# Three Hundred Million Points Suffice

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Received February 3, 1987

There is a graph G with 300,000,000 vertices and no clique on four points, such that if its edges are two colored these must be a monochromatic triangle. Acadenuc Press. Inc.

#### HISTORY AND SUMMARY

In the late 1960s Paul Erdös asked what graphs G, other than  $K_6$ , had the property that  $G \rightarrow (K_3)$ . We use the Rado arrow notation:  $G \rightarrow (H)$  is the statement that if the edges of  $G$  are two colored there exists a monochormatic H and, more generally,  $G \rightarrow (H)$ , is the statement that if the edges of G are r-colored there exists a monochromatic  $H$ . In particular, Erdös asked if there is a graph  $G$  satisfying

$$
G \to (K_3) \tag{(*)}
$$

$$
\omega(G) = 3. \tag{(*)}
$$

A proof of the existence of such a G was first given by Jon Folkman  $[2]$ . This supremely ingenious proof had two drawbacks. First, the graph G given was extremely large. Second, the proof did not generalize to give for all r a graph G with  $\omega(G) = 3$  and  $G \rightarrow (K_3)$ . At the Combinatorial Conference in Kesthely, Hungary 1973 this problem was given to the Czechoslovakian mathematician Jarik Nesetril and his young student Vojtech Rodl. They [4] found a completely different argument that for all r graphs G exist with  $\omega(G) = 3$  and  $G \rightarrow (K_3)$ . Those of us at that meeting (see  $[5]$  for an anecdotal account) recall the sense of excitement accompanying that discovery and I feel it played a critical role in the development of modern Ramsey Theory. The graphs given by the Nesetril-Rod1 methods were still extremely large and Erdös offered a reward for the discovery of a G satisfying  $(*)$  having less than 10<sup>10</sup> vertices. Here we claim this reward.

The method used has been known for seveal years to Szemeredi, Nesetril, Rodl, Frankl, and others. Frankl and Rodl [3] calculated that a graph G datisfying (\*) with roughly  $7 \times 10^{11}$  vertices exists. Our note may be considered a case study in the application of asymptotic methods to give precise bounds. The method is extremely case specific. It does not give, for example, graphs G of moderate size satisfying  $\omega(G) = 3$  and  $G \rightarrow (K_3)_{3}$ . This remains an intriguing open problem.

#### 1. THE METHOD

Let  $G = G(n, p)$  be the random graph on *n* vertices with edge probability p. For each  $K_4$  in G randomly select an edge. Delete these edges from  $G$ , giving  $G^*$ . We show that for appropriate n, p (\*) is satisfied by  $G^*$  with positive probability. It shal be convenient to write  $p = cn^{-1/2}$ . In the end we will minimize *n* by taking  $c$  roughly 6, and *n* roughly 3E8. Set

$$
U = \{(x, xyz): xyz \text{ is a triangle of } G\}
$$
  

$$
U^* = \{(x, xyz): xyz \text{ is a triangle of } G^*\}.
$$

*Note.* xy, xyz shall denote the sets  $\{x, y\}$ ,  $\{x, y, z\}$  throughout. For each vertex  $x$  set

$$
N(x) = \{ y : xy \in G \}
$$

and

$$
A(x) = \text{maximum over all partitions } N(x) = T \cup B \text{ of the number of edges } yz \in G \text{ with } y \in T \text{ and } z \in B.
$$

THEOREM. If

$$
\sum_{x} A(x) < \frac{2}{3} |U^*| \tag{**}
$$

then  $G^*$  satisfies  $(*)$ .

*Proof.* Clearly  $G^*$  has no  $K_4$ ; suppose there is a coloring  $\chi$  with no monochromatic triangle. We count pairs  $(x, xyz)$  such that  $xyz$  is a triangle of  $G^*$  and  $\chi(xy) \neq \chi(xz)$ . For each triangle xyz the coloring is essentially unique (two red edges and a blue edge or vice versa) and there are two choices of x so that  $(x, xyz)$  is counted so the number of pairs is precisely  $\frac{3}{15}$  is  $\frac{1}{15}$  in  $\frac{1}{15}$  is unique nature of the matricer of pairs is precisely  $\frac{1}{3}$  |  $\sigma$  |  $\frac{1}{3}$  |  $\sigma$  |  $\frac{1}{3}$  |  $\sigma$  and case  $\frac{1}{3}$  |  $\sigma$  |  $\frac{1$ 

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 $\{ y \in N(x) : y(xy) = blue \}, T(x) = N(x) - B(x)$ . Then the number of  $(x, xyz)$ counted is precisely the number of edges  $yz \in G^*$  with  $y \in T(x)$ ,  $z \in B(x)$ . Replacing  $G^*$  by the larger G can only increase this number, and replacing the partition  $T(x)$ ,  $B(x)$  by the optimal partition T, B can only increase this number so that the number of  $(x, xyz)$  is at most  $A(x)$  and the total number of such pairs is at most  $\sum A(x)$  which would contradict (\*\*).

