

# **CMSC330 Spring 2015 Final Exam**

## **9:30am/11:00am/12:30pm**

Name:

UID:

**Discussion Time (circle one):**      10am    11am    12pm    1pm    2pm    3pm  
**Discussion TA (circle one):** Amelia Casey Chris Mike Elizabeth Eric Tommy

### **Instructions**

- The exam has 15 pages; make sure you have them all.
- Do not start this test until you are told to do so!
- You have 120 minutes to take this midterm.
- This exam has a total of 130 points, so allocate 55 seconds for each point.
- This is a closed book exam. No notes or other aids are allowed.
- Answer essay questions concisely in 2-3 sentences. Longer answers are not needed.
- For partial credit, show all of your work and clearly indicate your answers.
- Write neatly. Credit cannot be given for illegible answers.

	Problem	Score
1	PL Concepts	/22
2	Lambda Calculus	/6
3	OCaml Types	/9
4	OCaml Execution	/9
5	OCaml Programming	/14
6	Ruby Programming	/15
7	Prolog Execution	/15
8	Prolog Programming	/10
9	Multithreading in Ruby	/18
10	Regexp, FAs, CFGs	/12
	<b>TOTAL</b>	/130

**1. PL Concepts (22 pts)**

- a. (3 pts) Describe how the execution of the following OCaml program would differ with Call-by-Name parameter passing, compared to Call-by-Value:

```
let f x y =
  if x>0 then (y,y) else (1,1);;
f 1 (print_string "hello"; 2);;
```

**Answer: With CBV it would print the string “hello” and return (2,2), but CBN would not, simply returning the tuple ((print\_string “hello”; 2), (print\_string “hello”; 2)). Alternatively, you could also say that CBN would print “hello” twice, and return (2,2) for full credit.**

- b. (3 pts) Circle the following memory errors that can happen in a Ruby program.

*Dangling pointer dereference*      *Double-free*      **Memory leak**

- c. (4 pts) Circle the following languages below that have (mostly) static typing.

**OCaml**      **Ruby**      **Prolog**      **Java**

- d. (4 pts) Draw or describe the closure returned by calling (g 0), where g is defined as follows:

```
let g n =
  let r = ref 0 in
  let f m = !r + n + m in
  let x = f 0 in
  f
```

**Answer: the closure contains the code for f, and its environment contains the mapping n=0 and r=ref 0; it should not contain a mapping for x (or m).**

- e. (2 pts) True or false: if  $e_1$  and  $e_2$  are physically equal, then they are also structurally equal. (You may optionally justify your answer.)

**Answer: true**

- f. (6 pts) Consider the following Java code. Circle (and label) an example each of the use of *ad hoc polymorphism*, *subtype polymorphism*, and *parametric polymorphism* (generics) in the code.

```
public interface Foo {  
    public void shout();  
}  
  
public class Bar implements Foo {  
    public void shout() { System.out.println("foo!"); }  
    public void shout(String exclaim) {  
        System.out.println("foo says "+exclaim+"!");  
    }  
  
    public static <T> int count(Collection<T> stuff) {  
        int count = 0;  
        for (T o : stuff) { count++; }  
        return count;  
    }  
    public static void main(String args[]) {  
        Collection<Foo> c = new ArrayList<Foo>(10);  
        Bar b = new Bar();  
        b.shout();  
        b.shout("again");  
        c.add(b);  
        System.out.println(count(c));  
    }  
}
```

**Answer: Ad hoc is yellow, Parametric is purple, and subtyping is blue**

Notes:  $b.shout()$  is not subtyping, since  $b$  has type  $Bar$  (not  $Foo$ ).  $System.out.println$ , alone, is arguably *ad hoc polymorphism*, since there are multiple  $println$  functions.

**2. Lambda Calculus (6 pts).**

- a. (3 pts) What do you get if you fully reduce the following lambda calculus term?

