Review for CMSC 452 Midterm: Grammars

Context Free Languages

Examples of Context Free Grammars

$S \rightarrow aSb$
$S \rightarrow e$

The set of all strings **Generated** is

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$S \rightarrow aSb$
$S \rightarrow e$

The set of all strings **Generated** is

$$L = \{a^n b^n : n \in \mathbb{N}\}$$

Examples of Context Free Grammars

$S \to aSb$
$S \to e$

The set of all strings **Generated** is

$$L = \{a^n b^n : n \in \mathbb{N}\}$$

Note L is context free lang that is not regular.

Context Free Grammar for $\{a^{2n}b^n : n \in \mathbb{N}\}$

$S \to aaSb$
$S \to e$

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$$T \rightarrow aTb$$
$$T \rightarrow e$$
$$A \rightarrow Aa$$
$$A \rightarrow a$$

Context Free Grammars

Def A **Context Free Grammar** is a tuple $G = (N, \Sigma, R, S)$

- N is a finite set of **nonterminals**.
- Σ is a finite **alphabet**. Note $\Sigma \cap N = \emptyset$.
- $R \subseteq N \times (N \cup \Sigma)^*$ and are called **Rules**.
- $S \in N$, the **start symbol**.

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Examples:

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- $A \Rightarrow aB$

Then, if w is string of **non-terminals only**, we define $L(G)$ by:

$$L(G) = \{w \in \Sigma^* \mid S \Rightarrow w\}$$

Number of a's = Number of b's

Is

$$L = \{w \mid \#_a(w) = \#_b(w)\}$$

context free?

YES

Let G be the CFG

$S \rightarrow aSb$

$S \rightarrow bSa$

$S \rightarrow SS$

$S \rightarrow e$

YES

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Thm $L(G) = \{w \mid \#_a(w) = \#_b(w)\}$.

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Let G be the CFG
$S \rightarrow aSb$
$S \rightarrow bSa$
$S \rightarrow SS$
$S \rightarrow e$

Thm $L(G) = \{w \mid \#_a(w) = \#_b(w)\}$.

Note This Theorem is **not obvious**. Deserves a proof! But I won't give one.

Example of a Lang that is NOT a CFL

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3) If $L \subseteq a^*$ and L is not regular than L is not a CFL.

One proves theorems NON CFL using the PL for CFL's (next slide).

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1. $w = uvxyz$ and either $v \neq e$ or $y \neq e$.

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2. $|vxy| \leq n_1$.

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2. $|vxy| \leq n_1$.
3. For all $i \geq 0$, $uv^ixy^iz \in L$.

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Proof involves looking at the Parse Tree for w and finding some nonterminal T twice in the tree. We will not be doing the proof.

Closure Properties and REG$\subset$ CFL

L_1, L_2 CFL $\rightarrow$ $L_1 \cup L_2$ CFL

L_1 is CFL via CFG (N_1, Σ, R_1, S_1).
L_2 is CFL via CFG (N_2, Σ, R_2, S_2).

L_1, L_2 CFL $\rightarrow L_1 \cup L_2$ CFL

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L_2 is CFL via CFG (N_2, Σ, R_2, S_2).

CFL for $L_1 \cup L_2$:
Just add $S \rightarrow S_1$ and $S \rightarrow S_2$ to union of grammars.

L_1, L_2 CFL $\rightarrow L_1 \cap L_2$ CFL

NOT TRUE: $a^n b^n c^* \cap a^* b^n c^n = a^n b^n c^n$.

L_1, L_2 CFL $\rightarrow L_1 \cdot L_2$ CFL

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CFL for $L_1 \cup L_2$:
Just add $S \rightarrow S_1 S_2$ to union of grammars.

L CFL $\rightarrow \overline{L}$ CFL

FALSE.
Let

$$L = \overline{\{a^n b^n c^n : n \in \mathbb{N}\}}$$

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This is a CFL. This will a HW.

L CFL $\rightarrow L^*$ CFL

L is CFL via CFG (N, Σ, R, S).

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This one I leave to you to look up my slides on it.

REG contained in CFL

For every **regex** α, $L(\alpha)$ is a CFL.

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Prove by ind on the length of α.

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Prove by ind on the length of α.

We omit from this review.

Examples of CFL's and Size of CFG's

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Def CFG G is in **Chomsky Normal Form** if the rules are all of the following form:
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2) $A \rightarrow \sigma$ (where $A \in N$ and $\sigma \in \Sigma$).

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Def CFG G is in **Chomsky Normal Form** if the rules are all of the following form:
1) $A \rightarrow BC$ where $A, B, C \in N$ (nonterminals).
2) $A \rightarrow \sigma$ (where $A \in N$ and $\sigma \in \Sigma$).
3) $S \rightarrow e$ (where S is the start state).

Example of Chomsky Normal Form

Chomsky Normal form CFG that generates $\{aaaaaaaa\}$
$S \rightarrow AA$

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$A \rightarrow BB$

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Chomsky Normal form CFG that generates $\{aaaaaaaa\}$
$S \rightarrow AA$
$A \rightarrow BB$
$B \rightarrow CC$

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So $\{aaaaaaaa\}$ has a CFG of size 4.

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Chomsky Normal form CFG that generates $\{aaaaaaaa\}$
$S \rightarrow AA$
$A \rightarrow BB$
$B \rightarrow CC$
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So $\{aaaaaaaa\}$ has a CFG of size 4.

By the same trick $\exists$ a CFG for $\{a^n\}$ of size $O(\log n)$.

