Roth's Theorem A Dense Enough Set Has a 3-AP

Exposition by William Gasarch and Kelin Zhu

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2) The k = 3 case which involves the Discrete Fourier Transform.

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He proved Hales-Jewitt Thm which implies VDW's Thm.

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$$W(k,c) \le 2^{2^{c^{2^{k+9}}}}$$

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