# Universal computation by multi-particle quantum walk

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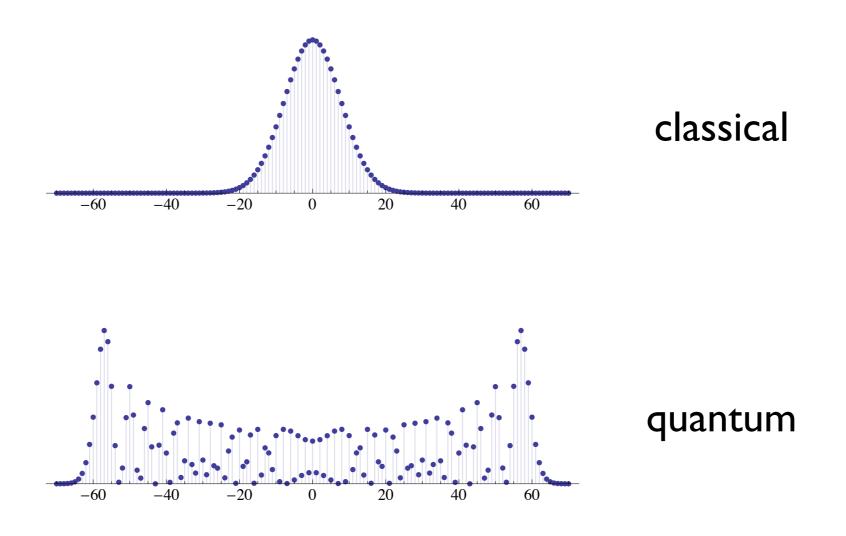
# UNIVERSITY OF WATERLOO

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#### Quantum walk

Quantum analog of a random walk on a graph.

Idea: Replace probabilities by quantum amplitudes. Interference can produce radically different behavior!



## Quantum walk algorithms

Quantum walk is a major tool for quantum algorithms (especially query algorithms with polynomial speedup).

- Exponential speedup for black-box graph traversal [CCDFGS 02]
- Quantum walk search framework [Szegedy 05], [Magniez et al. 06]
  - Spatial search [Shenvi-Kempe-Whaley 02], [CG 03, 04], [Ambainis-Kempe-Rivosh 04]
  - Element distinctness [Ambainis 03]
  - Subgraph finding [Magniez, Santha, Szegedy 03], [CK 10]
  - Matrix/group problems [Buhrman, Špalek 04], [Magniez, Nayak 05]
- Evaluating formulas/span programs
  - AND-OR formula evaluation [Farhi, Goldstone, Gutmann 07], [ACRŠZ 07]
  - Span programs for general query problems [Reichardt 09]
  - Learning graphs [Belovs 11] → new upper bounds (implicitly, quantum walk algorithms), new kinds of quantum walk search

## Universality of quantum walk

Quantum walk can be efficiently simulated by a universal quantum computer.

N-vertex graph  $\max \text{ degree } \operatorname{poly}(\log N) \qquad \qquad \Rightarrow \quad \operatorname{poly}(\log N) \text{ qubits}$ efficiently computable neighbors poly(log N) gates

circuit with

Conversely, quantum walk is a *universal computational primitive*: any quantum circuit can be simulated by a quantum walk. [C 09]

		graph with
N dimensions g gates	$\Rightarrow$	$\operatorname{poly}(N,g)$ vertices
		max degree 3
		walk for time $poly(g)$

Note: The graph is necessarily exponentially large in the number of qubits! Vertices represent basis states.

#### Quantum walk experiments



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Quantum random walks are the quantum counterpart of classical random walks, and were recently

Quantum random warks are the quantum counterpart of classical random warks, and were recently studied in the context of quantum computation. Physical implementations of quantum walks have only been made in very small scale systems severely limited by decoherence. Here we show that the been made in very small scale systems severely innited by deconerence. Here we snow that the propagation of photons in waveguide lattices, which have been studied extensively in recent years, are propagation of photons in waveguide lattices, which have been studied extensively in recent years, are essentially an implementation of quantum walks. Since waveguide lattices are easily constructed at large essentiative an implementation of quantum walks. Since waveguide lattices are easily constructed at large scales and display negligible decoherence, they can serve as an ideal and versatile experimental scales and utsplay negligible deconference, they can serve as an ideal and versatile experiment playground for the study of quantum walks and quantum algorithms. We experimentally observe quantum in the second se playeround for the study of quantum warks and quantum algorithms, we experimentary observe quantum walks in large systems ( $\sim$  100 sites) and confirm quantum walks effects which were studied theoretically,

implemented, using linear optical elements [7]. For CQWs,

a few suggestions have been made [8,9], yet only one

experimental method have been implemented by realizing

a small scale cyclic system (4 states) using a nuclear

magnetic resonance system [10]. Such systems are difficult

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to scale to much larger configurations. Moreover, even at these very small scales, errors attributed to decoherence Here we suggest a very different implementation of

