

Problem

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Cocke-Younger-Kasimi Algorithm

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Cocke-Younger-Kasimi Algorithm

Sakai (1962), Kasimi (1965), Younger (1967),
Cocke/Schwartz (1970)

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Let L be a CFL. Let G be the Chomsky Normal Form CFG for L .

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We will obtain $p(n) = O(n^3)$.

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We will assume that $e \notin L$ and hence we do not have the rule

$$S \rightarrow e.$$

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Our proof can be modified to accommodate this case.

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$$\text{GEN}[1, n] = \{A : A \Rightarrow \sigma_1 \cdots \sigma_n\}$$

We are really asking: Is $S \in \text{GEN}[1, n]$?

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All we want to know is:

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We will solve a harder problem:

Find $\text{GEN}[1, n]$

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Hence we will find if $S \in \text{GEN}[1, n]$.

Why solve this harder problem?

We will use Dynamic programming so having some $\text{GEN}[i, j]$ solved will help us solve later ones.

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$$\text{GEN}[i, i] = \{A : A \rightarrow \sigma_i\}$$

$$\sigma_1 \cdots \sigma_{i-1} \overbrace{\sigma_i}^B \overbrace{\sigma_{i+1}}^C \sigma_{i+2} \cdots \sigma_n$$

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$$\text{GEN}[i, i + 1] = \{A : A \rightarrow BC \wedge B \rightarrow \sigma_i \wedge C \rightarrow \sigma_{i+1}\}$$

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$$\begin{aligned} \text{GEN}[i, i+1] &= \{A : A \rightarrow BC \wedge B \rightarrow \sigma_i \wedge C \rightarrow \sigma_{i+1}\} \\ &= \{A : A \rightarrow BC \\ &\quad \wedge B \in \text{GEN}[i, i] \wedge C \in \text{GEN}[i+1, i+1]\} \end{aligned}$$

Bottom Up View (Continued)

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$$\begin{aligned} \text{GEN}[i, i+2] &= \{A : A \rightarrow BC \\ &\quad \wedge ((B \rightarrow \sigma_i \wedge C \Rightarrow \sigma_{i+1}\sigma_{i+2}) \\ &\quad \vee (B \Rightarrow \sigma_i\sigma_{i+1} \wedge C \rightarrow \sigma_{i+2})) \} \end{aligned}$$

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Bottom Up View (Continued)

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GEN[i,i+3]

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The Algorithm

```
for i = 1 to n do
  for j = i to n do
    GEN[i,j] ← ∅

for i = 1 to n do
  for all rules  $A \rightarrow \sigma_i$  do
    GEN[i,i] ← GEN[i,i] with A

for s = 2 to n do
  for i = 1 to n-s+1 do
    j ← i+s-1 do
      for k = i to j-1 do
        for all rules  $A \rightarrow BC$ 
          where  $B \in \text{GEN}[i,k]$  and  $C \in \text{GEN}[k+1,j]$ 
            GEN[i,j] ← GEN[i,j] with A
```