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Tricks for Divisibility and DFA's

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Divisibility tricks for 2,3,5,9,10.



Divisibility tricks for 2,3,5,9,10. Let $x = x_n \cdots x_0$.



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▶ 2: $x \equiv 0 \pmod{2}$ iff $x_0 \equiv 0 \pmod{2}$.

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What is a trick? We come back to that later.

We don't just learn divisibility. Let $x = x_n \cdots x_0$.



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1) We don't just get divisibility, we get mod.

2) Still have not defined trick carefully.

For this Slide Packet $\Sigma = \{0, \ldots, 9\}$.

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Strings are numbers in base 10.



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is the number

$$d_{n-1} \times 10^{n-1} + \cdots + d_1 \times 10^1 + d_0 \times 10^0.$$

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We feed a number into a DFA right-to-left: d_0 , then d_1 then d_2 then

Proof of Trick for Mod. All \equiv are mod 2.

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Thm $d_{n-1}\cdots d_0 \equiv d_0$.



Thm
$$d_{n-1} \cdots d_0 \equiv d_0$$
.
Pf

$$d_{n-1}\times 10^{n-1}+\cdots+d_1\times 10+d_0$$

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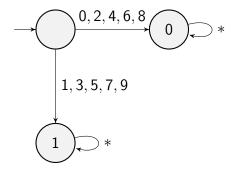
Thm
$$d_{n-1} \cdots d_0 \equiv d_0$$
.
Pf

$$d_{n-1} \times 10^{n-1} + \dots + d_1 \times 10 + d_0$$

= 10(d_{n-1} \times 10^{n-2} + \dots + d_1) + d_0
= d_0

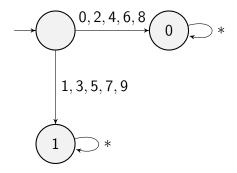
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We don't regard the empty string as even or odd.



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We don't regard the empty string as even or odd. If pressed, I would say even.

Thm $d_{n-1}\cdots d_0 \equiv d_{n-1} + \cdots + d_0$.



Thm $d_{n-1}\cdots d_0 \equiv d_{n-1}+\cdots+d_0$. Pf

 $d_{n-1} \times 10^{n-1} + \cdots + d_1 \times 10 + d_0 \times 10^0$

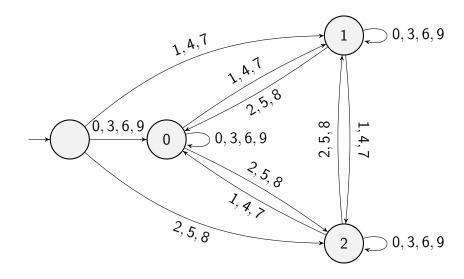
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Thm
$$d_{n-1} \cdots d_0 \equiv d_{n-1} + \cdots + d_0$$
.
Pf

$$egin{aligned} & d_{n-1} imes 10^{n-1} + \dots + d_1 imes 10 + d_0 imes 10^0 \ & \equiv & d_{n-1} imes 1 + \dots + d_1 imes 1 + d_0 imes 1 \end{aligned}$$

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Do you know the Mod 4 trick??

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Do you know the Mod 4 trick?? $n \equiv 0$ iff

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Do you know the Mod 4 trick?? $n \equiv 0$ iff last 2 digits are a number $\equiv 0$.

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Do you know the Mod 4 trick?? $n \equiv 0$ iff last 2 digits are a number $\equiv 0$. **Thm** $d_{n-1} \cdots d_0 \equiv 2d_1 + d_0$.

