## Announcements

#### • reading

- for Thursday 5.5

#### • midterms

- Mt #1 Tuesday March 6
- Mt #2 Tuesday April 15
- Final project design due April 11

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## **Network Layer**

#### • Responsibility

- end-to-end delivery of packets to the network
- selecting routes for the packets to take
  - implies knowledge of the network topology
- managing utilization of the links
  - provide flow control (across multiple links)
  - spread load among different routes
- Interface Design
  - should be independent of subnet technology
  - hide number, type, and topology of network from upper layers
  - export a common number plan for entire network

## Connection vs. Connectionless

- Two possible designs for network layer
  - connection oriented service (ATM)
    - based on experience of telcos
  - connectionless service (IP)
    - based on packet switching (ARPANET)
- Connectionless
  - transport datagrams from source to destination
    - end-point addresses in every datagram
  - less complex network layer, more complex transport
- Connection oriented
  - also called virtual circuits
  - establish an end-to-end connection with network state
    - can use VCI (global or next hop) in each packet

## Datagram vs. VC Addresses

#### • Datagrams

- must include full address in each packet
- addresses must be unqiue for entire network
  - don't re-use too often
  - addresses per src/dest pair
- Virtual Circuit
  - globally unique
    - requires allocation scheme to ensure its unique
    - consumes many bits per packet
  - per link
    - requires translation at each switch
    - uses fewer bits (important for small packets like ATM)

# Link Failue in Virtual Circuits • Re-establish virtual circuit - router near failure can patch up link - original host/router creates new virtual circuit Virtual circuit is dropped - transport layer can handle recovery

# Virtual Circuit vs. Datagram

Issue	Datagram	Virtual Circuit
Circuit setup	not needed	necessary
Addresses	full source/dest per packet	next hop vc sufficient
state	no state in network	per connection data at each router
routing	each packet individually	once at VC setup
router/link failure	a few packets may be lost	all VCs through router are terminated
congestion control	difficult	many pre-allocation and policing policies permitted

# Routing: Goals

#### • Correctness

- packets get where they are supposed
- Simplicity
  - easy to implement correctly
  - possible to make routing choices fast (or updates easy)
- Robustness
  - failures in the network still permit communication
- Stability
  - small changes in link availability results in a small change in the routing information

#### • Fairness

- each host, VC, or datagram has the same chance
- Optimality
  - best possible route
  - best utilization of bandwidth

# Do Routes Change During Network Operation?

- nonadaptive routing (static routing)
  - information loaded a boot time
  - never changes during network operation
- adaptive routing
  - changes in network operation alter routes
  - issue: where to get this data to make choices
    - locally from neighbors
    - globally from all routers (or a NIC network information center)
  - issue: when to change routes
    - only on topology changes (links or routers change)
    - in response to changes in load
  - issue: metric to optimize
    - distance, number of hops, estimated latency

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