BILL, RECORD LECTURE!!!!

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Ramsey Fails for $\ensuremath{\mathbb{R}}$

Exposition by William Gasarch

March 29, 2025

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Ramsey Over the Reals

We restate Ramsey's Theorem over $\ensuremath{\mathbb{N}}$ in a different way.

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Here is natural generalization to \mathbb{R} : **Conjecture** $\forall \text{COL} \colon \binom{\mathbb{R}}{2} \to [2] \exists$ homog set H such that $|H| = |\mathbb{R}|$. We will show that this Conj is **false**. The proof uses AC by using WOP.

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- **The Reals** Let ω_1 be the least uncountable ordinal.
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- By WOP there is a function $\mathbb{R} \to \omega_1$.
- This map induces a well-ordering \prec on $\mathbb R$
- ${\mathbb R}$ can be well ordered. Is that strange?

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Odd Fact 2: $(\forall x \in \mathbb{R})(\exists x^+)$ such that $x \prec x^+$ and $\neg \exists y[x \prec y \prec x^+]$.

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Odd? Do these two odd facts make your doubt WOP?

Lets look at

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Ramsey over $\mathbb R$ Does not hold

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See next page for how to make Gary happy, which will also make Meatloaf happy.

Making Gary Happy

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2) Work in ZF, not ZFC. Perhaps add some other axioms. I think there has been some work on this but its not stated this way.

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CON You call that an axiom? (TELL STORY)

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So far ZFC has not been able to show that X does not exist. Most set theorists think that $ZFC + \exists X$ is consistent.

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CON Is LC1 so obvious as to be a axiom?

Back to Ramsey Theory

Def A Ramsey Cardinal (RC) X is such that if $\forall \text{COL}: {X \choose 2} \rightarrow [2] \exists \text{ homog } H, |H| = |X|.$

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Thm If X is a RC then X is inaccessible. Hence we cannot prove RC's exist.