- Any DFA or NFA that recognizes $\{a^n\}$ has $n + \Omega(1)$ states.
- There is a CFG that generates $\{a^n\}$ with $O(\log n)$ rules.

$\{a, b\}^* a \{a, b\}^n$

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$\{a, b\}^*$ CONCAT $a\{a, b\}^n$

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$\{a, b\}^*$ CONCAT $a\{a, b\}^n$ has $O(\log n)$ rule CFG.

Any CFG can be Put Into Chomsky Normal Form

Recall the CFG for $\{a^m b^n : m > n\}$. We put it into Chomsky Normal Form.

1) $S \rightarrow AT$
2) $T \rightarrow aTb$
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4) $A \rightarrow Aa$
5) $A \rightarrow a$

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Use nonterminals $[aT]$, $[b]$, $[a]$. Replace $T \rightarrow aTb$ with:

$T \rightarrow [aT][b]$
$[aT] \rightarrow [a]T$
$[b] \rightarrow b$.
$[a] \rightarrow a$

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$[b] \rightarrow b$.
$[a] \rightarrow a$
Repeat the process with the other rules.

MISC

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2) Recall: DFA's are **Recognizers**, Regex are **Generators**.
CFG's are **Generators**. There is a **Recognizer** equivalent to it:

PDA: Push Down Automata

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They are NFAs with a stack.

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2) Recall: DFA's are **Recognizers**, Regex are **Generators**.
CFG's are **Generators**. There is a **Recognizer** equivalent to it:

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They are NFAs with a stack.
The proof that PDA-recognizers and CFG-generators are equivalent is messy so we won't be doing it. We won't deal with PDA's in this course at all.

CNF for $\{w\}$

CNF for $\{aabbbab\}$

Example CNF for $\{aabbbab\}$

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$S \to [A][ABBBAB]$

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1. You can do something similar for any w.
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3. Question we will come back to LATER:
   $(\exists w)$ such that $\{w\}$ requires large CFG?

CFL $\subset$ P

Poly Time Algorithm for CFG Membership

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For $i \leq j$ let

$$\mathrm{GEN}[i,j] = \{A : A \Rightarrow \sigma_i \cdots \sigma_j\}$$

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For $i \le j$ let

$$\mathrm{GEN}[i,j] = \{A : A \Rightarrow \sigma_i \cdots \sigma_j\}$$

We will find **all** $\mathrm{GEN}[i,j]$. Hence we will find $\mathrm{GEN}[1,n]$. Hence we will find if $S \in \mathrm{GEN}[1,n]$.

Bottom Up View

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$$\sigma_1 \cdots \sigma_{i-1} \overbrace{\sigma_i}^{A} \sigma_{i+1} \cdots \sigma_n$$

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$$\mathrm{GEN}[i,i] \; = \; \{A : A \rightarrow \sigma_i\}$$

Bottom Up View

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$$\mathrm{GEN}[i,i] \; = \; \{A : A \to \sigma_i\}$$

$$\sigma_1 \cdots \sigma_{i-1} \overbrace{\sigma_i}^{B} \overbrace{\sigma_{i+1}}^{C} \sigma_{i+2} \cdots \sigma_n$$

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$$\mathrm{GEN}[i,i+1] \;=\; \{A : A \to BC \;\wedge\; B \to \sigma_i \;\wedge\; C \to \sigma_{i+1}\}$$

Bottom Up View

$$\sigma_1 \cdots \sigma_{i-1} \overbrace{\sigma_i}^{A} \sigma_{i+1} \cdots \sigma_n$$

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$$\sigma_1 \cdots \sigma_{i-1} \overbrace{\sigma_i}^{B} \overbrace{\sigma_{i+1}}^{C} \sigma_{i+2} \cdots \sigma_n$$

$$\begin{aligned}
\mathrm{GEN}[i,i+1] &= \{A : A \to BC \ \wedge \ B \to \sigma_i \ \wedge \ C \to \sigma_{i+1}\} \\
&= \{A : A \to BC \\
&\qquad \wedge \ B \in GEN[i,i] \ \wedge \ C \in GEN[i+1,i+1]\}
\end{aligned}$$

Recurrence

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$$\mathrm{GEN}[i,j] \quad = \quad \bigcup_{i \le k < j} \{A : A \to BC \ \wedge \ B \Rightarrow \sigma_i \cdots \sigma_k \ \wedge \ C \Rightarrow \sigma_{k+1} \cdots \sigma_j\}$$

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$$\begin{aligned} \mathrm{GEN}[i,j] &= \bigcup_{i \le k < j} \{A : A \to BC \ \wedge\ B \Rightarrow \sigma_i \cdots \sigma_k \ \wedge\ C \Rightarrow \sigma_{k+1} \cdots \sigma_j\} \\ &= \bigcup_{i \le k < j} \{A : A \to BC \ \wedge\ B \in \mathrm{GEN}[i,k] \ \wedge\ C \in \mathrm{GEN}[k+1,j]\} \end{aligned}$$

The Algorithm

```
for i = 1 to n do
    for j = i to n do
        GEN[i,j] ← ∅

for i = 1 to n do
    for all rules A → σ_i do
        GEN[i,i] ← GEN[i,i] with A

for s = 2 to n do
    for i = 1 to n-s+1 do
        j ← i+s-1 do
        for k = i to j-1 do
            for all rules A → BC
                where B ∈ GEN[i,k] and C ∈ GEN[k+1,j]
                    GEN[i,j] ← GEN[i,j] with A
```