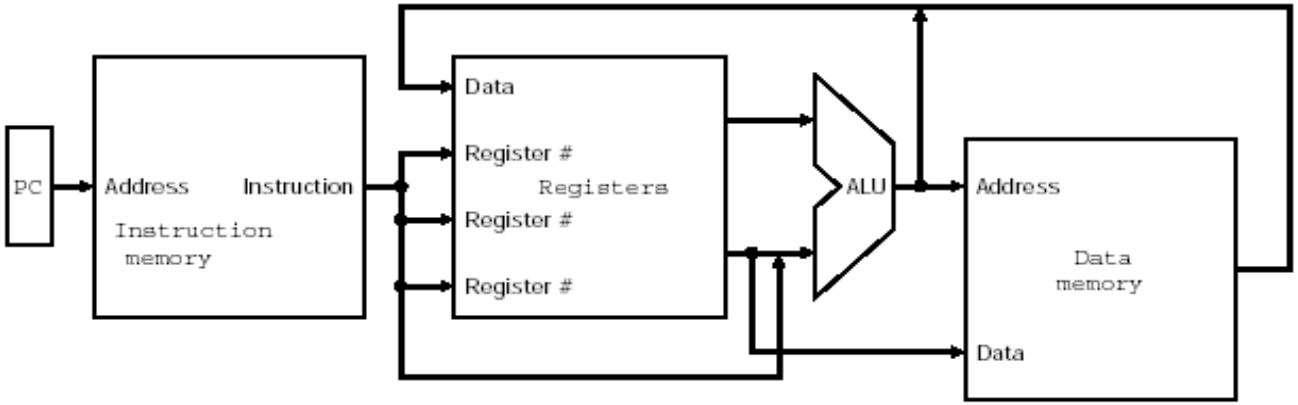
CPU



CPU = datapath + control

ALU: arithmetic/logic unit

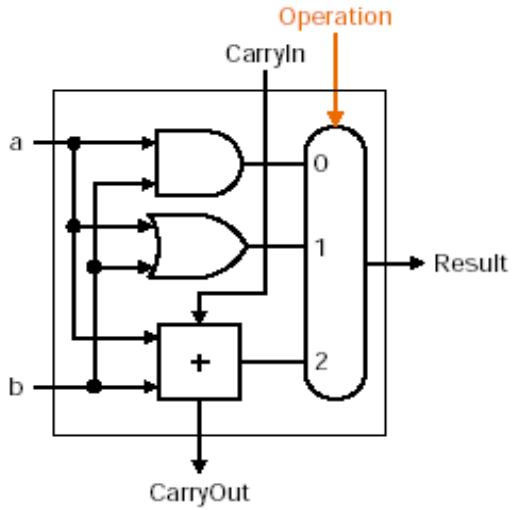
performs operations to execute arithmetic and logical instructions

ALU: 1-bit

We now have the ingredients for a simple 1-bit arithmetic-logic unit (ALU)

Operations:	**ADD**	**a + b + c_{in}**
	AND	**a AND b**
	OR	**a OR b**
Inputs:	**data:**	**a, b, c_{in}**
	control:	**con_1, con_2**
Outputs:		**result, c_{out}**

Components:	**AND gate**
	OR gate
	Full adder
	4-1 MUX



1-bit ALU (Fig. 4.14)

(Operation is 2-bit control con_1, con_2)

Can construct k-bit ALU by combining k 1-bit ALUs

What other operations could we have?

ALU: 1-bit

How about subtraction?

