### Note

# A Note on Constructive Lower Bounds for the Ramsey Numbers R(3, t)

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We present a simple explicit construction, in terms of t, of a graph that is triangle-free, has independence number t, and contains more than  $\frac{1}{6}((t-1)/2)^{\log_6(\log_4 t} \in \Omega(t^{1.29}))$  vertices. This result is a (feasibly) constructive proof that the Ramsey number  $R(3, t) \in \Omega(t^{1.29})$ . This improves the best previous constructive lower bound of  $R(3, t) > t^{(2 \log_2 t)/3(\log_3 - \log_2 t)} \in \Omega(t^{1.13})$ , due to P. Erdős (1966, J. Combin. Theory 17, 149–153). Also, our result yields a simple explicit construction, in terms of t, of a triangle-free t-chromatic graph whose size is  $O(t^{\log_2 t/3}(\log_3 t)) \in O(t^{4.42})$ .

The Ramsey number R(s, t) is the smallest integer for which every graph on R(s, t) vertices contains either a clique of size s or an independent set of size t. Ramsey [12] shows that, for all s and t, R(s, t) is well defined. The determination of R(s, t) has proven to be extremely difficult for all but a few values of s and t.

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Several upper and lower bounds on R(s, t) for various values of s and t are known. Many of the lower bounds are proven by nonconstructive methods. These methods, introduced by Erdős [6], establish the existence of graphs of size n(s, t) with clique size s and independence number t (which implies that R(s, t) > n(s, t)), but they yield no explicit constructions of these graphs. If a proof also contains an algorithm that, on input s and t, produces an instance of a graph of size n(s, t) with the required properties, and the algorithm runs in time polynomial in n(s, t) then the proof is (feasibly) constructive. Many constructive lower bounds are known, but they are considerably weaker than their nonconstructive counterparts.

The problem of finding good constructive lower bounds for the "diagonal" Ramsey numbers R(t, t) has received considerable attention [1, 4, 9–11]. The most recent result, due to Frankl and Wilson [10], yields the strongest constructive lower bound to date,

$$R(t, t) \ge \exp\left(\frac{(1-\varepsilon)(\log t)^2}{4\log\log t}\right),$$

for all  $\varepsilon > 0$  and sufficiently large t. There is still a significant gap between this and the best nonconstructive lower bound, which is  $R(t, t) \in 2^{\Omega(t)}$ .

For the Ramsey numbers of the form R(3, t), Erdős [8] constructively proves that

$$R(3, t) > t^{(2 \log 2)/3(\log 3 + \log 2)} \in \Omega(t^{1.13}),$$

whereas nonconstructive lower bounds of the form  $R(3, t) \in \Omega((t/\ln t)^2)$  are known (originally shown by Erdős [7]; improved by a constant factor by Spencer [14]). Also, upper bounds are of the form  $R(3, t) \in O(t^2/\log t)$  are known (originally shown by Ajtai, Komlós, and Szemerédi [2, 3]; improved by a constant factor by Shearer [13]).

In this note, we present a simple constructive proof that

$$R(3, t) > \frac{5}{6} \left(\frac{t-1}{2}\right)^{\log 6/\log 4} \in \Omega(t^{1.29}).$$

This is an improvement of the earlier version of these results in [5].

In addition to the connection with Ramsey theory, our result contributes to the problem of explicitly constructing small triangle-free graphs with large chromatic numbers. Our result yields immediately an explicit construction, in terms of k, of a triangle-free k-chromatic graph of size  $O(k^{\log 6/(\log 6 - \log 4)}) \subset O(k^{4.42})$ .

Our method is based on a construction that transforms a graph G to a graph consisting of six disjoint copies of G connected by additional edges

in a particular way. This construction preserves the triangle-freeness property and increases the independence number of the graph by a factor of four, while increasing the size of the graph by a factor of six. By repeatedly applying the transformation, we obtain an explicit construction in terms of t of triangle-free graphs with independence number t and size  $\Omega(t^{\log 6/\log 4})$ , constructively proving that  $R(3, t) \in \Omega(t^{\log 6/\log 4})$ .

DEFINITION 1. For any graph G = (V(G), E(G)), the fibration of G is the graph H = (V(H), E(H)) defined below. Roughly speaking, H consists of six disjoint copies of G with extra edges that connect the vertices of different copies of G together. Formally,  $V(H) = V(G) \times \{0, 1, 2, 3, 4, 5\}$ , and E(H) consists of precisely the following edges.

