

15-213

"The course that gives CMU its Zip!"

Machine-Level Programming II: Control Flow Sept. 12, 2002

Topics

- Condition Codes
 - Setting
 - Testing
- Control Flow
 - If-then-else
 - Varieties of Loops
 - Switch Statements

class06.ppt

Setting Condition Codes (cont.)

Explicit Setting by Compare Instruction

`cmp1 Src2,Src1`

- `cmp1 b,a` like computing $a-b$ without setting destination
- CF set if carry out from most significant bit
 - Used for unsigned comparisons
- ZF set if $a == b$
- SF set if $(a-b) < 0$
- OF set if two's complement overflow
 - $(a>0 \&& b<0 \&& (a-b)<0) \mid\mid (a<0 \&& b>0 \&& (a-b)>0)$

Condition Codes

Single Bit Registers

CF	Carry Flag	SF	Sign Flag
ZF	Zero Flag	OF	Overflow Flag

Implicitly Set By Arithmetic Operations

`addl Src,Dest`

C analog: $t = a + b$

- CF set if carry out from most significant bit
 - Used to detect unsigned overflow
- ZF set if $t == 0$
- SF set if $t < 0$
- OF set if two's complement overflow
 - $(a>0 \&& b>0 \&& t<0) \mid\mid (a<0 \&& b<0 \&& t>=0)$

Not Set by `leal` instruction

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Setting Condition Codes (cont.)

Explicit Setting by Test instruction

`testl Src2,Src1`

- Sets condition codes based on value of `Src1 & Src2`
 - Useful to have one of the operands be a mask
- `testl b,a` like computing $a \& b$ without setting destination
- ZF set when $a \& b == 0$
- SF set when $a \& b < 0$

Reading Condition Codes

SetX Instructions

- Set single byte based on combinations of condition codes

\$	SetX	Condition	Description
		je	ZF Equal / Zero
		jne	ZF Not Equal
		jl	ZF Less Than
		jge	ZF Greater Than or Equal
		jle	ZF Less Than or Equal
		jg	ZF Greater Than
		ja	ZF Above
		jna	ZF Above or Equal

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Reading Condition Codes (Cont.)

SetX Instructions

- Set single byte based on combinations of condition codes
- One of 8 addressable byte registers
 - Embedded within first 4 integer registers
 - Does not alter remaining 3 bytes
 - Typically use movzbl to finish job

```
int gt (int x, int y)
{
    return x > y;
}
```

Body

```
movl 12(%ebp),%eax # eax = y
cmpb %ah,8(%ebp) # Compare x : y
setg %al # al = x > y
movzbl %al,%eax # Zero rest of %eax
```

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%eax	%ah	%al
%edx	%dh	%dl
%ecx	%ch	%cl
%ebx	%bh	%bl
%esi		
%edi		
%esp		
%ebp		

Note
inverted
ordering!

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Jumping

jX Instructions

- Jump to different part of code depending on condition codes

j	X	Condition	Description
		jmp	Unconditional
		je	ZF Equal / Zero
		jne	ZF Not Equal
		jl	ZF Less Than
		jge	ZF Greater Than or Equal
		jle	ZF Less Than or Equal
		jg	ZF Greater Than
		ja	ZF Above
		jna	ZF Above or Equal

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Conditional Branch Example

```
max:
    pushl %ebp
    movl %esp,%ebp
    movl 8(%ebp),%edx
    movl 12(%ebp),%eax
    cmpl %eax,%edx
    jle L9
    movl %edx,%eax
L9:
    movl %ebp,%esp
    popl %ebp
    ret
```

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Conditional Branch Example (Cont.)

```
int goto_max(int x, int y)
{
    int rval = y;
    int ok = (x <= y);
    if (ok)
        goto done;
    rval = x;
done:
    return rval;
}
```

