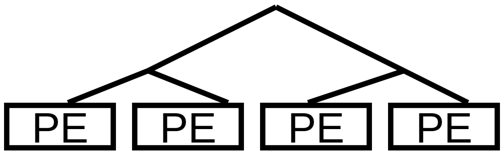
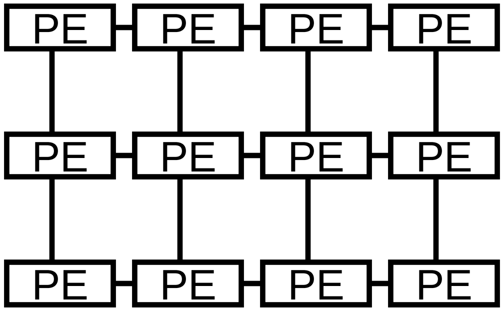
Announcements

Topologies (cont)

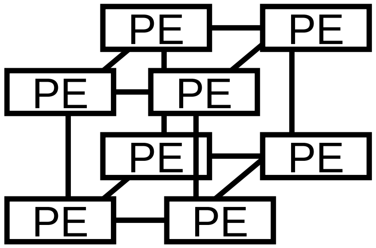
– tree - needs to increase bandwidth near the top



–mesh - two or three dimensions



–hypercube - needs a power of number of nodes



Memory Systems

- Key Performance Issues
  - latency: time for first byte
  - throughput: average bytes/second
- Design Issues
  - Where is the memory
    - divided among each node
    - centrally located (on communication network)
  - Access by processors
    - can all processors get to all memory?
    - is the access time uniform?

Coordination

- Synchronization
  - protection of a single object (locks)
  - coordination of processors (barriers)
- Size of a unit of work by a processor
  - need to manage two issues
    - load balance - processors have equal work
    - coordination overhead - communication and sync.
  - often called “grain” size - large grain vs. fine grain

Sources of Parallelism

- Statements
  - called “control parallel”
  - can perform a series of steps in parallel
- Loops
  - called “data parallel”
  - most common source of parallelism
  - each processor gets one (or more) iterations to perform

Example of Parallelism

- Easy (embarrassingly parallel)
  - multiple independent jobs (i.e..., different simulations)
- Scientific
  - Largest users of parallel computing
  - dense linear algebra (divide up matrix)
  - physical system simulations (divide physical space)
- Databases
  - biggest commerical success of parallel computing (divide tuples)
    - exploits semantics of relational calculus
- AI
  - search problems (divide search space)
  - pattern recognition and image processing (divide image)

Metrics in Application Performance

- Speedup (often call strong scaling)
  - ratio of time on n nodes to time on a single node
  - hold problem size fixed
  - should really compare to best serial time
  - goal is linear speedup
  - super-linear speedup is possible due to:
    - adding more memory
    - search problems
- Weak Scaling (also called Iso-Speedup)
  - scale data size up with number of nodes
  - goal is a flat horizontal curve
- Amdahl's Law
  - max speedup is 1/(serial fraction of time)
- Computation to Communication Ratio
  - goal is to maximize this ratio

How to Write Parallel Programs

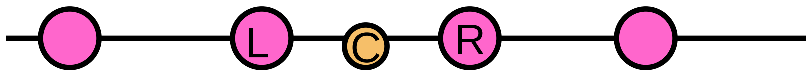
- Use old serial code
  - compiler converts it to parallel
  - called the dusty deck problem
- Serial Language plus Communication Library
  - no compiler changes required!
  - PVM and MPI use this approach
- New language for parallel computing
  - requires all code to be re-written
  - hard to create a language that provides performance on different platforms
- Hybrid Approach
  - HPF - add data distribution commands to code
  - add parallel loops and synchronization operations

Application Example - Weather

- Typical of many scientific codes
  - computes results for three dimensional space
  - compute results at multiple time steps
  - uses equations to describe physics/chemistry of the problem
  - grids are used to discretize continuous space
    - granularity of grids is important to speed/accuracy
- Simplifications (for example, not in real code)
  - earth is flat (no mountains)
  - earth is round (poles are really flat, earth buldges at equator)
  - second order properties

Grid Points

- Divide Continuous space into discrete parts
  - for this code, grid size is fixed and uniform
    - possible to change grid size or use multiple grids
  - use three grids
    - two for latitude and longitude
    - one for elevation
    - Total of M * N * L points
- Design Choice: where is the grid point?
  - left, right, or center of the grid



  - in multiple dimensions this multiples:
    - for 3 dimensions have 27 possible points

Variables

- One dimensional
  – m - geo-potential (gravitational effects)
- Two dimensional
  – pi - “shifted” surface pressure
  – sigmadot - vertical component of the wind velocity
- Three dimensional (primary variables)
  – <u,v> - wind velocity/direction vector
  – T - temperature
  – q - specific humidity
  – p - pressure
- Not included
  – clouds
  – precipitation
  – can be derived from others

Serial Computation

- Convert equations to discrete form
- Update from time t to t + delta t

```
foreach longitude, latitude, altitude
        ustar[i,j,k] = n * pi[i,j] * u[i,j,k]
        vstar[i,j,k] = m[j] * pi[i,j] * v[i,j,k]
        sdot[i,j,k] = pi[i,j] * sigmadot[i,j]
end
foreach longitude, latitude, altitude
        D = 4 * ((ustar[i,j,k] + ustar[i-1,j,k]) * (q[i,j,k] + q[i-1,j,k]) +
                        terms in {i,j,k}{+,-}{1,2}
        piq[i,j,k] = piq[i,j,k] + D * delat
        similar terms for piu, piv, piT, and pi
end
foreach longitude, latitude, altitude
        q[i,j,k] = piq[i,j,k]/pi[i,j,k]
        u[i,j,k] = piu[i,j,k]/pi[i,j,k]
        v[i,j,k] = piv[i,j,k]/pi[i,j,k]
        T[i,j,k] = piT[i,j,k]/pi[i,j,k]
end
```

Shared Memory Version

- in each loop nest, iterations are independent
- use a parallel for-loop for each loop nest
- synchronize (barrier) after each loop nest
  - this is overly conservative, but works
  - could use a single sync variable per item, but would incur excessive overhead
- potential parallelism is M * N * L
- private variables: D, i, j, k
- Advantages of shared memory
  - easier to get something working (ignoring performance)
- Hard to debug
  - other processors can modify shared data

Distributed Memory Weather

- decompose data to specific processors
  - assign a cube to each processor
    - maximize volume to surface ratio
    - minimizes communication/computation ratio
  - called a <block,block,block> distribution
- need to communicate {i,j,k}{+,-}{1,2} terms at boundaries
  - use send/receive to move the data
  - no need for barriers, send/receive operations provide sync
    - sends earlier in computation too hide communication time
- Advantages
  - easier to debug?
  - consider data locality explicitly with data decomposition
- Problems
  - harder to get the code running

Seismic Code

- Given echo data, compute under sea map
- Computation model
  - designed for a collection of workstations
  - uses variation of RPC model
  - workers are given an independent trace to compute
    - requires little communication
    - supports load balancing (1,000 traces is typical)
- Performance
  - max mfops = $O((F * nz * B^*)^{1/2})$
  - F - single processor MFLOPS
  - nz - linear dimension of input array
  - B^* - effective communication bandwidth
    - $B^* = B/(1 + BL/w) \approx B/7$ for Ethernet (10msec lat., w=1400)
  - real limit to performance was latency **not** bandwidth