We can use the adder to add the negated form of the operand

`a - b = a + (-b)`

Add an inverter to the circuit to negate b

This gives 1's complement

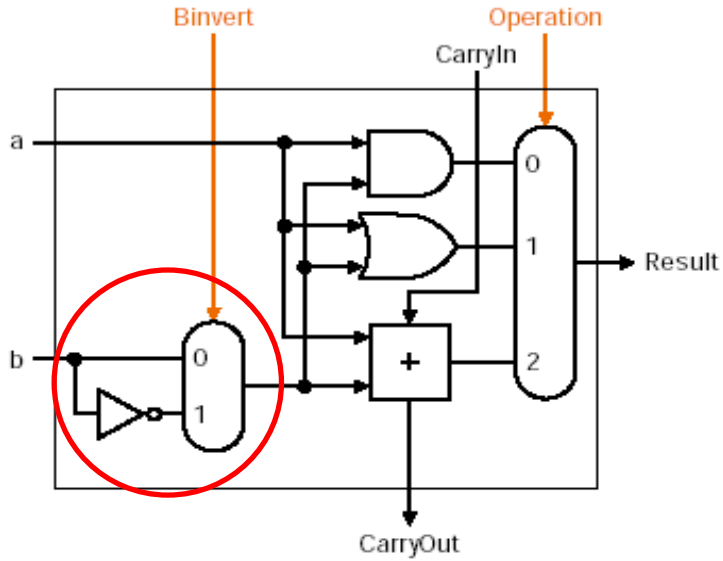
How do we get 2C value?

Use c_{in} = 1 for least significant bit

`a + ~b + 1 = a + (~b + 1) = a + (-b) = a - b`

Another MUX with control input Binvert

can select b or -b



1-bit ALU with subtraction

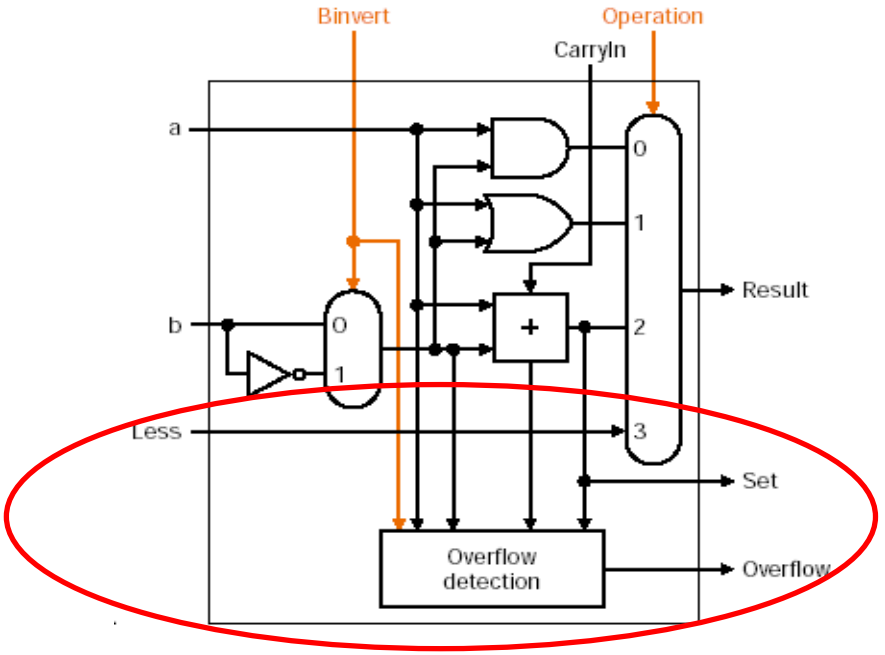
(Fig. 4.16)

ALU: 1-bit

This ALU can perform most of the data operations in the MIPS instruction set
Another operation, useful for branching: set on less than (`slt`)
Set the lsb to 1 if rs < rt, and 0 otherwise
If (a-b) is negative, then a < b: `(a - b) < 0`
`(a - b) + b < 0 + b`
`a < b`
Result is same as sign bit from subtraction: Connect sign bit from adder to lsb of output
Unfortunately, we can only do 1 ALU operation at a time (add or `slt`)
Need a new 1-bit ALU for the msb
- **with an extra output from adder**
- **Extra output: `Set`**

Additional MUX data input: `Less`
- **0 for all except lsb**
- **`Set` value for lsb**