- C allows “goto” as means of transferring control
 - Closer to machine-level programming style
- Generally considered bad coding style

```
movl 8(%ebp),%edx    # edx = x
movl 12(%ebp),%eax  # eax = y
cmpl %eax,%edx      # x : y
jle L9                # if <= goto L9
movl %edx,%eax       # eax = x } Skipped when x ≤ y
L9:                  # Done:
```

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“Do-While” Loop Example

C Code

```
int fact_do
    (int x)
{
    int result = 1;
    do {
        result *= x;
        x = x-1;
    } while (x > 1);
    return result;
}
```

Goto Version

```
int fact_goto(int x)
{
    int result = 1;
loop:
    result *= x;
    x = x-1;
    if (x > 1)
        goto loop;
    return result;
}
```

- Use backward branch to continue looping
- Only take branch when “while” condition holds

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“Do-While” Loop Compilation

Goto Version

```
int fact_goto
    (int x)
{
    int result = 1;
loop:
    result *= x;
    x = x-1;
    if (x > 1)
        goto loop;
    return result;
}
```

Registers

```
%edx  x
%eax  result
```

Assembly

```
_fact_goto:
    pushl %ebp          # Setup
    movl %esp,%ebp      # Setup
    movl $1,%eax         # eax = 1
    movl 8(%ebp),%edx    # edx = x

    L11:
        imull %edx,%eax   # result *= x
        decl %edx           # x--
        cmpl $1,%edx         # Compare x : 1
        jg L11                # if > goto loop

        movl %ebp,%esp      # Finish
        popl %ebp             # Finish
        ret                  # Finish
```

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General “Do-While” Translation

C Code

```
do
    Body
    while (Test);
```

Goto Version

```
loop:
    Body
    if (Test)
        goto loop
```

- **Body** can be any C statement
 - Typically compound statement:

```
{
    Statement1;
    Statement2;
    ...
    Statementn;
}
```

- **Test** is expression returning integer
 - = 0 interpreted as false ≠ 0 interpreted as true

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“While” Loop Example #1

C Code

```
int fact_while
    (int x)
{
    int result = 1;
    while (x > 1) {
        result *= x;
        x = x-1;
    };
    return result;
}
```

First Goto Version

```
int fact_while_goto
    (int x)
{
    int result = 1;
loop:
    if (!(x > 1))
        goto done;
    result *= x;
    x = x-1;
    goto loop;
done:
    return result;
}
```

- Is this code equivalent to the do-while version?
- Must jump out of loop if test fails

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Actual “While” Loop Translation

C Code

```
int fact_while(int x)
{
    int result = 1;
    while (x > 1) {
        result *= x;
        x = x-1;
    };
    return result;
}
```

Second Goto Version

```
int fact_while_goto2
    (int x)
{
    int result = 1;
    if (!(x > 1))
        goto done;
loop:
    result *= x;
    x = x-1;
    if (x > 1)
        goto loop;
done:
    return result;
}
```

- Uses same inner loop as do-while version
- Guards loop entry with extra test

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General “While” Translation

C Code

```
while (Test)
    Body
```



Do-While Version



Goto Version

```
if (!Test)
    goto done;
do
    Body
    while (Test);
done:
```

```
if (!Test)
    goto done;
loop:
    Body
    if (Test)
        goto loop;
done:
```

“For” Loop Example

```
/* Compute x raised to nonnegative power p */
int ipwr_for(int x, unsigned p) {
    int result;
    for (result = 1; p != 0; p = p>>1) {
        if (p & 0x1)
            result *= x;
        x = x*x;
    }
    return result;
}
```