$(\lambda x.x) (\lambda x.\lambda y.x) w z$

**Answer: w**

- b. (3 pts) Which of the following pairs of terms are alpha equivalent?

i.  $(\lambda x.\lambda y.x y) y z$        $(\lambda x.\lambda y.x y) w z$

ii.  $(\lambda x.\lambda y.x y) y z$        $(\lambda x.\lambda w.x w) y z$  ← **Answer**

iii.  $(\lambda x.\lambda y.x y) y z$        $(\lambda x.\lambda y.x x) y z$

**3. OCaml typing (9 pts).**

- a. (3 pts) What is the type of f in the following definition?

```
let f x = (x 1)+1
```

**Answer: (int -> int) -> int**

- b. (3 pts) What is the type of the following OCaml expression?

```
(fun x -> fun y -> y::x)
```

**Answer: 'a list -> 'a -> 'a list**

- c. (3 pts) What is the type of f in the following definition?

```
let f x a b = if x>0 then (a x) else (b x)
```

**Answer: int -> (int -> 'a) -> (int -> 'a) -> 'a**

4. **OCaml execution (9 pts):** What is the value of the variable `result` after executing the following code? If an error will occur, say what the problem is. (Note there are no syntax errors in these programs.)

a. (3 pts)

```
let rec aux m x =
  match m with
    [] -> []
  | h::t ->
    let t' = if h > x then [] else aux t x in
    h::t'
let result = aux [1;2;4;7;7;0] 4;;
```

**Answer:** [ 1; 2; 4; 7 ]

b. (3 pts)

```
type t = Int of int | String of string
let f x =
  match x with
    Int i -> string_of_int (i+1)
  | String _ -> "2";;
let result = (f (String "hi"), f (Int 1))
```

**Answer:** ("2","2")

c. (3 pts)

```
let co f g x =
  let z = f x in
  g z;;
let result = co (fun x -> x + 1) (fun y -> 0) 2
```

**Answer:** 0

5. **OCaml coding (14 pts).** For the following you may write helper functions if you like. You may also use standard library functions, e.g., from the List and Pervasives modules, unless otherwise noted.

- a. (4 pts) Write a function flat of type

```
(‘a -> ‘b -> ‘c) -> (‘a * ‘b) -> ‘c.
```

This function takes a curried function as its argument and returns a function that takes the arguments as a tuple instead. So:

```
let f = flat (fun x y -> x + y);; // f has type int*int -> int  
let result = f (1,2);; // result contains 3
```

The answer is actually very simple (one line): Don't overthink it!

**Answer:**

```
let flat f (x,y) = f x y
```

- b. (10 pts) **Either do this problem or the next one. If you do both we'll average your scores.** Write a function splast of type ‘a list -> ‘a list \* ‘a which which when given a non-empty list I will return a pair where the first component of the pair is the same as the input list but missing the last element, and the second component of the pair is the last element of the input list.

Examples:

```
splast [0;3] = ([0],3)  
splast [1;2;3;4] = ([1;2;3],4)
```

**Answer:**

```
let rec splast l =  
  match l with  
    [x] -> ([],x)  
  | h::t ->  
    let (m,x) = splast t in (h::m,x)
```

- c. (10 pts) **Either do this problem or the previous one. If you do both we'll average your scores.** Using either map or fold (code below) and an anonymous function, write a curried function partialProduct which when given a list of ints lst and an integer x, returns the product of each element in lst which is greater than x. The partial product of an empty list is 1. (If you use a recursive function and not map/fold you can earn at most 80% partial credit).

```
let rec map f l = match l with
  [] -> []
  | (h::t) -> (f h)::(map f t)
```

```
let rec fold f a l = match l with
  [] -> a
  | (h::t) -> fold f (f a h) t
```

Examples:

```
partialProduct [ ] 0 = 1
partialProduct [ ] 5 = 1
partialProduct [0;3] 1 = 3
partialProduct [1;2;3;4] 2 = 12
```