Also add overflow detection

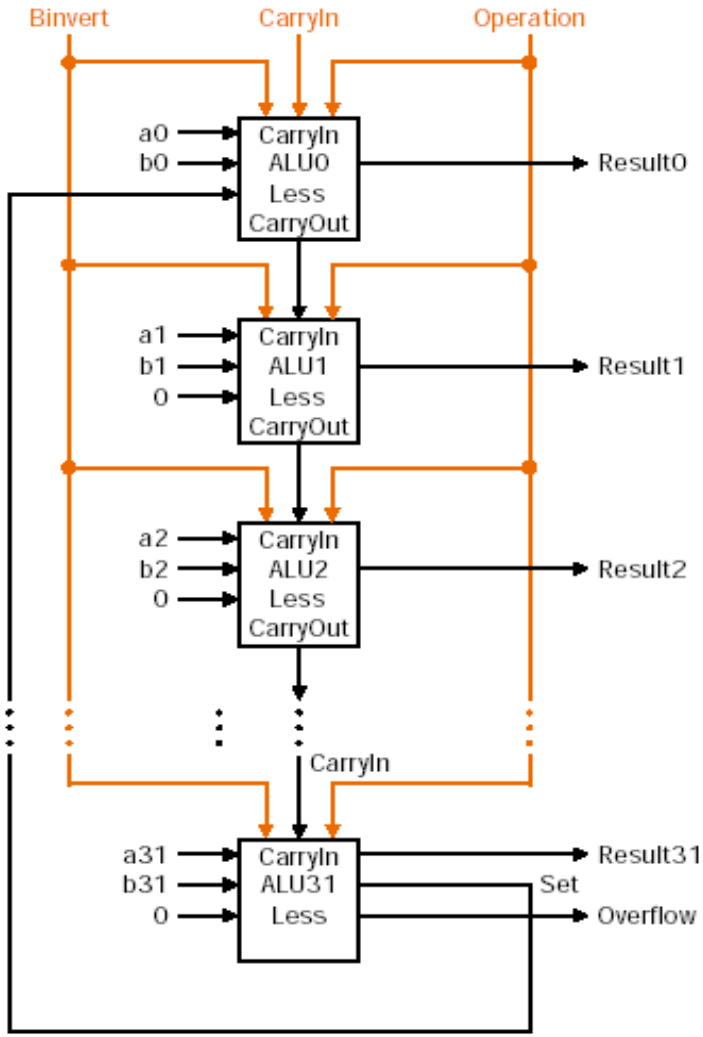


1-bit ALU with set on less than
(Fig. 4.17b)

ALU: k-bit

To operate on k-bit values,
we can connect k 1-bit ALU's

32-bit ALU is constructed using 32
1-bit ALU's
Input bits are connected in pairs
Control bits (Binvert, Operation)
are connected to each ALU
c_{out} from each ALU is connected to
c_{in} of next most significant bit ALU
(ripple carry)
c_{in} for lsb is 1 for subtract operation
`Set` from ALU31 (msb) is connected
to `Less` input of ALU0 (lsb)
(0 input for all other ALUs)
`Overflow` from ALU31 is additional output



32-bit ALU (Fig. 4.18)

ALU: k-bit

What about conditional branch?

- **Branch if 2 values are either equal or not equal**
- **Easiest way to test if a == b: subtract, test result:**
  - **OR all result bits together and complement**

One more refinement: combine Binvert and CarryIn control values into Bnegate:

