k-1)=k total

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Contrast

Old Method:  $2 \times n^{1/2}$  drops

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Find least k with 1+2+\cdots+k \geq n. Note: k \sim 2^{1/2} \times n^{1/2}. Algorithm
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kth floor. If SPLAT, k-1 left, 1+(k-1)=k+1 **total** (k+(k-1))th Floor. If SPLAT, k-2 left, 2+(k-1)=k **total** k:  $(k+\cdots+1)$ th Floor.

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Is  $\sim 2^{1/2} \times n^{1/2}$  optimal?

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Is  $\sim 2^{1/2} \times n^{1/2}$  optimal? Vote YES or NO

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Find least k with 1+2+\cdots+k \geq n. Note: k \sim 2^{1/2} \times n^{1/2}. Algorithm
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$$k$$
th floor. If SPLAT,  $k-1$  left,  $1+(k-1)=k+1$  total  $(k+(k-1))$ th Floor. If SPLAT,  $k-2$  left,  $2+(k-1)=k$  total  $(k+\cdots+1)$ th Floor.

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Old Method:  $2 \times n^{1/2}$  drops New Method:  $2^{1/2} \times n^{1/2}$  drops

Is  $\sim 2^{1/2} \times n^{1/2}$  optimal? Vote YES or NO Answer is YES.

# **Optimal**

Two eggs. Given n, the optimal number of drops is least k with  $1 + 2 + \cdots + k \ge n$ .

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**Proof Sketch** Any algorithm that deviates from the one we give has to do worse. Formally you would look at the first step where the algorithm differs from ours.

# Three Eggs

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- 1. How many drops needed if 100 floors, 3 eggs.
- 2. ow many drops needed if n floors, 3 eggs. (Can ignore +O(1) terms.)

Work on in groups.

First need to know some two-egg answers

n	Sum	number of drops
1		0
2, 3	$1+2 \ge 3$	2
4, 5, 6	$1 + 2 + 3 \ge 6$	3
7, 8, 9, 10	$1+2+3+4 \ge 10$	4
$11,\ldots,15$	$1+2+3+4+5 \ge 15$	5
$16,\ldots,21$	$1+2+3+4+5+6 \geq 21$	6
21, , 28	$1+2+3+4+5+6+7 \ge 28$	7
28, , 36	$1+2+3+4+5+6+7+8 \geq 36$	8

How many drops needed if 100 floors, 3 eggs. 36th floor. If SPLAT use 2-egg sol. 1+8=9 (36+28)th floor. If SPLAT use 2-egg sol. 2+7=9 (36+28+21)th floor. If SPLAT use 2-egg sol. 3+6=9 (36+28+21+15+10))th floor. If SPLAT use 2-egg sol. 3+6=9

Recall 2-egg for n least k such that

$$1+2+\cdots+k \geq n \ k \sim 2^{1/2} n^{1/2}$$

Rephrase:

Recall 2-egg for n least k such that

$$\binom{1}{1} + \binom{2}{1} + \dots + \binom{k}{1} \ge n \ k \sim 2^{1/2} n^{1/2}$$

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**2-egg for** n least k such that

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Approximate:

$$\binom{2}{2} + \dots + \binom{k}{2} \sim \sum_{i=2}^{k} \frac{i^2}{2} \sim \frac{k^3}{6}$$

Recall 2-egg for *n* least *k* such that

$$1+2+\cdots+k > n \ k \sim 2^{1/2} n^{1/2}$$

Rephrase:

**Recall 2-egg for** *n* least *k* such that

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Approximate:

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$$\binom{e-1}{e-1} + \cdots + \binom{k}{e-1} = \sum_{i=e-1}^{k} \frac{i^{e-1}}{(e-1)!} \sim \frac{k^e}{e!}$$

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$$\frac{k^e}{e!} \ge n$$
$$k \ge (e!)^e n^{1/e}$$