**Answer:**

```
let partialProduct lst x =
  fold (fun a m -> if m>x then m*a else a) 1 lst
```

6. **Ruby execution (15 pts).** Give the output of executing the following programs. If executing the program produces a failure of some sort, write the output up to that failure, and then write FAIL.

a. (3 pts)

```
a = "CMSC330"  
if a =~ /([\d]+)/ then  
    puts $1  
end
```

**Answer:** 330

b. (3 pts)

```
a = {}  
a[1] = "me"  
puts a[0]
```

**Answer:** (empty line, for nil)

c. (3 pts)

```
x = [3, 5, 8].map { |num| num * 2}  
puts x
```

**Answer:** 6  
10  
16

(Ok if all on same line)

d. (3 pts)

```
x = [["harry",6],["tom",3],["jack",9]]  
x.sort! { |x,y| x[1] <=> y[1]}  
puts x
```

**Answer:** tom  
3  
harry  
6  
jack  
9

(Ok if all on the same line)

e. (3 pts)

```
h = Hash.new(100)  
h["a"] = 10  
h["b"] = 20  
puts h["a"]  
puts h["c"]
```

**Answer:** 10  
100

**7. Prolog Execution (15 pts).**

- a. (6 pts total; 2 pts each.) Consider the following definitions. List all answers returned by the following queries (including substitutions of R that make the query true). If there is more than one answer, list them all; if there is no answer, write false. Answers must be listed in the correct order to receive full credit.

```
enrolled(steve,cmsc330).
enrolled(bill,cmsc351).
enrolled(alice,cmsc330).
isstudent(P) :- enrolled(P,_).
```

```
married(joe,alice).
married(bill,jane).
```

```
ismarried(P) :- married(P,_).
ismarried(P) :- married(_,P).
```

i.    enrolled(R,cmsc330)

**Answer: R = steve, R = alice**

ii.    isstudent(R).

**Answer: R = steve, R = bill, R = alice**

iii.    ismarried(steve).

**Answer: false**

- b. (3 pts) Suppose we extend the program above with the following definitions. What is returned by the query `ummarried_student(R)`?

```
hah(X) :- X, !, fail.
hah(_).
```

```
ummarried_student(X) :-
    isstudent(X),
    hah(ismarried(X)).
```

**Answer: R = steve**

- c. (3 pts) What if we changed `ummarried_student` to be defined as follows; would that change the result of the above `ummarried_student(R)` query? Why or why not?

```
unmarried_student(X) :-  
    hah(ismarried(X)),  
    isstudent(X).
```

**Answer: Yes. Since X would not be resolved prior to being used in hah, so it would be resolved there so that ismarried(X) succeeds, but this then causes hah(ismarried(X)) to fail overall.**

- d. (3 pts) Consider the following definitions. If check(L) is true, what can we say about the contents of L (explain in words)?

```
splast([X],[],X).  
splast([H|T],[H|T1],X1) :-  
    splast(T,T1,X1).
```

```
check([]).  
check([X]).  
check([H|T]) :-  
    splast(T,T1,H),  
    check(T1).
```

**Answer: L is a palindrome (a list that is its own mirror image)**

**8. Prolog Programming (10 pts).** Do either the next two problems. If you do both, we'll average the two scores.

- a. Define predicate `drop(L,X,P)` which holds when `P` is the same as the list `L` but with all occurrences of `X` in `L` removed. Thus, the following should hold:

```
drop([1,2],1,[2]).  
drop([1,1,2],1,[2]).  
drop([2,1,1],5,[2,1,1]).
```

**Answer:**

```
drop([],X,[X]).  
drop([H|T],H,T1) :-  
    drop(T,H,T1).  
drop([H|T],X,[H|T1]) :-  
    X \= H,  
    drop(T,X,T1).
```

- b. Define predicate `exp(N,M,P)` which holds when `P` is  $N^M$ . (You may assume that `M` is non-negative.) Thus, the following should hold:

```
exp(2,3,8).      %  $2^3 = 8$   
exp(3,2,9).      %  $3^2 = 9$  etc.
```