Algorithm

- Exploit property that $p = p_0 + 2p_1 + 4p_2 + \dots + 2^{n-1}p_{n-1}$
- Gives: $x^p = z_0 \cdot z_1^2 \cdot (z_2^2)^2 \cdot \dots \cdot (\dots((z_{n-1}^2)^2)\dots)^2$
- Complexity $O(\log p)$

$$z_i = 1 \text{ when } p_i = 0$$

$$z_i = x \text{ when } p_i = 1$$

$\underbrace{\quad}_{n-1 \text{ times}}$

Example

$$\begin{aligned} 3^{10} &= 3^2 * 3^8 \\ &= 3^2 * ((3^2)^2)^2 \end{aligned}$$

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ipwr Computation

```
/* Compute x raised to nonnegative power p */
int ipwr_for(int x, unsigned p) {
    int result;
    for (result = 1; p != 0; p = p>>1) {
        if (p & 0x1)
            result *= x;
        x = x*x;
    }
    return result;
}
```

result	x	p

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“For” Loop Example

General Form

```
for (Init; Test; Update )
    Body
```

```
int result;
for (result = 1;
     p != 0;
     p = p>>1) {
    if (p & 0x1)
        result *= x;
    x = x*x;
}
```

Init

```
result = 1
```

Test

```
p != 0
```

Update

```
p = p >> 1
```

Body

```
{
    if (p & 0x1)
        result *= x;
    x = x*x;
}
```

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“For”→“While”

For Version

```
for (Init; Test; Update )
    Body
```

While Version

```
Init;
while (Test) {
    Body
    Update ;
}
```

Do-While Version

```
Init;
if (!Test)
    goto done;
do {
    Body
    Update ;
} while (Test)
done:
```

Goto Version

```
Init;
if (!Test)
    goto done;
loop:
    Body
    Update ;
    if (Test)
        goto loop;
done:
```

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“For” Loop Compilation

Goto Version

```
Init;
if (!Test)
    goto done;
loop:
    Body
    Update ;
    if (Test)
        goto loop;
done:
```

Init

```
result = 1
```

Test

```
p != 0
```

Update

```
p = p >> 1
```

```
result = 1;
if (p == 0)
    goto done;
loop:
    if (p & 0x1)
        result *= x;
    x = x*x;
    p = p >> 1;
    if (p != 0)
        goto loop;
done:
```

Body

```
{
    if (p & 0x1)
        result *= x;
    x = x*x;
}
```

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```

typedef enum
{ADD, MULT, MINUS, DIV, MOD, BAD}
op_type;

char unparse_symbol(op_type op)
{
    switch (op) {
    case ADD :
        return '+';
    case MULT:
        return '*';
    case MINUS:
        return '-';
    case DIV:
        return '/';
    case MOD:
        return '%';
    case BAD:
        return '?';
    }
}

```

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Switch Statements

Implementation Options

- Series of conditionals
 - Good if few cases
 - Slow if many
- Jump Table
 - Lookup branch target
 - Avoids conditionals
 - Possible when cases are small integer constants
- GCC
 - Picks one based on case structure
- Bug in example code
 - No default given

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Switch Statement Example

Branching Possibilities

```

typedef enum
{ADD, MULT, MINUS, DIV, MOD, BAD}
op_type;

char unparse_symbol(op_type op)
{
    switch (op) {
    * *
    }
}

```

Setup:

```

unparse_symbol:
    pushl %ebp      # Setup
    movl %esp,%ebp  # Setup
    movl 8(%ebp),%eax # eax = op
    cmpl $5,%eax   # Compare op : 5
    ja .L49         # If > goto done
    jmp *.L57(,%eax,4) # goto Table[op]

```

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Enumerated Values

ADD	0
MULT	1
MINUS	2
DIV	3
MOD	4
BAD	5

Jump Table Structure

Switch Form

```

switch(op) {
    case val_0:
        Block 0
    case val_1:
        Block 1
    ...
    case val_n-1:
        Block n-1
}

```

Jump Table

jtab:	Targ0
	Targ1
	Targ2
	•
	•
	•
	Targn-1

Jump Targets

Targ0:	Code Block 0
Targ1:	Code Block 1
Targ2:	Code Block 2
•	•
•	•
Targn-1:	Code Block n-1

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Approx. Translation

