

Announcements

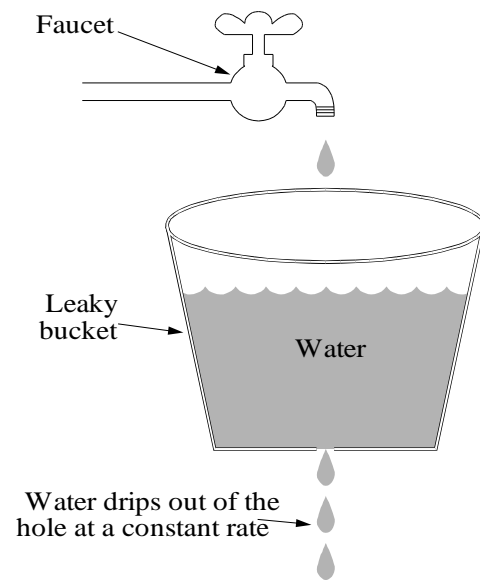
- This Lecture was presented by Dr. Shankar
- These are my lecture notes for similar topics, but not the ones used during this lecture.

Traffic Shaping

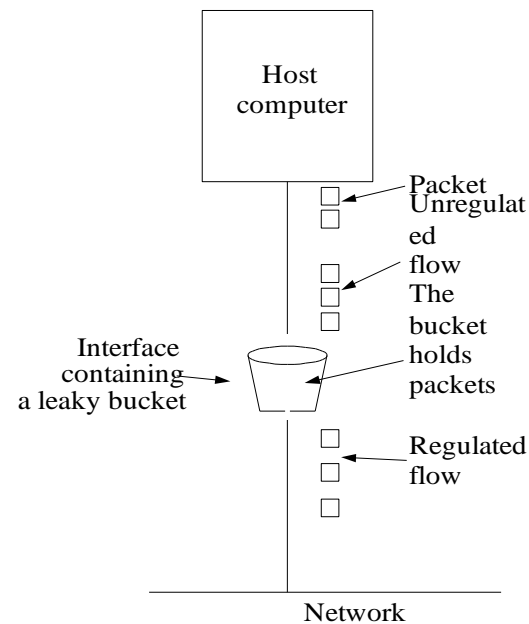
- Traffic tends to be bursty
 - great variation between min and max bandwidth used
 - this uncertainty leads to inefficient use of the network
- Flow Specification
 - user proposes a specific probability distribution
 - maximum packet size
 - transmission rate (min, max, or mean)
 - maximum delay
 - maximum delay variation (jitter)
 - quality guarantee (how strong is this agreement)
 - network can
 - agree to request
 - refuse it
 - counter offer

Leaky Bucket

- buffer accepts traffic at link rate
 - buffer has a bounded size (limits burst size that is accepted)
- output is limited to a lower rate
 - traffic is constrained to this rate



(a)

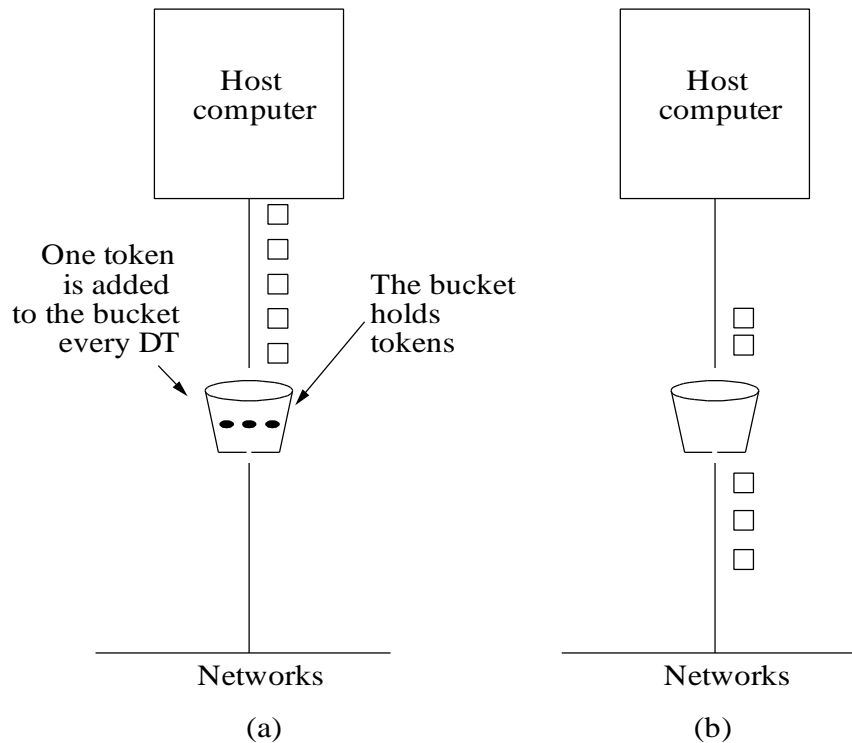


(b)

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Token Bucket

- Bucket hold tokens (generated one every T seconds)
- Can save up to a fixed limit of n tokens
- When traffic arrives, it must a have token to be sent



- Max burst rate

C - capacity of bucket

S - burst length in seconds

M - max output rate

p - token credit rate

$$C + pS = MS$$

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Congestion Control with Virtual Circuits

- Admission control
 - once traffic reaches a threshold, don't admit more VCs
 - doesn't correct current problem, but prevents additional congestion
- Alter routes
 - admit new connections
 - route them around "trouble" areas
- Negotiate traffic
 - establish parameters for volume and shape of traffic

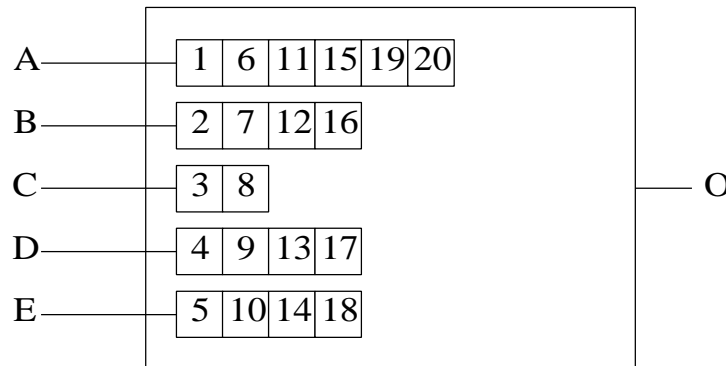
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Fair Queuing

- Local (per router) congestion control

- each output link has n queues, one for each sender
 - need to limit max queue size or buffers will be exhausted
- use round-robin to select next packet to queue
 - can use per-packet or per-byte



Packet	Finishing time
C	8
B	16
D	17
E	18
A	20

(b)

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- Weighted Fair Queuing

- can give different links different priorities
- give higher priority length multiple slots per round