We shall show for appropriate  $n, p$  that  $(**)$  holds with positive probability.

## 2. THE CALCULATION IGNORING VARIANCE

### Let

 $T =$  number of triangles in G

 $Q =$  number of  $K_4$  in G

 $R =$  number of  $(xy, uv, a)$  with x, y, u, v, a distinct,  $ax, ay \in G$ , xyuv forming a  $K_4$  in G, xy selected from xyuv to be removed from  $F^*$ .

Clearly  $|U| = 3T$ . Also  $|U-U^*| \le 2Q + R$ . For suppose  $(a, axy) \in U-U^*$ . Then xy was in a  $K_4$  of G and was deleted and  $ax, ay \in G$ . If the  $K_4$  does not contain a it is counted in R; those  $(a, \, \alpha xy)$ , where the  $K_4$  contains a are at most 2Q in number, since each  $K_4$  abxy chooses one edge xy and contributes  $axy, bxy$  to  $U - U^*$ . Together,

$$
|U^*| > 3T - 2Q - R.
$$

We find expectations

$$
E(T) = \binom{n}{3} p^3 \sim (c^3/6)n^{3/2} \tag{1}
$$

$$
E(Q) = \binom{n}{4} p^6 \sim (c^6/24)n \tag{2}
$$

$$
E(R) = 30 {n \choose 5} p^8/6 \sim (c^8/24)n
$$
 (3)

so that

$$
E(|U^*|) > \frac{1}{2}c^3n^{3/2} - (c^6/12 + c^8/24)n.
$$
 (4)

In the next section we examine variances and show that  $|U^*|$  is "very often" " very close" to its expectation.

Now we examine  $A(x)$ . Set

$$
d = d(x) = |N(x)|
$$
  

$$
e = e(x) = \text{number of edges of } G \text{ in } N(x).
$$

Conditioning on values  $d, e, N(x)$  becomes a random graph H with d vertices and e edges.

For a partition  $N(x) = T \cup B$  let  $X_T$  be the number of edges of H from T to B. Assume  $|T| = |B| = d/2$ , that being the extreme case. Then  $X_T$  has basically binomial distribution  $B(e, \frac{1}{2})$  as e edges are selected and each has probability  $\frac{1}{2}$  of "crossing." Employing the basic Chernoff bound

$$
\Pr[X_T > \frac{1}{2}e + \frac{1}{2}e^{1/2}s] < \exp(-s^2/2). \tag{5}
$$

We set  $s = (2d \ln 2)^{1/2}(1.01)$  so that this probability is  $\le 2^{-d}$ . But  $A(x) =$ max  $X_T$ , over  $2^d$  possible T, so

$$
\Pr[A(x) > \frac{1}{2}e + \frac{1}{2}e^{1/2}s] \leq 2^d 2^{-d} \leq 1. \tag{6}
$$

That is, "almost always," all

$$
A(x) < \frac{1}{2}e(x) + d(x)^{1/2} e(x)^{1/2} \left(\frac{1}{2} \ln 2\right)^{1/2} (1.01). \tag{7}
$$

Now  $\sum_{x} e(x) = 3T \sim \frac{1}{2}c^3 n^{3/2}$ , all  $d(x) \sim np$ , all  $e(x) \sim \frac{1}{2}c^3 n^{1/2}$  so

$$
\sum |A(x)| < c^3 n^{3/2}/4 + n(np)^{1/2} (c^3 n^{1/2}/2)^{1/2} (\ln 2/2)^{1/2}.
$$
 (8)