CQWs using optical waveguide lattices. These systems have been studied extensively in recent years [11], but not in the context of QWs and quantum algorithms. We show that these systems can serve as a unique and robust tool for the study of CQWs. For this purpose we demonstrate three fundamental QW effects that have been theoretically analyzed in the QW literature. These include ballistic propagation in the largest system reported to date (  $\sim$  100 sites), the effects of disorder on QWs, and QWs with reflecting boundary conditions (related to Berry's "particle in a box" and quantum carpets [12,13]). Waveguide lattices can be easily realized with even larger scales than shown here  $(10^2 - 10^4$  sites with current fabrication technologies), with practically no decoherence. The high level of engineering and control of these systems enable the study of a wide range of different parameters and initial conditions. Specifically it allows the implementation and study of a large variety of CQWs and show experimental observations of their unique behavior. The CQW model was first suggested by Farhi and Gutmann [6], where the intuition behind it comes from continuous time classical Markov chains. In the classical random walk on a graph, a step can be described by a matrix M which transforms the probability distribution for the particle position over the graph nodes (sites). The entries of the matrix  $M_{j,k}$  give the probability to go from site j to site k in one step of the walk. The idea was to carry

this construction over to the quantum case, where the Hamiltonian of the process is used as the generator matrix. The system is evolved using  $U(t) = \exp(-iHt)$ . If we start in some initial state  $|\Psi_{in}\rangle$ , evolve it under U for a time T and measure the positions of the resulting state, we obtain a 170506-1

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regard to energy transport; the asymmetry between the horizontal directions and the reduced overall energy flux reflect the constraints imappl posed on the convective motions by the presence of a strong and inclined magnetic field. The development of systematic outflows is a direct consequence of the anisotropy, and the similarities between granulation and penumbral flows strongly suggest that driving the Evershed flow does not require physical processes that go beyond the combination of convection and anisotropy intro-duced by the magnetic field. Weaker laterally overturning flows perpendicular to the main filament direction explain the apparent twisting motions observed in some filaments (15, 16) and lead to a weakening of the magnetic field in the flow channels through flux expulsion (6).

random walk

Galton's quin

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nterference phenomena with microscopic particles are a direct consequence of their uantum-mechanical wave nature (1-5). The prospect to fully control quantum properties of atomic systems has stimulated ideas to engineer quantum states that would be useful for applications in quantum information processing, for example, and also would elucidate fundamental questions, such as the quantum-to-classical transition (6). A prominent example of state engineering by controlled multipath interference is the quantum walk of a particle (7). Its classical

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aspects of our lives, providing insight fields: It forms the basis for algorith scribes diffusion processes in physic (8, 9), such as Brownian motion, used as a model for stock market Similarly, the quantum walk is exp

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counterpart, the random walk, is releve

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#### Multi-particle quantum walk

With many walkers, the Hilbert space can be much bigger. m distinguishable particles on an n-vertex graph:  $n^m$  dimensions (similar scaling for indistinguishable bosons/fermions)

Main result: Any *n*-qubit, *g*-gate quantum circuit can be simulated by a multi-particle quantum walk of n + 1 particles interacting for time poly(n,g) on a graph with poly(n,g) vertices.

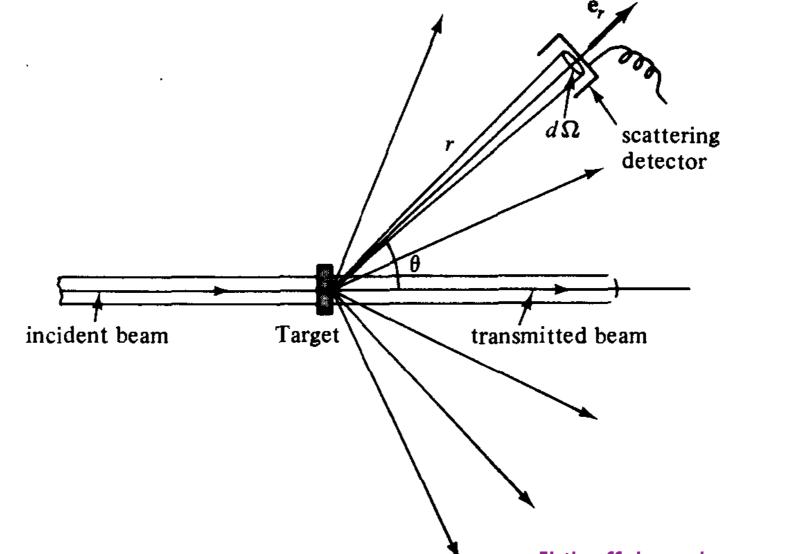
Consequences:

- Architecture for a quantum computer with no time-dependent control
- Simulating interacting many-body systems is BQP-hard (e.g., Bose-Hubbard model on a sparse, unweighted, planar graph)

# Outline

- Scattering theory on graphs
- Single-particle universality (review)
- Multi-particle universality
- Proof ideas
- Refinements and extensions
- Open questions

# Scattering theory on graphs



[Liboff, Introductory Quantum Mechanics]

#### Quantum walk

Quantum analog of a random walk on a graph G = (V, E).