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$$d_{n-1}\times 10^{n-1}+\cdots+d_1\times 10+d_0$$

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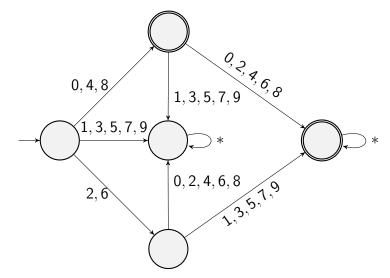
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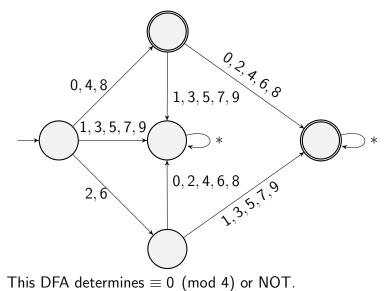
$$d_{n-1} imes 10^{n-1}+\cdots+d_1 imes 10+d_0 \ \equiv d_1 imes 10+d_0$$

Do you know the Mod 4 trick?? $n \equiv 0$ iff last 2 digits are a number $\equiv 0$. **Thm** $d_{n-1} \cdots d_0 \equiv 2d_1 + d_0$. **Pf**

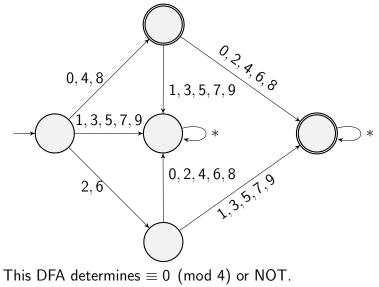
$$d_{n-1} imes 10^{n-1} + \dots + d_1 imes 10 + d_0$$

 $\equiv d_1 imes 10 + d_0$
 $\equiv 2d_1 + d_0.$





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A DFA classifier for mod 4 may be on the HW.

Key to all of these Problems

For all of these problems we need to find a pattern of $10^n \pmod{a}$.

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Key to all of these Problems

For all of these problems we need to find a pattern of 10^n (mod *a*). Mod 2: Pattern is 1,0,0,0,..., DFA cared about first digit.

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For all of these problems we need to find a pattern of 10ⁿ (mod *a*). Mod 2: Pattern is 1,0,0,0,..., DFA cared about first digit. Mod 3: Pattern is 1,1,1,1,..., DFA tracked sum mod 3.

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Proof of Tricks for Mod 5,9,10 and Trick for Mod 6

These may be on a HW.



Is there a trick for mod 11?



Is there a trick for mod 11? We derive it together!



Is there a trick for mod 11? We derive it together! $10^0 \equiv 1$



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Is there a trick for mod 11? We derive it together! 10^0 \equiv 110^1 \equiv 10 \equiv -1
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Is there a trick for mod 11?
We derive it together!
10^0 \equiv 1
10^1 \equiv 10 \equiv -1
10^2 \equiv 10 \times 10 \equiv -1 \times -1 \equiv 1.
```

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```
Is there a trick for mod 11?
We derive it together!
10^0 \equiv 1
10^1 \equiv 10 \equiv -1
10^2 \equiv 10 \times 10 \equiv -1 \times -1 \equiv 1.
10^3 \equiv 10^2 \times 10 \equiv 1 \times -1 \equiv -1.
Pattern is 1, -1, 1, -1, ....
```

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Is there a trick for mod 11?

We derive it together!

10^0 \equiv 1

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10^2 \equiv 10 \times 10 \equiv -1 \times -1 \equiv 1.

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Pattern is 1, -1, 1, -1, \dots

Thm d_n \cdots d_0 \equiv d_0 - d_1 + d_2 - \cdots \pm d_n.
```

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Trick for Mod 11. All \equiv are Mod 11

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We derive it together!

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Pattern is 1, -1, 1, -1, \dots

Thm d_n \cdots d_0 \equiv d_0 - d_1 + d_2 - \dots \pm d_n.
```

Proof may be on HW or Midterm or Final or some combination.

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Need to keep track of both the running weighted sum mod 11 and if you are reading an even or odd place.

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 $Q=\{0,\ldots,10\}\times\{0,1\}$

Need to keep track of both the running weighted sum mod 11 and if you are reading an even or odd place.

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$$Q = \{0, \dots, 10\} \times \{0, 1\}$$

 $s = (0, 0).$

Need to keep track of both the running weighted sum mod 11 and if you are reading an even or odd place.

$$Q = \{0, \dots, 10\} \times \{0, 1\}$$

$$s = (0, 0).$$

Final state: Not going to have these, this is DFA-classifier.