**Answer:**

```
exp(N,0,1).  
exp(N,X,P) :-  
    X > 0,  
    X1 is X-1,  
    exp(N,X1,P1),  
    P is N * P1.
```

9. **Multithreading in Ruby (18 points).**

- a. (8 pts) Write all possible outputs from executing the following multithreaded Ruby program.

```
def func1
    i=2
    while i<=3
        sleep(rand())
        puts "func1: #{i}"
        i=i+1
    end
end

def func2
    j=0
    while j<=1
        sleep(rand())
        puts "func2: #{j}"
        j=j+1
    end
end

puts "Start"
t1=Thread.new{func1()}
t2=Thread.new{func2()}
t1.join
t2.join
puts "End"
```

**Answer:** There are 6 possible output sequences, basically capturing the combination of possible interleavings between the two threads.

- b. (10 pts.) The following code, implementing a thread-shared, fixed-sized Stack, has two problems: it has data races and it uses busy-waiting. Fix the code by using proper synchronization and condition variables. Write your changes to the side, clearly marking what you are inserting/deleting/changing. You may use the library functions listed here:

```
m = Monitor.new          // creates a new monitor
m.synchronize { ... }    // acquires m, enters the block, releases m
c = m.new_cond           // returns new condition variable (condvar)
for m
c.wait_while { ... }     // waits on condvar c while code block is true
c.wait_until { ... }      // waits on condvar c until code block is true
c.broadcast              // signals all threads waiting on condvar c
```

```
class Stack
  def initialize(sz)
    @idx = 0
    @buf = Array.new(sz)
  end

  def push(x)
    while @idx >= (@buf.size - 1) do
      # wait
    end
    @buf[@idx] = x
    @idx = @idx + 1
  end

  def pop
    while @idx == 0 do
      # wait
    end
    @idx = @idx - 1
    x = @buf[@idx]
    @buf[@idx] = nil
    x
  end
end
```

Example: s = Stack.new(2)  
t1 = Thread.new { s.push("hi"); x = s.pop; puts x }  
t2 = Thread.new { s.push("there"); y = s.pop; puts y }  
t1.join; t2.join  
*# should print hi\n there\n or there\n hi\n*

**Answer (stuff changed/added from what was given is in red):**

```
class Stack
  def initialize(sz)
    @idx = 0
    @buf = Array.new(sz)
    @m = Monitor.new
    @c = @m.new_cond
  end
  def push(x)
    @m.synchronize {
      @c.wait_while { @idx >= (@buf.size - 1) }
      # while @idx >= (@buf.size - 1) do
      # end
      @buf[@idx] = x
      @idx = @idx + 1
      @c.broadcast
    }
  end
  def pop
    @m.synchronize {
      @c.wait_while { @idx == 0 }
      # while @idx == 0 do
      # end
      @idx = @idx - 1
      x = @buf[@idx]
      @buf[@idx] = nil
      @c.broadcast
      x
    }
  end
end
```

**10. Regular expressions, Context-free Grammars, Automata (12 pts).**

- a. (3 pts) Write a context-free grammar that accepts all strings  $a^n b^m$  where  $n = 2+2m$ , for all  $m > 0$ . For example, the strings aaaab, aaaaaabb, and aaaaaaaaaabbb are all in the language.

**Answer:**

$$S \rightarrow aaT$$

$$T \rightarrow aaTb \mid aab$$

- b. (3 pts) Give a regexp that denotes the language of all strings over  $\{a,b\}$  which start with a, end with b, and contain the substring abb.

**Answer:**  $a(a|b)^*abb(a|b)^*b \mid abbb^* \mid aaa^*bb$

- c. Circle all of the inputs accepted by the NFA below (1 pt each)

- i.  $\epsilon$
- ii. aba       $\leftarrow$  Accepted
- iii. aaaaaaaaa       $\leftarrow$  Accepted
- iv. abab       $\leftarrow$  Accepted
- v. aaaab       $\leftarrow$  Accepted
- vi. aa