Idea: Replace probabilities by quantum amplitudes.

$$\begin{split} \psi(t) \rangle &= \sum_{v \in V} a_v(t) |v\rangle \\ & \swarrow \\ & \checkmark \\ & \text{amplitude for vertex } v \text{ at time } t \end{split}$$

Define time-homogeneous, local dynamics on G.

$$\mathrm{i} \frac{\mathrm{d}}{\mathrm{d}t} |\psi(t)\rangle = H |\psi(t)\rangle$$

Adjacency matrix: 
$$H = \sum_{(u,v)\in E(G)} |u\rangle\langle v|$$

#### Momentum states

Consider an infinite path:

$$-7 \quad -6 \quad -5 \quad -4 \quad -3 \quad -2 \quad -1 \quad 0 \quad 1 \quad 2 \quad 3 \quad 4 \quad 5 \quad 6 \quad 7$$

Hilbert space:  $\operatorname{span}\{|x\rangle: x \in \mathbb{Z}\}$ 

Eigenstates of the adjacency matrix:  $|\tilde{k}\rangle$  with  $\langle x|\tilde{k}\rangle := e^{\mathrm{i}kx} \qquad k \in [-\pi,\pi)$ 

We have 
$$\langle x|A|\tilde{k}\rangle = \langle x-1|\tilde{k}\rangle + \langle x+1|\tilde{k}\rangle$$
  
$$= e^{ik(x-1)} + e^{ik(x+1)}$$
$$= (2\cos k)\langle x|\tilde{k}\rangle$$

so this is an eigenstate with eigenvalue  $2\cos k$ .

#### Wave packets

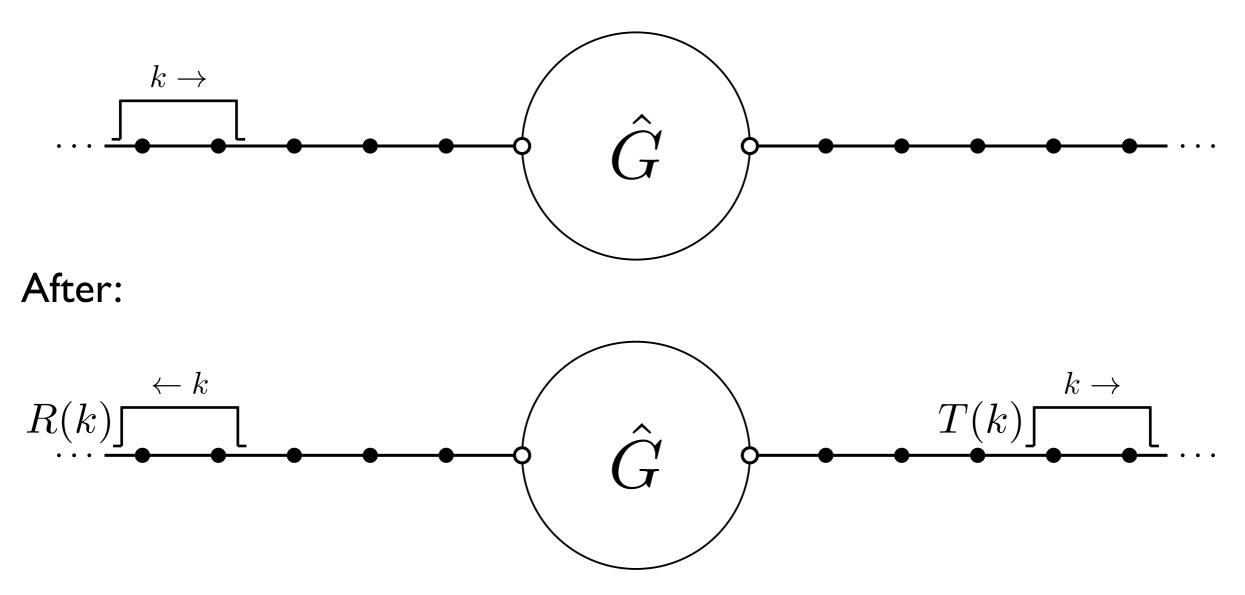
A wave packet is a normalized state with momentum concentrated near a particular value k.

