BILL START RECORDING

HW 09 Solutions

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We do this problem without using O-of. We will assume t is large.

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 $A \rightarrow BC$ Number of rules of this type: t^3 .

 $A \rightarrow \sigma$ Number of rules of this type: 2t.

 $A \rightarrow e$ Number of rules of this type: t.

Number of rules: $t^5 + t^4 + t^3 + 2t + t \le 2t^5$.

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There are $\leq 2t^5$ rules.

Hence there are $\leq 2^{2t^5}$ grammars.

We let $g(t) = 2^{2t^5}$.

Find a function h such that the following is true: $\exists w \in \{a,b\}^*$ of length $\leq h(t)$ such that there is NO Leo-Grammar G with t nonterminals such that $L(G) = \{w\}$.

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Hence there must be some w of length $3t^5$ that is not generated by ANY Leo grammar with t NTs.

Problem 2a

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$$\Sigma = \{a, b, c, d\}.$$

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 $AB \rightarrow BA$

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 $BA \rightarrow AB$
 $AC \rightarrow CA$
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 $BC \rightarrow CB$

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$$S o SABC \mid e$$
 $AB o BA$
 $BA o AB$
 $AC o CA$
 $CA o AC$
 $BC o CB$
 $CB o BC$
 $A o a$

 $B \rightarrow b$

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$$C \rightarrow c$$

Number of rules: 2 + 6 + 3 = 11.



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S \rightarrow SA_1A_2A_3\cdots A_n \mid e
(for future: 2 rules)
For all \leq i \leq j \leq n
A_iA_i \rightarrow A_iA_i.
A_iA_i \rightarrow A_iA_n.
(for future: 2\binom{n}{2} = n(n-1) rules)
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A_i \rightarrow a_i
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Problem 2c

Find a function r such that your CSG from the last part has $\leq r(n)$ rules.

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Let $r(n) = n^2 + 2$.