- **subtract: both are 1**
- **add or logical ops: both are 0**

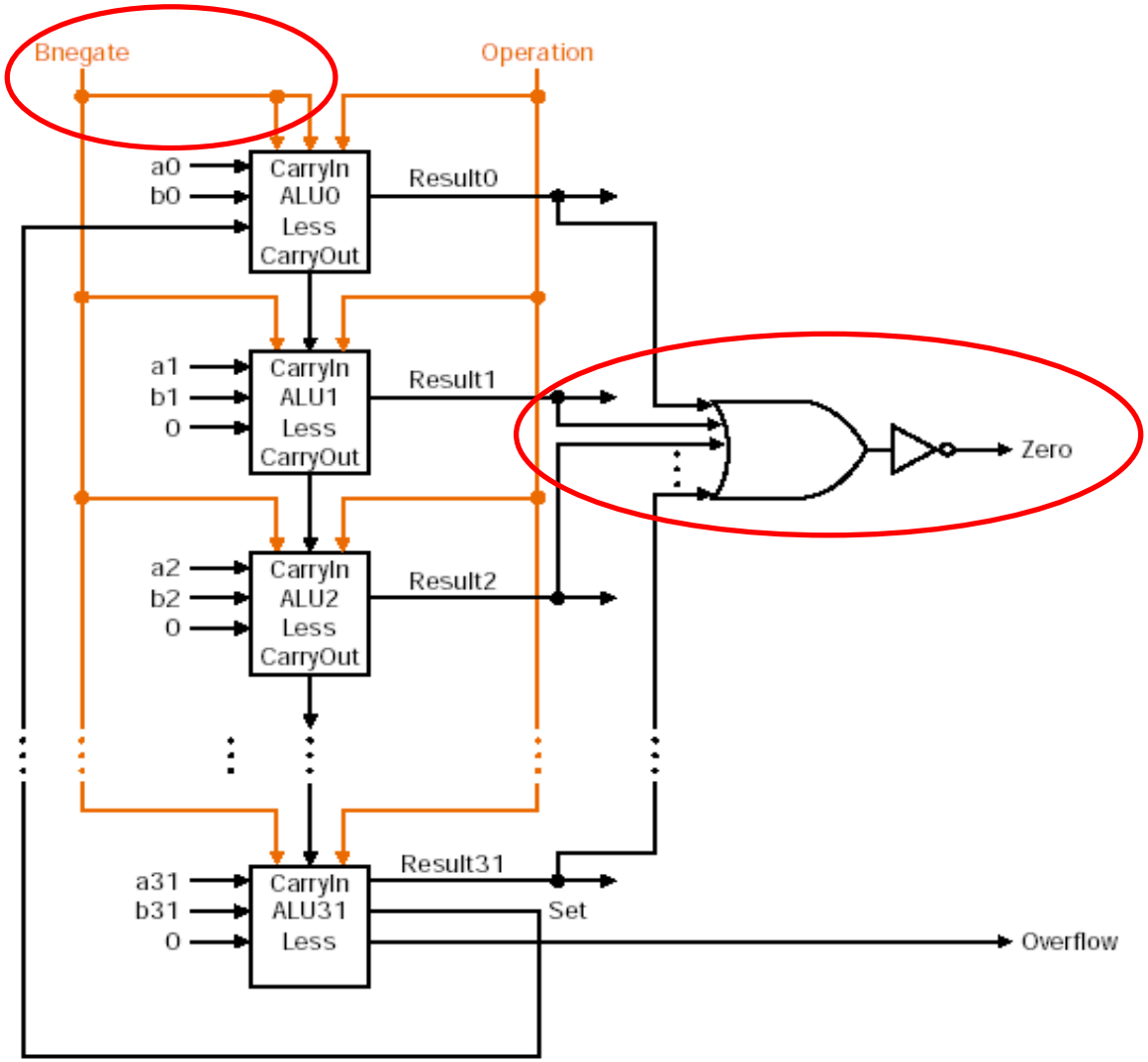
Bnegate (1 bit) and Operation (2 bits) are 3-bit control for MUX:

control	function
000	AND
001	OR
010	ADD
110	SUB
111	SLT

ALU: k-bit

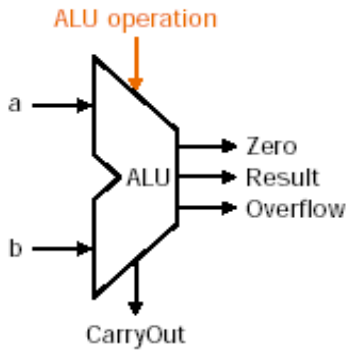
Additions:
Bnegate control input
to subtract
Zero output:
inverted OR of
all outputs
Input: a, b
Control:
Bnegate
Operation
Output:
Result
Zero
Overflow

32-bit ALU with zero detection
(Fig. 4.19)



ALU: k-bit

Universal symbol to represent ALU
Can also be used for adder alone, so labeled accordingly



(Fig. 4.21)