```

target = JTab[op];
goto *target;

```

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Assembly Setup Explanation

Symbolic Labels

- Labels of form .LXX translated into addresses by assembler

Table Structure

- Each target requires 4 bytes
- Base address at .L57

Jumping

- ```

jmp .L49

```
- Jump target is denoted by label .L49

```

jmp *.L57(,%eax,4)

```

  - Start of jump table denoted by label .L57
  - Register %eax holds op
  - Must scale by factor of 4 to get offset into table
  - Fetch target from effective Address .L57 + op\*4

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# Jump Table

## Table Contents

```
.section .rodata
.align 4
.L57:
.long .L51 #Op = 0
.long .L52 #Op = 1
.long .L53 #Op = 2
.long .L54 #Op = 3
.long .L55 #Op = 4
.long .L56 #Op = 5
```

## Enumerated Values

|       |   |
|-------|---|
| ADD   | 0 |
| MULT  | 1 |
| MINUS | 2 |
| DIV   | 3 |
| MOD   | 4 |
| BAD   | 5 |

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## Targets & Completion

```
.L51:
 movl $43,%eax # '+'
 jmp .L49
.L52:
 movl $42,%eax # '*'
 jmp .L49
.L53:
 movl $45,%eax # '-'
 jmp .L49
.L54:
 movl $47,%eax # '//'
 jmp .L49
.L55:
 movl $37,%eax # '%'
 jmp .L49
.L56:
 movl $63,%eax # '?'
 # Fall Through to .L49
```

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# Switch Statement Completion

```
.L49: # Done:
 movl %ebp,%esp # Finish
 popl %ebp # Finish
 ret # Finish
```

## Puzzle

- What value returned when op is invalid?

## Answer

- Register %eax set to op at beginning of procedure
- This becomes the returned value

## Advantage of Jump Table

- Can do k-way branch in  $O(1)$  operations

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# Object Code

## Setup

- Label .L49 becomes address 0x804875c
- Label .L57 becomes address 0x8048bc0

```
08048718 <unparse_symbol>:
08048718: 55 pushl %ebp
08048719: e5 movl %esp,%ebp
0804871b: 8b 45 08 movl 0x8(%ebp),%eax
0804871e: f8 05 cmpl $0x5,%eax
08048721: 77 39 ja 804875c <unparse_symbol+0x44>
08048723: ff 24 85 c0 8b jmp *0x8048bc0(,%eax,4)
```

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# Object Code (cont.)

## Jump Table

- Doesn't show up in disassembled code
- Can inspect using GDB

gdb code-examples

(gdb) x/6wx 0x8048bc0

- Examine 6 hexadecimal format "words" (4-bytes each)
- Use command "help x" to get format documentation

0x8048bc0 <\_fini+32>:

0x08048730  
0x08048737  
0x08048740  
0x08048747  
0x08048750  
0x08048757

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## Extracting Jump Table from Binary

### Jump Table Stored in Read Only Data Segment (.rodata)

- Various fixed values needed by your code

### Can examine with objdump

```
objdump code-examples -s --section=.rodata
```

- Show everything in indicated segment.

### Hard to read

- Jump table entries shown with reversed byte ordering

Contents of section .rodata:

```
8048bc0 30870408 37870408 40870408 47870408 0...7...@...G...
8048bd0 50870408 57870408 46616374 28256429 P...W...Fact(%d)
8048be0 203d2025 6c640a00 43686172 203d2025 = %ld..Char = %
...
...
```

- E.g., 30870408 really means 0x08048730

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## Disassembled Targets