Combining  $(4)$ ,  $(8)$ ,  $(**)$  holds if

$$
c^{3}n^{3/2}/4 + n(cn^{1/2})^{1/2}(c^{3}n^{1/2}/2)^{1/2}(\ln 2/2)^{1/2}
$$
  
< 
$$
< c^{3}n^{3/2}/3 - [c^{6}/18 + c^{8}/36]n;
$$
 (9)

i.e., if

$$
\left[\frac{c^6}{18}\left(1+\frac{c^2}{2}\right)\middle/\left(\frac{c^3}{12}-\frac{c^2(\ln 2)^{1/2}}{2}\right)\right]^2 < n,\tag{10}
$$

where the LHS must have positive denominator. We take  $c \sim 6$  to minimize this inequality so that  $n \sim 2.7 \times 10^8$ . We allow ourselves a little room and set  $c = 6$ ,  $n = 3 \times 10^8$  in the next section. We know that (\*\*) holds "almost always"-i.e., with probability approaching unity as  $n$  approaches  $\mu$ <sub>with</sub>  $\mu$ <sub>b</sub>  $\mu$ <sub>b</sub> $\mu$ <sub>b</sub>  $\mu$ <sub>b</sub> $\mu$ <sub>b</sub> $\mu$ <sub>b</sub> $\mu$ <sub>b</sub> $\mu$ <sub>b</sub> $\mu$ <sub>b</sub> $\mu$ <sub>b</sub> $\mu$ <sub>b</sub> $\mu$ <sub>b</sub> $\mu$ <sub>b</sub> $\mu$ <sub>b</sub> $\mu$ <sub>b</sub> $\mu$ <sub>b</sub> holds with positive probability.

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# 3. THE CALCULATION

Set  $c=6$ ,  $n=3E8$ ,  $p=cn^{-1/2}$ . We find (to three significant decimals)

$$
E(E) = 1.87E14 \qquad \text{Var}(T) < 5E16 \tag{11}
$$

$$
E(Q) = 5.83E11 \quad \text{Var}(Q) < E12 \tag{12}
$$

$$
E(R) = 2.10E13 \qquad \text{Var}(R) < 6E16. \tag{13}
$$

The variance calculations are cumbersome though elementary exercises. We employ Chebyschev's inequality in the form

$$
Pr[|X - E(X)| > tE(X)] < t^{-2} \operatorname{Var}(X)/E(X)^{2}.\tag{14}
$$

Taking  $t = 10^{-3}$  with  $X = T$ , Q, and R above we find

$$
Pr[1.88E14 > T > 1.86E14] > 0.999
$$
 (15)

$$
Pr[Q < 5.84E11] > 0.999 \tag{16}
$$

$$
Pr[R < 2.11E13] > 0.999. \tag{17}
$$

Let BAD(x) be the event, setting  $e = e(x)$ ,  $d = d(x)$  given by

$$
\text{BAD}(x)\colon A(x) > \frac{1}{2}e\left(\frac{d}{d-1}\right) + e^{1/2} \frac{d^{1/2}(\frac{1}{2}\ln 2)^{1/2}(1.01)}{18} \tag{18}
$$

and let BAD be the disjunction of the events  $BAD(x)$  over all vertices x. We show

$$
Pr[BAD] < 0.01,
$$
 (19)

for which it suffices to show

$$
\Pr[\text{BAD}(x)] < 3E - 10. \tag{20}
$$

The degree  $d(x)$  has distribution  $B(n-1, p)$  which has mean  $(n-1)p =$ 1.04E5 and variance  $(n - 1)p(1 - p) = 1.04E5$ . We use the Chernoff bounds (see, e.g., [6; or 1, sect. 1.31)

$$
Pr[B(m, p) < mp - a] < exp[-a^2/2pm] \tag{a > 0} \tag{21}
$$

$$
Pr[B(m, p) > mp + a] < exp[-a^{2}/2pm + a^{3}/2(pm)^{2}] \qquad (a > 0). (22)
$$

First, quite roughly, take  $a = E4$  and note

$$
Pr[d(x) < 0.9E5] < exp[-10^8/2p(n-1)] < 10^{-100}.\tag{23}
$$

To show (20) it suffices to show

$$
Pr[ BAD(x) | d(x) = d, e(x) = e] < 3 \times 10^{-10} - 10^{-100}
$$
 (24)

for every d, e with  $d \ge 0.9E5$ . Conditioning on d, e we may consider  $N(x)$  as a random graph  $H = (V(H), E(H))$  with d vertices and e edges. For each  $S \subset V(H)$  let  $Y_S$  be the number of  $yz \in E(H)$  with  $y \in S$ ,  $z \notin S$ . Let HYP[N, M,  $r$ ] denote the hypergeometric distribution of the number of red balls from an urn of M red and  $(N-M)$  nonred balls selected in r trials without replacement. Letting  $|S| = s$ ,  $Y_s$  has precisely the distribution HYP $[(\frac{d}{2}), s(d-s), e]$ . Set