Example: 
$$\frac{1}{\sqrt{L}} \sum_{x=1}^{L} e^{-ikx} |x\rangle$$
 (large L)  
 $\frac{k \rightarrow}{-1 \ 0 \ 1 \ 2 \ 3 \ 4 \ 5 \ 6 \ 7 \ 8 \ 9 \ 10 \ 11 \ 12 \ 13} \cdots$   
Propagation speed:  $\left|\frac{dE}{dk}\right| = 2|\sin k|$ 

## Scattering on graphs

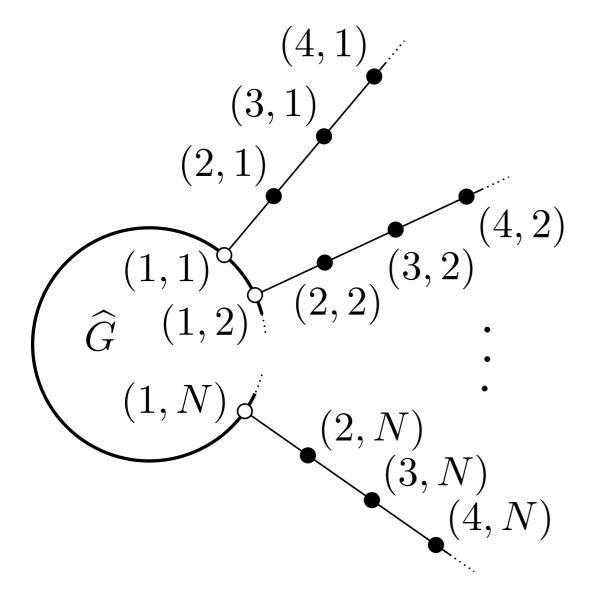
Now consider adding semi-infinite lines to two vertices of an arbitrary finite graph.

Before:



#### The S-matrix

This generalizes to any number N of semi-infinite paths attached to any finite graph.



Incoming wave packets of momentum near k are mapped to outgoing wave packets (of the same momentum) with amplitudes corresponding to entries of an  $N \times N$  unitary matrix S(k), called the S-matrix.

# Single-particle universality

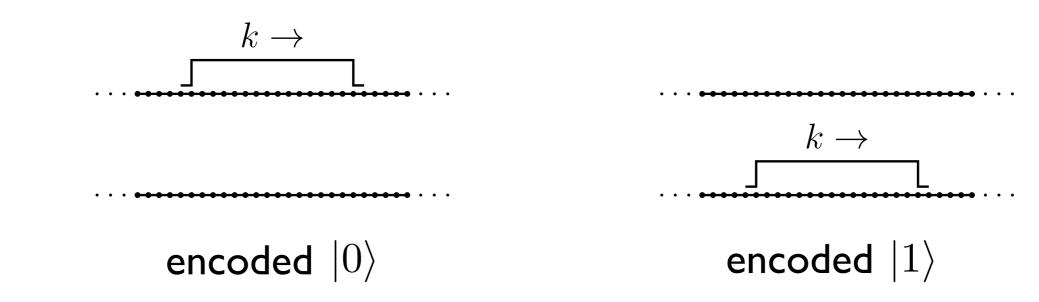
arXiv:0806.1972 Physical Review Letters 102, 180501 (2009)

## Encoding a qubit

Encode quantum circuits into graphs.

Computational basis states correspond to paths ("quantum wires").

For one qubit, use two wires ("dual-rail encoding"):

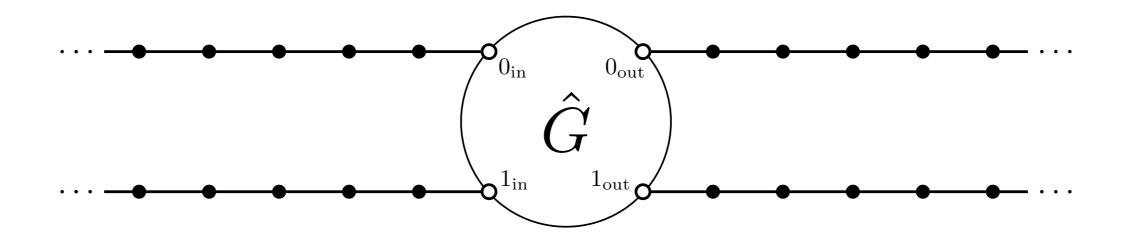


Fix some value of the momentum (e.g.,  $k = \pi/4$ ).

Quantum information propagates from left to right at constant speed.