- 1. For all  $i \in \{0, 1, 2, 3, 4, 5\}$ , and for all  $(u, v) \in E(G)$ ,  $((u, i), (v, i)) \in E(H)$ .
- 2. For all  $i, j \in \{0, 1, 2, 3, 4, 5\}$  with  $j \equiv i + 1 \pmod{6}$ , and for all  $(u, v) \in E(G)$ ,  $((u, i), (v, j)) \in E(H)$  and  $((u, j), (v, i)) \in E(H)$ .
- 3. For all  $i, j \in \{0, 1, 2, 3, 4, 5\}$  with  $j \equiv i + 3 \pmod{6}$ , and for all  $u \in V(G)$ ,  $((u, i), (u, j)) \in E(H)$ .

For each  $i \in \{0, 1, 2, 3, 4, 5\}$ , define  $V_i = \{(u, i) \in V(H)\}$ . For any set  $U \subset V(H)$ , let H[U] be the subgraph of H induced by U. We define the projection  $\pi: V(H) \to V(G)$  as  $\pi(u, i) = u$ , for all  $(u, i) \in V(H)$ .

LEMMA 1. If G is triangle-free and H is the fibration of G then H is triangle-free.

**Proof.** Let (u, i), (v, j), and (w, k) be any three distinct vertices in H. If i = j = k then, since G is triangle-free, the three vertices do not form a triangle in H. Also, if  $\{i, j, k\}$  contains two values whose difference modulo 6 is more than 1 then, by the structure of H, the three vertices cannot form a triangle in H. The remaining cases to consider are those where  $\{i, j, k\}$  contains exactly two distinct values whose difference modulo 6 is 1. By the symmetry of H, no generality is lost if we assume that i = j = 0 and k = 1. If (u, 0) and (v, 0) are both joined to (w, 1) in H then, by Definition 1, (u, 0) and (v, 0) are both also joined to (w, 0). Thus, if (u, 0), (v, 0), and (w, 1) form a triangle in H then (u, 0), (v, 0), and (w, 0) form a triangle in H, which contradicts the fact that G is triangle-free.

DEFINITION 2. For a graph G, the independence number of G, denoted by  $\alpha(G)$ , is the size of the largest independent set in G.

LEMMA 2. If H is the fibration of G then  $|H| = 6 \cdot |G|$  and  $\alpha(H) \leq 4 \cdot \alpha(G)$ .

*Proof.* Obviously,  $|H| = 6 \cdot |G|$ . The interesting part is to show that  $\alpha(H) \leq 4 \cdot \alpha(G)$ .

Now, let S be an arbitrary independent set in H. We need to show that  $|S| \le 4 \cdot \alpha(G)$ . For each  $i \in \{0, 1, 2, 3, 4, 5\}$ , let  $S_i = \{(u, i) \in S\}$ .

We shall first show that

$$|S_0| + |S_1| \le \alpha(G) + |\pi(S_0) \cap \pi(S_1)|$$

$$|S_2| + |S_3| \le \alpha(G) + |\pi(S_2) \cap \pi(S_3)|$$

$$|S_4| + |S_5| \le \alpha(G) + |\pi(S_4) \cap \pi(S_5)|.$$

Fix  $i \in \{0, 2, 4\}$ . The equality

$$|\pi(S_i)| + |\pi(S_{(i+1) \bmod 6})| = |\pi(S_i) \cup \pi(S_{(i+1) \bmod 6})| + |\pi(S_i) \cap \pi(S_{(i+1) \bmod 6})|$$

follows trivially because set cardinality is a modular function. It is therefore sufficient to prove  $\pi(S_i) \cup \pi(S_{(i+1) \bmod 6})$  is an independent set in G, since this implies that  $|\pi(S_i) \cup \pi(S_{(i+1) \bmod 6})| \leq \alpha(G)$ . Clearly, both  $S_i$  and  $S_{(i+1) \bmod 6}$  are independent sets in  $H[V_i]$  and  $H[V_{(i+1) \bmod 6}]$  and, therefore,  $\pi(S_i)$  and  $\pi(S_{(i+1) \bmod 6})$  are independent sets in G. Now assume that some  $u \in \pi(S_i)$  is adjacent to some  $v \in \pi(S_{(i+1) \bmod 6})$  in G. Then, by Definition 1, there is an edge in E(H) joining  $(u,i) \in S_i$  and  $(v,(i+1) \bmod 6) \in S_{(i+1) \bmod 6}$ . This is a contradiction, since S is assumed to be independent. Therefore,  $\pi(S_i) \cup \pi(S_{(i+1) \bmod 6})$  is independent in G and thus the three inequalities hold.