$$
b = \frac{1}{2}e(d/(d-1)) + e^{1/2} d^{1/2}(\frac{1}{2}\ln 2)^{1/2}(1.01),
$$
 (25)

for convenience. Clearly  $Pr[Y_s > b]$  is maximized when  $s(d-s)$  is maximized, i.e., at  $s = \lfloor d/2 \rfloor$ . Setting

$$
q' = \left[\frac{d}{2}\right](d - \left[\frac{d}{2}\right]) \left/ \binom{d}{2},\right.\tag{26}
$$

for convenience,

$$
\Pr[Y_S > b] \le \Pr\left[\text{HYP}\left[\binom{d}{2}, q'\binom{d}{2}, e\right] > b\right] \tag{27}
$$

W. Uhlmann [7] has made a systematic comparison between  $HYP[N]$ ,  $Nq, r$ ] and the corresponding binomial  $B(r, q)$ —the distribution given by electing balls with replacement. For our values,

$$
\Pr\left[\text{HYP}\left[\binom{d}{2}, q'\binom{d}{2}, e\right] > b\right] \leqslant \Pr[B(e, q') > b] \leqslant \Pr[B(e, q) > b],\tag{28}
$$

setting  $q = \frac{1}{2}(d/(d-1))$ , a convenient upper bound on q'. We use the bound (again see, e.g.,  $\lceil 6 \text{ or } 1 \rceil$ )

$$
Pr[B(e, q) > eq + a] < exp(-2a^2/e) \qquad (a > 0),
$$
 (29)

valid for all  $e$ ,  $q$ . Then

$$
\Pr[\,Y_{S} > b\,] < \exp[\,-2(1.01)^2\,d(\ln 2)/2\,] < 2^{-d(1.02)}.\tag{30}
$$

**Hence** 

$$
\Pr[\text{BAD}(x) | d(x) = d, e(x) = e] < \sum \Pr[Y_S] < 2^{d} 2^{-d(1.02)} \\
= 2^{-0.02d} < 2^{-1800}, \tag{31}
$$

giving (24) with "plenty of room."

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Application of (21), (22) with precise values give

$$
Pr[d(x) > 1.06E5] < 0.2/n
$$
 (32)

$$
Pr[d(x) < 1.01E5] < 0.1/n,\tag{33}
$$

so that, with room to spare,

$$
Pr[1.01E5 \le d(x) \le 1.06E5 \text{ for all } x] > 0.7. \tag{34}
$$

Combining  $(15)-(17)$ ,  $(19)$ ,  $(34)$  we have that, with probability at least 0.65, the pair  $G, G^*$  satisfy

$$
1.86E14 < T < 1.88E14
$$
  
\n
$$
Q < 5.84E11
$$
  
\n
$$
R < 2.11E13
$$
  
\n
$$
A(x) < b, \quad \text{all } x
$$
  
\n
$$
101000 \le d(x) \le 106000, \quad \text{all } x.
$$
\n(35)

Let  $G, G^*$  be a specific graph pair satisfying the above. Then

$$
\sum A(x) = \frac{1}{2}(1.00001) \sum e(x) + (1.01)(\frac{1}{2}\ln 2)^{1/2} \sum e(x)^{1/2} d(x)^{1/2}.
$$
 (36)

We note  $\sum e(x) = 3T$  and bound

$$
\sum e(x)^{1/2} d(x)^{1/2} \le (106000)^{1/2} \sum e(x)^{1/2}
$$
  
 
$$
\le (106000)^{1/2} (37n)^{1/2}
$$
 (37)

as, in general<sup>t</sup>  $y_1^{1/2} + \cdots + y_n^{1/2} \leq (y_1 + \cdots + y_n)^{1/2} n^{1/2}$ . Plugging in values

$$
\sum A(x) < 2.83E14. \tag{38}
$$

On the other side.

$$
2 |U^*|/3 \ge 2T - (2/3)(2Q + R) > 3.57E14,
$$
 (39)

so that, indeed, the conditions of the theorem hold and  $G^* \to (K_3)$ .

There was plenty of room in our variance arguments. But even if all variances were zero without further argumentation we could not improve on the value  $c = 6.0157$  and a graph G with 266, 930, 400 vertices.

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