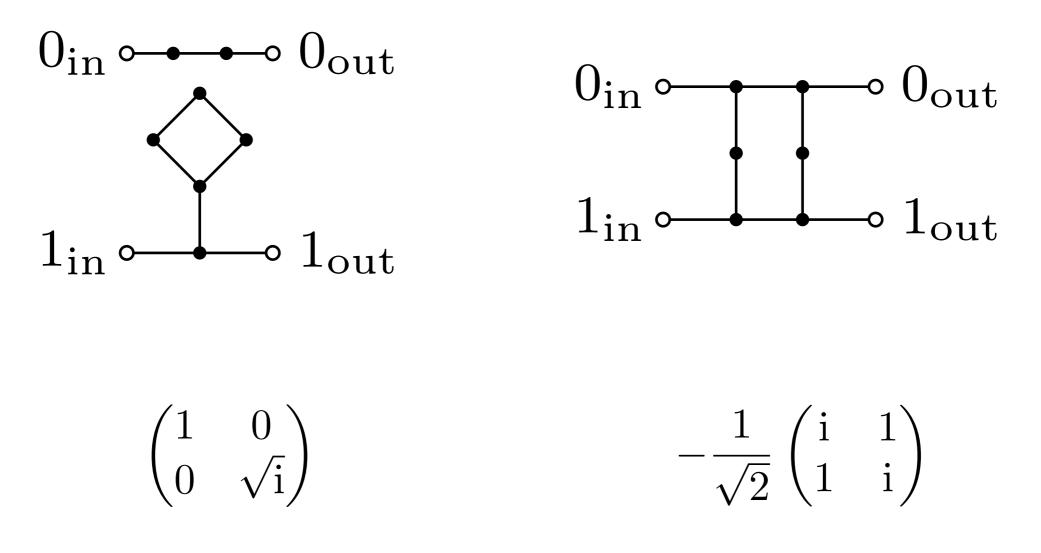
#### Implementing a gate

To perform a gate, design a graph whose S-matrix implements the desired transformation U at the momentum used for the encoding.



$$S(k) = \begin{pmatrix} 0 & V \\ U & 0 \end{pmatrix}$$

#### Universal set of single-qubit gates



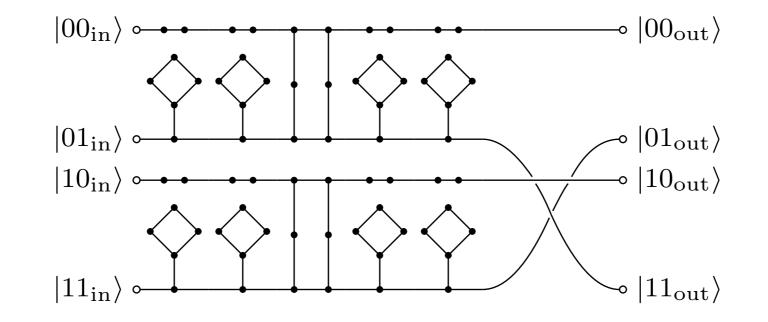
momentum for logical states:  $k = \pi/4$ 

#### Single-particle universality construction

With an appropriate encoding of *n*-qubit states, two-qubit gates are trivial.

Implement sequences of gates by concatenation.

Any n-qubit circuit can be simulated by some graph. The number of vertices is (necessarily) exponential in n.



# Multi-particle universality

#### Multi-particle quantum walk

With m distinguishable particles:

states: 
$$|v_1, \dots, v_m\rangle$$
  $v_i \in V(G)$   
Hamiltonian:  $H_G^{(m)} = \sum_{i=1}^m \sum_{(u,v)\in E(G)} |u\rangle\langle v|_i + \mathcal{U}$ 

Indistinguishable particles:

bosons: symmetric subspace fermions: antisymmetric subspace

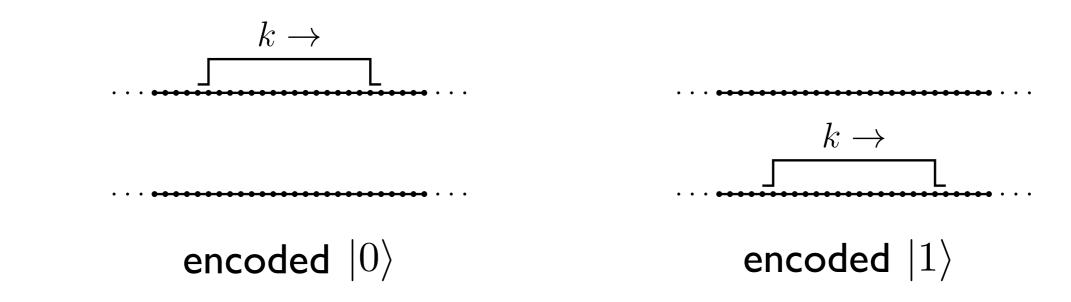
Many possible interactions:

$$\begin{array}{ll} \text{on-site:} & \mathcal{U} = J \sum_{v \in V(G)} \hat{n}_v (\hat{n}_v - 1) & \hat{n}_v = \sum_{i=1} |v\rangle \langle v|_i \\ \\ \text{nearest-neighbor:} & \mathcal{U} = J \sum_{(u,v) \in E(G)} \hat{n}_u \hat{n}_v \end{array}$$

m

# Logical encoding

Encode each qubit in a single-particle state as before:



For n qubits, use n particles on 2n paths.