```
8048730:b8 2b 00 00 00 movl $0x2b,%eax
8048735:eb 25 jmp 804875c <unparse_symbol+0x44>
8048737:b8 2a 00 00 00 movl $0x2a,%eax
804873c:eb 1e jmp 804875c <unparse_symbol+0x44>
804873e:89 f6 movl %esi,%esi
8048740:b8 2d 00 00 00 movl $0x2d,%eax
8048745:eb 15 jmp 804875c <unparse_symbol+0x44>
8048747:b8 2f 00 00 00 movl $0x2f,%eax
804874c:eb 0e jmp 804875c <unparse_symbol+0x44>
804874e:89 f6 movl %esi,%esi
8048750:b8 25 00 00 00 movl $0x25,%eax
8048755:eb 05 jmp 804875c <unparse_symbol+0x44>
8048757:b8 3f 00 00 00 movl $0x3f,%eax
```

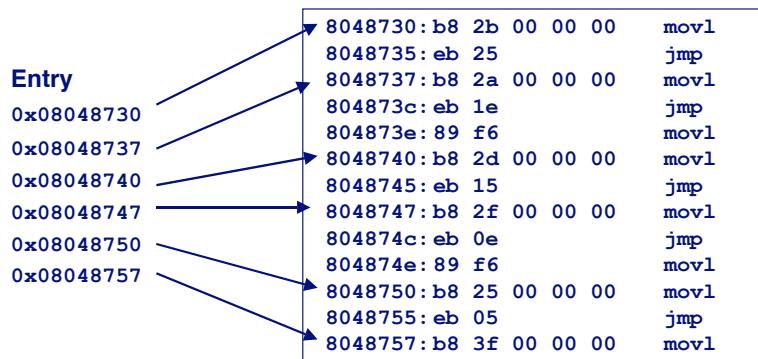
- movl %esi,%esi does nothing

- Inserted to align instructions for better cache performance

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## Matching Disassembled Targets



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## Sparse Switch Example

```
/* Return x/111 if x is multiple
 && <= 999. -1 otherwise */
int div111(int x)
{
 switch(x) {
 case 0: return 0;
 case 111: return 1;
 case 222: return 2;
 case 333: return 3;
 case 444: return 4;
 case 555: return 5;
 case 666: return 6;
 case 777: return 7;
 case 888: return 8;
 case 999: return 9;
 default: return -1;
 }
}
```

- Not practical to use jump table
  - Would require 1000 entries
- Obvious translation into if-then-else would have max. of 9 tests

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## Sparse Switch Code

```

movl 8(%ebp),%eax # get x
cmpl $444,%eax # x:444
je L8
jg L16
cmpl $111,%eax # x:111
je L5
jg L17
testl %eax,%eax # x:0
je L4
jmp L14

...

```

- Compares x to possible case values
- Jumps different places depending on outcomes

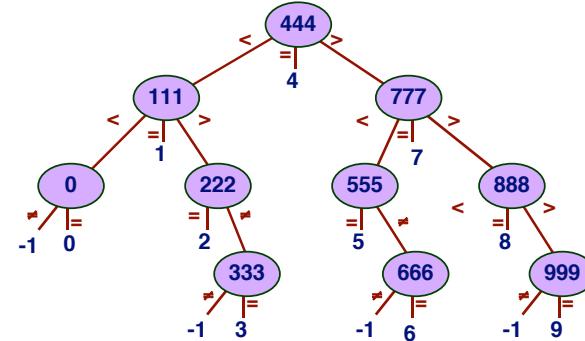
```

 . . .
L5:
 movl $1,%eax
 jmp L19
L6:
 movl $2,%eax
 jmp L19
L7:
 movl $3,%eax
 jmp L19
L8:
 movl $4,%eax
 jmp L19
 . . .

```

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## Sparse Switch Code Structure



- Organizes cases as binary tree
- Logarithmic performance

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## Summarizing

### C Control

- if-then-else
- do-while
- while
- switch

### Assembler Control

- jump
- Conditional jump

### Compiler

- Must generate assembly code to implement more complex control

### Standard Techniques

- All loops converted to do-while form
- Large switch statements use jump tables

### Conditions in CISC

- CISC machines generally have condition code registers

### Conditions in RISC

- Use general registers to store condition information
  - Special comparison instructions
  - E.g., on Alpha:
- ```
cmple $16,1,$1
    • Sets register $1 to 1 when Register $16 <= 1
```

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