

BILL, RECORD LECTURE!!!!

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Kruskal's Tree Theorem and Two More Fast Growing Functions

Exposition by William Gasarch

February 27, 2025

Today's Lesson

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Warning Part of this talk will be on the whiteboard.

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Note If (X, \preceq) is a wqo then its both well founded and has no infinite antichains.

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Will use this lemma without pointing it out.

wqo's Closed Under \times

Def If (X, \preceq_1) and (Y, \preceq_2) are wqo then we define \preceq on $X \times Y$ as $(x, y) \preceq (x', y')$ if $x \preceq_1 x'$ and $y \preceq_2 y'$.

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If color has DOWN or INC in it then violates wqo.

The color must be UP-UP. This shows that there is an infinite ascending sequence.

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THIS IS SIMILAR TO THE KTT. MIGHT BE A HW.

tree: A Fast Growing Functions

HW Proof that \exists function $\text{tree}(n)$ such that:
 $\text{tree}(n)$ is largest number such that \exists a sequence of trees
 $T_1, T_2, \dots, T_{\text{tree}(n)}$
with the following properties.

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Known:

$$\text{tree}(1) = 1 \quad \text{tree}(2) = 5$$

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Known:

$$\text{tree}(1) = 1 \quad \text{tree}(2) = 5 \quad \text{tree}(3) \geq 844,424,930,181,960$$

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tree grows **much faster** than Ackermann's function.

TREE: A Faster Growing Functions

HW Show that there is a function $\text{TREE}(n)$ such that the following holds:

$\text{TREE}(n)$ is the largest number such that there exists a sequence of n -colored trees

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Known:

$\text{TREE}(1) = 1$ $\text{TREE}(2) = 3$ $\text{TREE}(3)$: See Next Slide.

tree vs TREE

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$\text{tree}^{\text{tree}^{\text{tree}^{\text{tree}^{\text{tree}^8(7)(7)(7)(7)(7)}}}}$

Suffice to say that $TREE(n)$ grows much faster than $\text{tree}(n)$.

TREE

TREE might be the faster growing **natural** function in mathematics.