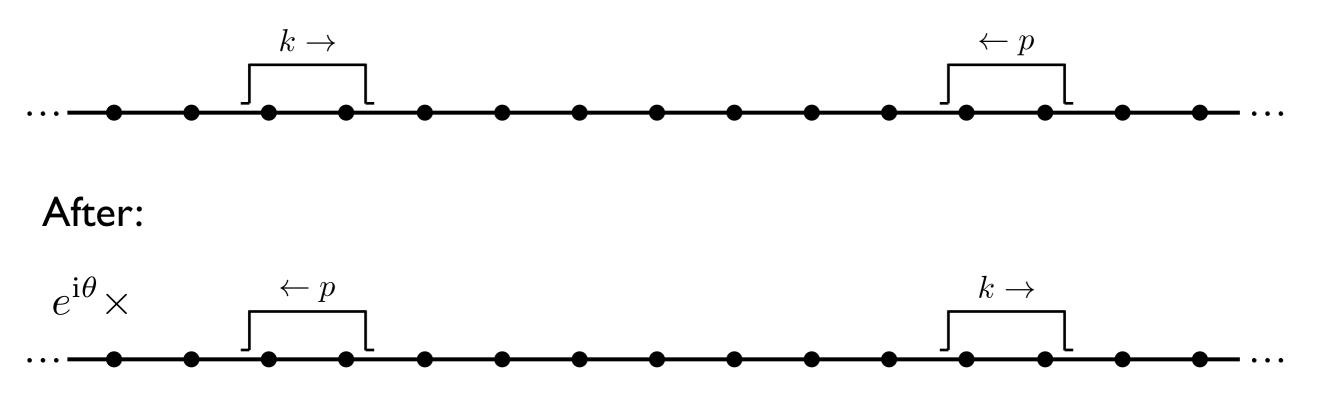
Perform single-qubit gates as before. How to perform interactions?

#### Two-particle scattering

In general, multi-particle scattering is complicated.

But scattering of indistinguishable particles on an infinite path is simple.

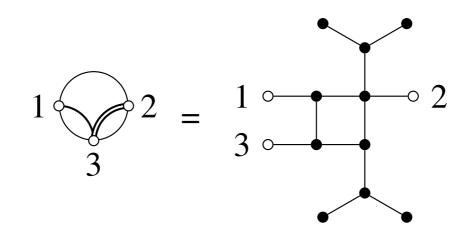
Before:



Phase  $\theta$  depends on momenta and interaction details.

#### Momentum switch

To selectively induce the two-particle scattering phase, we route particles depending on their momentum.

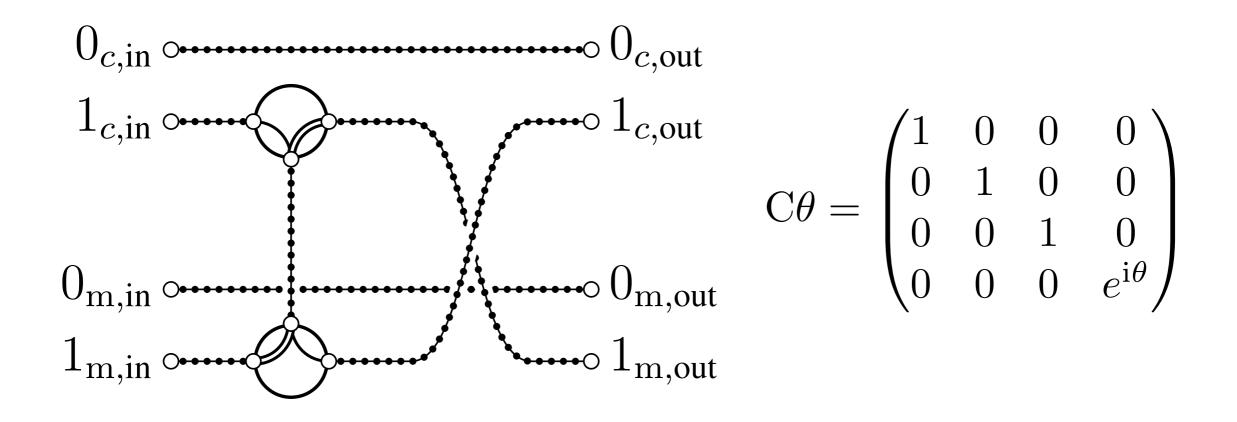


Particles with momentum  $\pi/4$  follow the single line.

Particles with momentum  $\pi/2$  follow the double line.

#### Controlled phase gate

Computational qubits have momentum  $\pi/4$ . Introduce a "mediator qubit" with momentum  $\pi/2$ . We can perform an entangling gate with the mediator qubit.