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$$s = (0, 0).$$

Final state: Not going to have these, this is DFA-classifier.

$$\delta((i,j),\sigma) \begin{cases} (i+\sigma \pmod{11}, j+1 \pmod{2}) & \text{if } j=0\\ (i-\sigma \pmod{11}, j+1 \pmod{2}) & \text{if } j=1\\ \end{cases}$$
(1)

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We keep track of a running weighted sum mod 11 and position of the symbol mod 2.

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We keep track of a running weighted sum mod 11 and position of the symbol mod 2.

22 states.

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Final state: Not going to have these, this is DFA-classifier.

$$\delta((i,j),\sigma) \begin{cases} (i+\sigma \pmod{11}, j+1 \pmod{2}) \text{ if } j=0\\ (i-\sigma \pmod{11}, j+1 \pmod{2}) \text{ if } j=1\\ \end{cases}$$
(1)

We keep track of a running weighted sum mod 11 and position of the symbol mod 2.

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22 states.

Classifier If end in (i, 0) or (i, 1) then number is $\equiv i$.

Is there a trick for mod 7?



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Is there a trick for mod 7? Answer Depends what you call a trick.

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Is there a trick for mod 7? Answer Depends what you call a trick. We need to spot a pattern.

Is there a trick for mod 7? Answer Depends what you call a trick. We need to spot a pattern. $10^0 \equiv 1$

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Is there a trick for mod 7? Answer Depends what you call a trick. We need to spot a pattern. $10^0 \equiv 1$ $10^1 \equiv 3$

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Is there a trick for mod 7? Answer Depends what you call a trick. We need to spot a pattern. $10^0 \equiv 1$ $10^1 \equiv 3$ $10^2 \equiv 10 \times 10 \equiv 3 \times 3 \equiv 9 \equiv 2$

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Is there a trick for mod 7? **Answer** Depends what you call a trick. We need to spot a pattern. $10^0 \equiv 1$ $10^1 \equiv 3$ $10^2 \equiv 10 \times 10 \equiv 3 \times 3 \equiv 9 \equiv 2$ $10^3 \equiv 10^2 \times 10 \equiv 2 \times 3 \equiv 6$

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Is there a trick for mod 7? Answer Depends what you call a trick. We need to spot a pattern. $10^0 \equiv 1$ $10^1 \equiv 3$ $10^2 \equiv 10 \times 10 \equiv 3 \times 3 \equiv 9 \equiv 2$ $10^3 \equiv 10^2 \times 10 \equiv 2 \times 3 \equiv 6$ $10^4 \equiv 10^3 \times 10 \equiv 6 \times 3 \equiv 18 \equiv 4$

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Is there a trick for mod 7? **Answer** Depends what you call a trick. We need to spot a pattern. $10^0 = 1$ $10^{1} \equiv 3$ $10^2 \equiv 10 \times 10 \equiv 3 \times 3 \equiv 9 \equiv 2$ $10^3 = 10^2 \times 10 = 2 \times 3 = 6$ $10^4 \equiv 10^3 \times 10 \equiv 6 \times 3 \equiv 18 \equiv 4$ $10^5 = 10^4 \times 10 = 4 \times 3 = 12 = 5$

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Is there a trick for mod 7? **Answer** Depends what you call a trick. We need to spot a pattern. $10^0 = 1$ $10^{1} \equiv 3$ $10^2 \equiv 10 \times 10 \equiv 3 \times 3 \equiv 9 \equiv 2$ $10^3 = 10^2 \times 10 = 2 \times 3 = 6$ $10^4 \equiv 10^3 \times 10 \equiv 6 \times 3 \equiv 18 \equiv 4$ $10^5 = 10^4 \times 10 = 4 \times 3 = 12 = 5$ $10^6 \equiv 10^5 \times 10 \equiv 5 \times 3 \equiv 15 \equiv 1$

