

Announcements

- Reading
 - Today:3.1-3.3
 - Tuesday: 3.5-3.6
 - section 3.4 was covered before (during session)

Data Link Layer

- Goal: transmit error free frames over the physical link
- Sample Issues:
 - how big is a frame?
 - can I detect an error in sending the frame?
 - what demarks the end of the frame?
 - how to control access to a shared channel?
- Examples:
 - Ethernet framing

Frames

- Slice Raw bit stream up into frames
 - need to have manageable unit of transmission
- Frame Boundary
 - How do we know when a frame ends?
 - Character count
 - header indicates number of bytes
 - problem: what if the header is corrupt, can't tell end of frame
 - Special character
 - ASCII: DLE STX ... DLE STE
 - need to use character stuffing to send DLE characters
 - send two DLE to indicate a DLE
 - Special bit pattern - no longer tied to ASCII
 - 01111110 - indicates end of frame
 - need to use bit stuffing to send 01111110 as data
 - insert 0 after 5 1's
 - use link level invalid bit patterns
 - some bits may not be valid

Other Link Functions

- Error Control

- may want to do sequence numbers and re-transmission
- this introduces overhead, but useful if probability of failure is high

- Flow Control

- provide rate matching between sender and receiver
- sender has rules about when it can send: credits, etc.

Error Correcting Codes

- Idea: add redundant information to permit recovery
 - this is the dual of data compression (remove redundancy)
- Hamming distance (n)
 - number of bit positions that differ in two words
 - key idea: need n single bit errors to go from one word to the other
 - to detect d errors, need a hamming distance of $d+1$ from **any other valid word.**
 - to recover d errors, need a hamming distance of $2d + 1$
 - any error of d bits is still closer to correct word
- Parity bit
 - ensure that every packet has an odd (or even) # of 1's
 - permits detection of one 1 bit error

Error Codes (cont.)

- Error Recovery

- Given m bits of data and r bits of error code
- Want to correct any one bit error
- There are n words one bit from each valid message
 - so need $n+1$ words for each valid message
 - thus $(n + 1) 2^m \leq 2^n$
 - but $n = m + r$ so $(m + r + 1) \leq 2^r$

- Hamming Code

- recovers from any one bit error
- number bits from left (starting at 1)
 - power of two bits are parity
 - rest contain data
- bit is checked by all parity bits in its sum of power expansion
 - bit 11 is used to compute parity bits 1, 2, and 8

Hamming Code Example

Char	ASCII	Hamming
H	1001000	00110010000
a	1100001	10111001001
m	1101101	11101010101
l	1101001	01101011001

- **Burst Errors**

- can send hamming codes by column rather than row
- if use k rows, then can detect any burst error up to k bits
 - uses kr bits to check a block km bits long

Error Detection

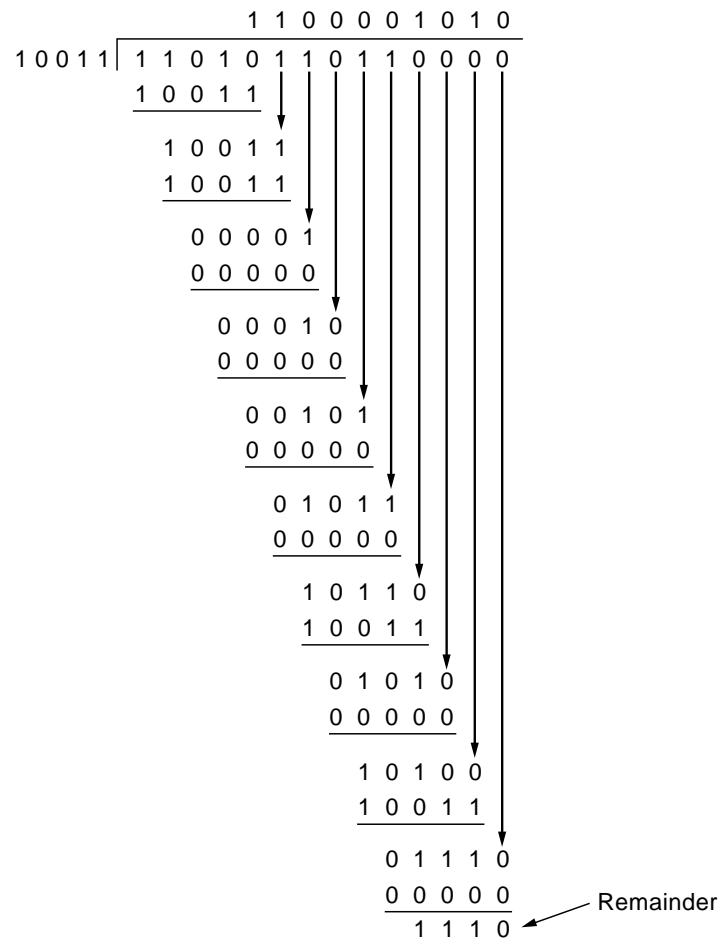
- Less bits are required
 - if errors are infrequent, then then this works better
 - assumes that re-transmission is possible
- Cyclic Redundancy Codes (CRC)
 - Use a generator function $G(x)$ of degree r
 - let M' be the message with r 0's on the end of it
 - divide M' into $G(x)$ and compute remainder
 - use this as the r bit CRC code
 - a code with r bits will detect all burst errors less than r bits
 - several G 's are standardized
 - $\text{CRC-12} = x^{12} + x^{11} + x^3 + x^2 + x + 1$
 - $\text{CRC-16} = x^{16} + x^{15} + x^2 + 1$
 - $\text{CRC-CCITT} = x^{16} + x^{12} + x^5 + 1$
 - 16 bit CRC will catch
 - all single and double bit errors
 - all errors with an odd number of bits
 - all burst errors of length less than 16

CRC Example

Frame : 1101011011

Generator: 10011

Message after appending 4 zero bits: 11010110000



Transmitted frame: 1101011011110

Data Link Protocols

- Stop And Wait
 - send a frame
 - wait for ACK
 - need sequence number to tell re-transmission from next packet
 - lost ACK vs. lost frame
- Sliding Window
 - sequence numbers
 - can send up to window size number frames
 - Retransmission
 - Go Back N
 - Selective repeat