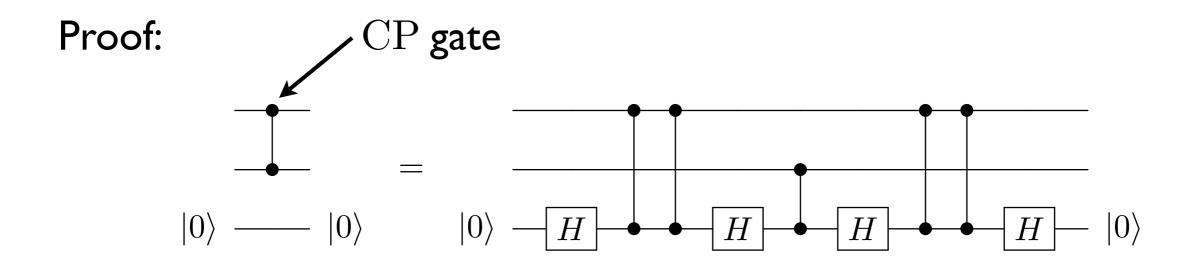


#### Canonical form of a circuit

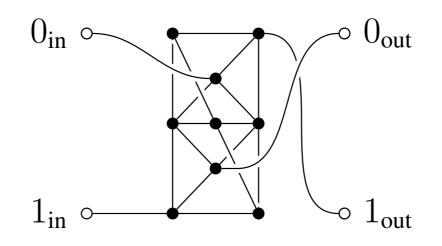
For almost all  $\theta$ , we can use  $C\theta$  to do  $CP = \begin{pmatrix} 1 & 0 & 0 & 0 \\ 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 \\ 0 & 0 & 0 & -i \end{pmatrix}$ .

Any circuit can be implemented using

- Single-qubit gates on computational qubits
- CP gates between a computational and the mediator qubit
- Hadamard gates on the mediator qubit

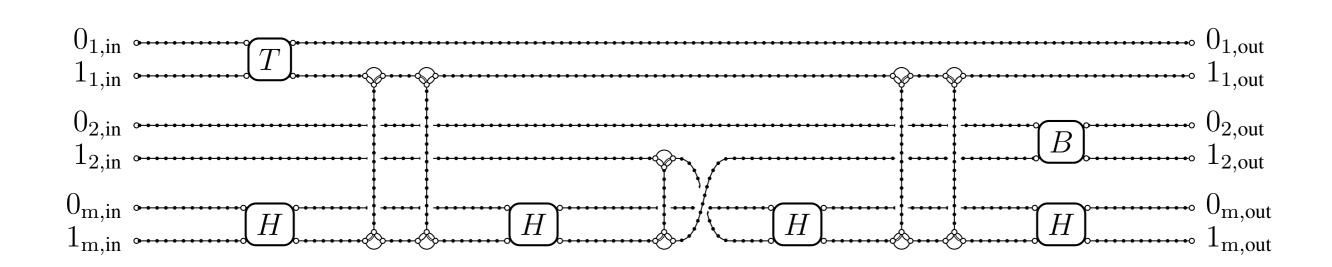


#### Hadamard on mediator qubit



[Blumer-Underwood-Feder 11]

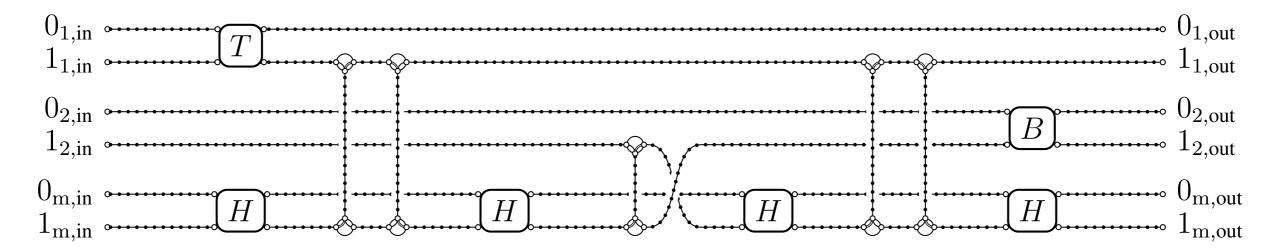
## Example



#### Error bound

Initial state: each particle is a square wave packet of length L

Consider a g-gate, n-qubit circuit:



2(n+1) paths, O(gL) vertices on each path Evolution time O(gL)Total # of vertices O(ngL)

Theorem: The error can be made arbitrarily small with L = poly(n, g).