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Is there a trick for mod 7? **Answer** Depends what you call a trick. We need to spot a pattern. $10^0 = 1$ $10^{1} \equiv 3$ $10^2 \equiv 10 \times 10 \equiv 3 \times 3 \equiv 9 \equiv 2$ $10^3 = 10^2 \times 10 = 2 \times 3 = 6$ $10^4 \equiv 10^3 \times 10 \equiv 6 \times 3 \equiv 18 \equiv 4$ $10^5 = 10^4 \times 10 = 4 \times 3 = 12 = 5$ $10^6 \equiv 10^5 \times 10 \equiv 5 \times 3 \equiv 15 \equiv 1$ Pattern is 1, 3, 2, 6, 4, 5, 1, 3, 2, 6, 4, 5, 1,

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Is there a trick for mod 7? **Answer** Depends what you call a trick. We need to spot a pattern. $10^0 = 1$ $10^{1} \equiv 3$ $10^2 \equiv 10 \times 10 \equiv 3 \times 3 \equiv 9 \equiv 2$ $10^3 = 10^2 \times 10 = 2 \times 3 = 6$ $10^4 \equiv 10^3 \times 10 \equiv 6 \times 3 \equiv 18 \equiv 4$ $10^5 = 10^4 \times 10 = 4 \times 3 = 12 = 5$ $10^6 \equiv 10^5 \times 10 \equiv 5 \times 3 \equiv 15 \equiv 1$ Pattern is 1, 3, 2, 6, 4, 5, 1, 3, 2, 6, 4, 5, 1, Can we use this?

Want to know what 3876554 is mod 7.

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Want to know what 3876554 is mod 7.

3876554

 $= \ 3\cdot 10^6 + 8\cdot 10^5 + 7\cdot 10^4 + 6\cdot 10^3 + 5\cdot 10^2 + 5\cdot 10 + 4$

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Want to know what 3876554 is mod 7.

3876554

- $= \ 3\cdot 10^6 + 8\cdot 10^5 + 7\cdot 10^4 + 6\cdot 10^3 + 5\cdot 10^2 + 5\cdot 10 + 4$
- $\equiv 3 \cdot 1 + 8 \cdot 5 + 7 \cdot 4 + 6 \cdot 6 + 5 \cdot 2 + 5 \cdot 3 + 4 \pmod{7}$

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Want to know what 3876554 is mod 7.

3876554

- $= 3 \cdot 10^6 + 8 \cdot 10^5 + 7 \cdot 10^4 + 6 \cdot 10^3 + 5 \cdot 10^2 + 5 \cdot 10 + 4$
- $\equiv 3 \cdot 1 + 8 \cdot 5 + 7 \cdot 4 + 6 \cdot 6 + 5 \cdot 2 + 5 \cdot 3 + 4 \pmod{7}$
- $\equiv 3 \cdot 1 + 1 \cdot 5 + 0 \cdot 4 + -1 \cdot 6 + -2 \cdot 2 + -2 \cdot 3 + 4 \pmod{7}$

Want to know what 3876554 is mod 7.

3876554

- $= 3 \cdot 10^6 + 8 \cdot 10^5 + 7 \cdot 10^4 + 6 \cdot 10^3 + 5 \cdot 10^2 + 5 \cdot 10 + 4$
- $\equiv 3 \cdot 1 + 8 \cdot 5 + 7 \cdot 4 + 6 \cdot 6 + 5 \cdot 2 + 5 \cdot 3 + 4 \pmod{7}$
- $\equiv 3 \cdot 1 + 1 \cdot 5 + 0 \cdot 4 + -1 \cdot 6 + -2 \cdot 2 + -2 \cdot 3 + 4 \pmod{7}$

 $\equiv 3+5+0-6-4-6+4 \pmod{7}$

Want to know what 3876554 is mod 7.

