

Introduction

- Reading
 - for today Chapter 8
 - for Thursday Chapter 10
- Please email me the class you would like to lead discussion

Vector Timestamps in LRC

- Each Acquire/Release advances local clock
 - I.e. starts a new interval
- At Acquire:
 - compute new time at pair wise max of previous local and release.
 - Send old acquire timestamp to release node
 - Releasing node sends updates to node acquire node
- Maintain diffs for each node for each interval
- At Barrier, all participants in barrier see all changes
 - if barrier is global (all threads participate), can discard previous diffs

Interconnection Networks

● Performance Metrics

- latency: time for the first bits to arrive
- bandwidth: sustained volume of communication
- connectivity: how far is the furthest neighbor
- hardware cost: can we afford to build this thing?
- reliability: can links fail and the system still run?
- functionality: “for free” features
 - combining messages
 - fault tolerance

● Options

- topology
- operation mode (synchronous or asynchronous)
- switching strategy (circuit, store-and-forward, cut-through)
- control strategy (central or distributed)

Networks

● Static

- Processors only communicate with neighbors
 - non-neighbor communication requires help of intermediate processors
 - this is a simple structure to build
 - performance suffers due to using processors to forward
- Usual topologies are possible
 - hypercube, tree, mesh

● Dynamic

- “switches” take care of forwarding between nodes
 - requires extra hardware to forward messages
 - hardware has to be able to route data
- Topologies
 - bus, multi-stage, crossbar networks
 - any static topology

Switching Options

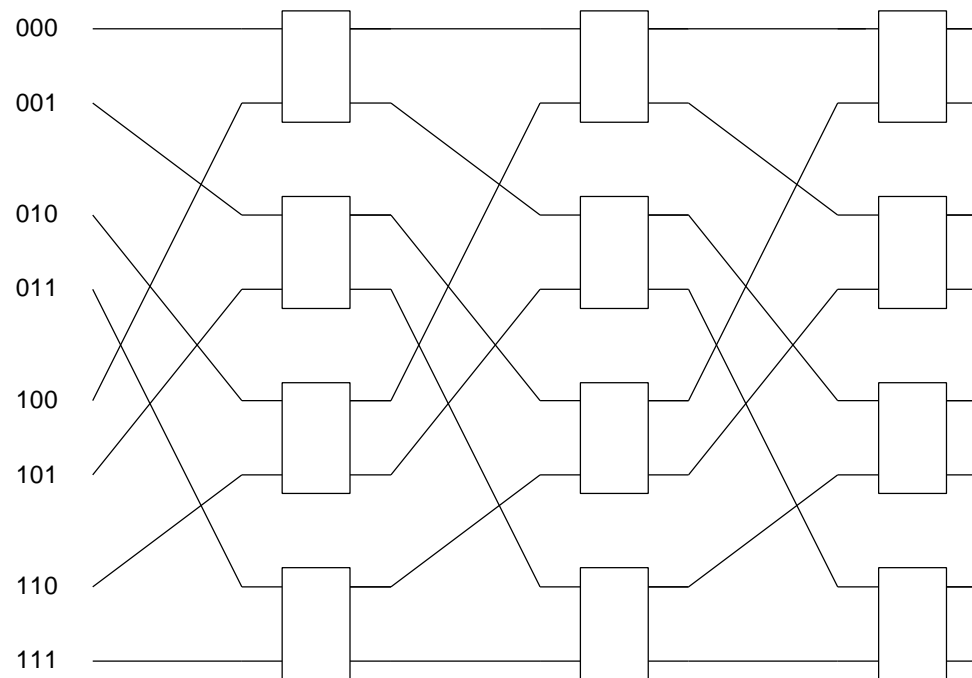
- **Circuit switching**
 - “connection” is established between end points
 - contention free connection is then used for data
 - latency is length + number of hops
- **Store-and-forward**
 - send hop by hop with the message being buffered
 - latency is message length * number of hops
- **Cut-through (hybrid)**
 - wormhole routing
 - first part of a message establishes a circuit
 - rest of message snakes its way long the path
 - possible to deadlock if not careful

Multistage Switching Networks

- several stages between source and destination
- provide different options in
 - cost
 - wires
 - blocking probability
- can be combined into many hybrid options

Omega Networks

- $\log_n(M)$ stages for M nodes and n by n cross bars
- message block if the required output link is not available



Benes networks

- additional stages to reduce blocking
- can route any disjoint src/dest pairs wo blocking
 - but must know requests in advance