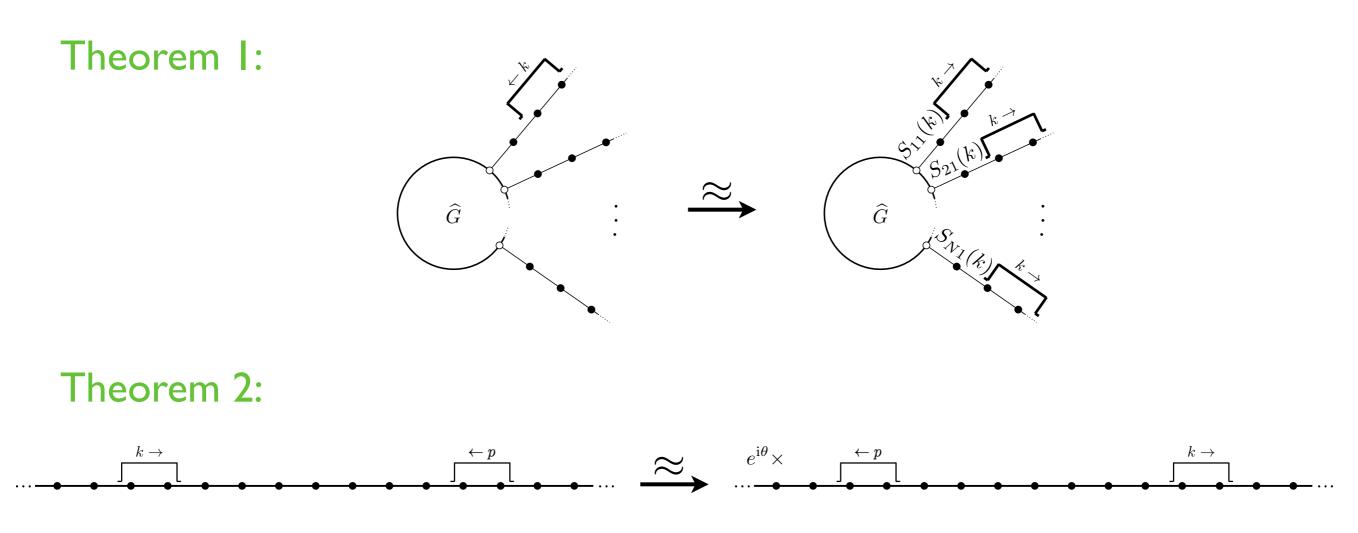
Example: For Bose-Hubbard model,  $L = O(n^{12}g^4)$  suffices.

**Proof ideas** 

#### Approximating wave packet scattering

Analysis of single-particle construction: method of stationary phase

Instead, we directly prove that, with long enough incoming square wave packets (length L), the outgoing wave packets are well-approximated by the effect of the S-matrix (error  $O(L^{-1/4})$ ).



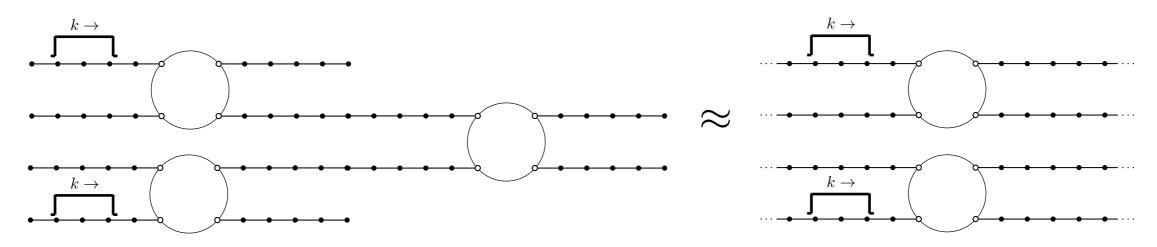
## Truncation

Truncation Lemma: If, for all times of interest, the state is wellapproximated by one that is well-localized to some region, then changing how the Hamiltonian acts outside that region has little effect.

This lets us

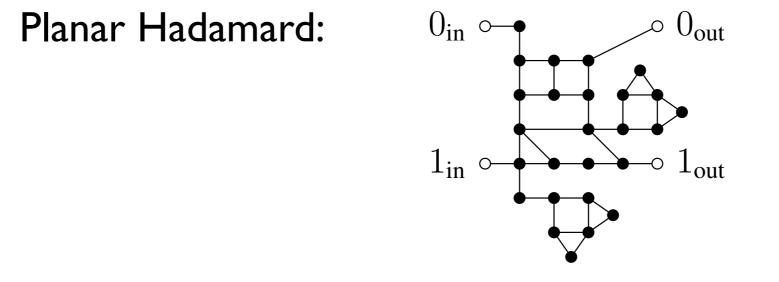
- approximate finite graphs by infinite ones
- approximate the evolution by piecewise scattering through separate gate gadgets
- focus exclusively on one- and two-particle scattering

Example: For small enough times,

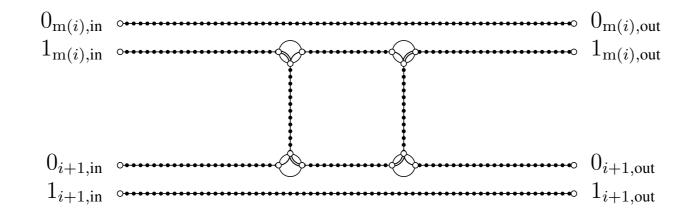


# **Refinements and extensions**





Put a mediator between every pair of computational qubits and unwind pairs of  $C\theta$  gates:



## Distinguishable particles

For two distinguishable particles on a long path, reflected and transmitted states are distinct:

Then the effect of scattering on a long path is not just to accumulate a phase.

But with an appropriate choice of parameters, can ensure (say) R = 0.

# **Open questions**

- Improved error bounds
- Simplified initial state
- Are generic interactions universal for distinguishable particles?
- New quantum algorithms
- Experiments
- Fault tolerance