Summing the above inequalities, we obtain

$$|S_0| + |S_1| + |S_2| + |S_3| + |S_4| + |S_5|$$

$$\leq 3 \cdot \alpha(G) + |\pi(S_0) \cap \pi(S_1)| + |\pi(S_2) \cap \pi(S_3)| + |\pi(S_4) \cap \pi(S_5)|.$$

Now, to complete the proof, it is sufficient to show that

$$|\pi(S_0) \cap \pi(S_1)| + |\pi(S_2) \cap \pi(S_3)| + |\pi(S_4) \cap \pi(S_5)| \le \alpha(G).$$

First, note that  $\pi(S_0) \cap \pi(S_1)$ ,  $\pi(S_2) \cap \pi(S_3)$ , and  $\pi(S_4) \cap \pi(S_5)$  are mutually exclusive. To see why, suppose, without loss of generality, that  $v \in \pi(S_0) \cap \pi(S_1)$  and  $w \in \pi(S_2) \cap \pi(S_3)$ . Then, in particular,  $v \in \pi(S_0)$  and  $w \in \pi(S_3)$ . Thus, if v = w then, by the structure of H,  $((v, 0), (w, 3)) \in H$ , contradicting the fact that S is independent in H.

Furthermore,  $(\pi(S_0) \cap \pi(S_1)) \cup (\pi(S_2) \cap \pi(S_3)) \cup (\pi(S_4) \cap \pi(S_5))$  is an independent set. This follows from the observation that if, without loss of generality,  $v \in \pi(S_0) \cap \pi(S_1)$  and  $w \in \pi(S_2) \cap \pi(S_3)$  then, in particular,

 $v \in \pi(S_1)$  and  $w \in \pi(S_2)$ . Thus, if  $(v, w) \in G$  then  $((v, 0), (w, 1)) \in H$ , contradicting the fact that S is independent in H.

Therefore, since  $\pi(S_0) \cap \pi(S_1)$ ,  $\pi(S_2) \cap \pi(S_3)$ , and  $\pi(S_4) \cap \pi(S_5)$  are mutually exclusive and their union is independent in G,

$$|\pi(S_0) \cap \pi(S_1)| + |\pi(S_2) \cap \pi(S_3)| + |\pi(S_4) \cap \pi(S_5)| \le \alpha(G),$$

which completes the proof.

Theorem 3. There exists a feasible method for constructing a triangle-free graph with independence number less than t, whose size is greater than  $\frac{5}{6}((t-1)/2)^{\log 6/\log 4}$ . This constructively proves that  $R(3,t) > \frac{5}{6}((t-1)/2)^{\log 6/\log 4} \in \Omega(t^{1.29})$ .

*Proof.* Construct a sequence of graphs  $G_0$ ,  $G_1$ ,  $G_2$ , ... as follows. Let  $G_0$  be a 5-cycle, and let  $G_{i+1}$  be the fibration of  $G_i$ .  $G_0$  is triangle-free so, by Lemma 1, for all i,  $G_i$  is triangle-free. Clearly,  $|G_0| = 5$  and  $\alpha(G_0) = 2$ . By Lemma 2, for all i,  $|G_{i+1}| = 6 \cdot |G_i|$  and  $\alpha(G_{i+1}) \leqslant 4 \cdot \alpha(G_i)$ . Therefore, for all i,  $|G_i| = 5 \cdot 6^i$  and  $\alpha(G_i) \leqslant 2 \cdot 4^i$ . If the sequence is constructed until  $i = \lfloor \log((t-1)/2)/\log 4 \rfloor$  then  $\alpha(G_i) \leqslant t-1$  and  $|G_i| > \frac{5}{6}((t-1)/2)^{\log 6/\log 4}$ . ▮

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