3876554

- $= \ 3\cdot 10^6 + 8\cdot 10^5 + 7\cdot 10^4 + 6\cdot 10^3 + 5\cdot 10^2 + 5\cdot 10 + 4$
- $\equiv 3 \cdot 1 + 8 \cdot 5 + 7 \cdot 4 + 6 \cdot 6 + 5 \cdot 2 + 5 \cdot 3 + 4 \pmod{7}$
- $\equiv 3 \cdot 1 + 1 \cdot 5 + 0 \cdot 4 + -1 \cdot 6 + -2 \cdot 2 + -2 \cdot 3 + 4 \pmod{7}$

- $\equiv 3+5+0-6-4-6+4 \pmod{7}$
- \equiv 3 (mod 7)

Want to know what 3876554 is mod 7.

3876554

- $= \ 3\cdot 10^6 + 8\cdot 10^5 + 7\cdot 10^4 + 6\cdot 10^3 + 5\cdot 10^2 + 5\cdot 10 + 4$
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- $\equiv 3+5+0-6-4-6+4 \pmod{7}$
- \equiv 3 (mod 7)

DFA States will keep track of

Want to know what 3876554 is mod 7.

3876554

- $= \ 3\cdot 10^6 + 8\cdot 10^5 + 7\cdot 10^4 + 6\cdot 10^3 + 5\cdot 10^2 + 5\cdot 10 + 4$
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- $\equiv 3+5+0-6-4-6+4 \pmod{7}$
- \equiv 3 (mod 7)

DFA States will keep track of Running weighted sum mod 7

Want to know what 3876554 is mod 7.

3876554

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- $\equiv 3 \cdot 1 + 8 \cdot 5 + 7 \cdot 4 + 6 \cdot 6 + 5 \cdot 2 + 5 \cdot 3 + 4 \pmod{7}$

$$\equiv 3 \cdot 1 + 1 \cdot 5 + 0 \cdot 4 + -1 \cdot 6 + -2 \cdot 2 + -2 \cdot 3 + 4 \pmod{7}$$

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- $\equiv 3+5+0-6-4-6+4 \pmod{7}$
- \equiv 3 (mod 7)

DFA States will keep track of Running weighted sum mod 7 Position of digit mod 6 so know which weights to use.

Want to know what 3876554 is mod 7.

3876554

- $= \ 3\cdot 10^6 + 8\cdot 10^5 + 7\cdot 10^4 + 6\cdot 10^3 + 5\cdot 10^2 + 5\cdot 10 + 4$
- $\equiv 3 \cdot 1 + 8 \cdot 5 + 7 \cdot 4 + 6 \cdot 6 + 5 \cdot 2 + 5 \cdot 3 + 4 \pmod{7}$

$$\equiv 3 \cdot 1 + 1 \cdot 5 + 0 \cdot 4 + -1 \cdot 6 + -2 \cdot 2 + -2 \cdot 3 + 4 \pmod{7}$$

- $\equiv 3+5+0-6-4-6+4 \pmod{7}$
- \equiv 3 (mod 7)

DFA States will keep track of Running weighted sum mod 7 Position of digit mod 6 so know which weights to use. So there are $7 \times 6 = 42$ states.

Is the Method a Trick?

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Is the Method a Trick?

YES A DFA can do it.



Is the Method a Trick?

YES A DFA can do it.

NO A human cannot do it easily. (The pattern is not like $1,1,1,\ldots$ or mostly 0's.)

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Too hard for me ...



Too hard for me ...

... but not for you.



Too hard for me ...

... but not for you.

Might make it a HW to do as a table.

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Possible Research Question

What is the fastest way to determine $n \pmod{7}$?



Possible Research Question

What is the fastest way to determine $n \pmod{7}$? Method One Divide and take remainder.

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What is the fastest way to determine $n \pmod{7}$? Method One Divide and take remainder. Method Two Use the DFA.

What is the fastest way to determine n (mod 7)?Method One Divide and take remainder.Method Two Use the DFA.Question Which is faster?

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What is the fastest way to determine n (mod 7)?
Method One Divide and take remainder.
Method Two Use the DFA.
Question Which is faster?
Might be hard to tell because today's computers are so fast!

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