

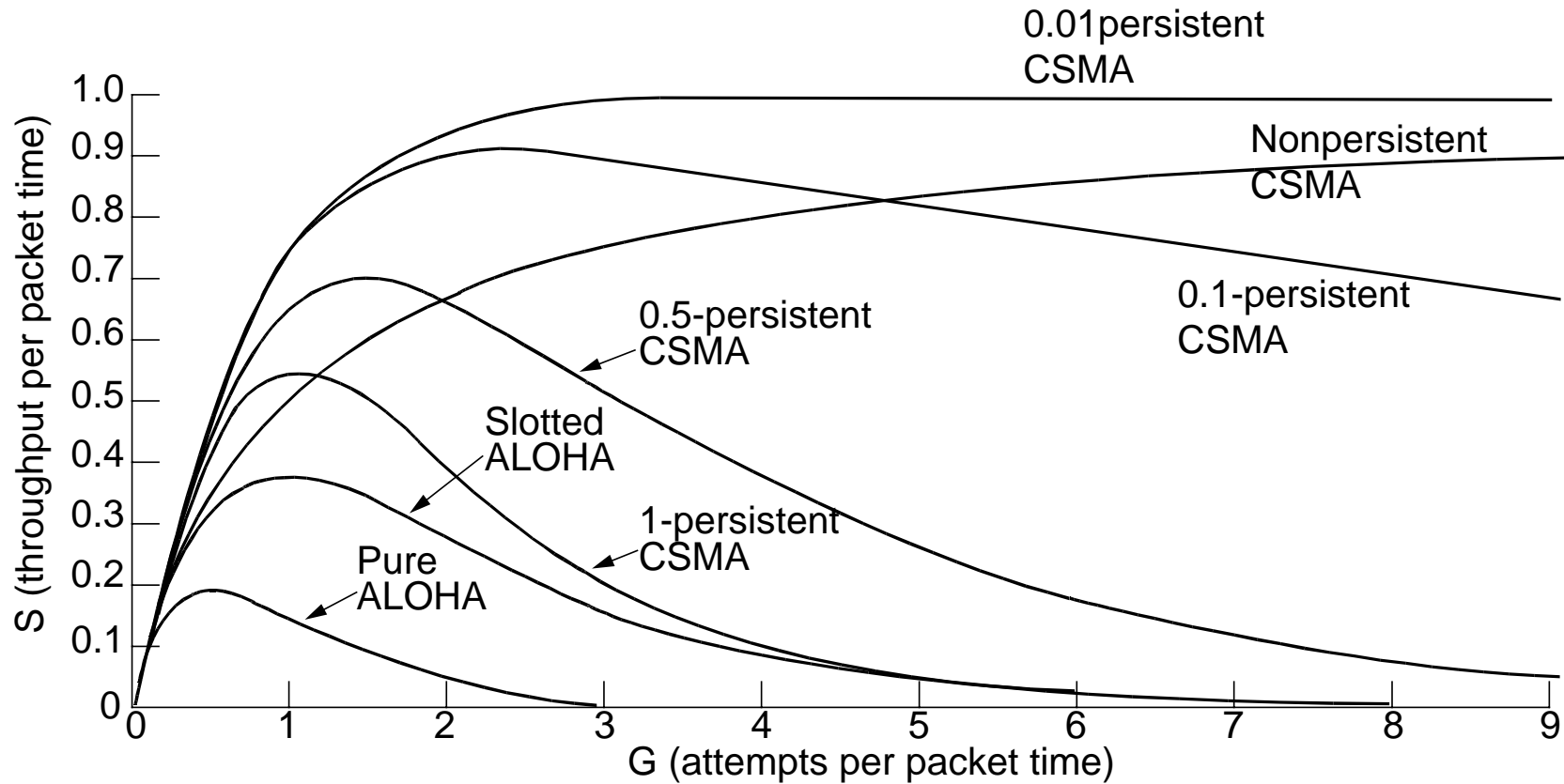
Announcements

- Reading
 - Chapter 4 (4.3.1, 4.4, 4.5)
- Homework #1
 - Due today
- Midterm #2
 - 11/8/01 in class
- No office hours next week

Carrier Sense Multiple Access

- look before you leap!
 - don't send if someone else is sending
- collisions are still possible
 - propagation delay induces uncertainty into sensing
 - possible two hosts both start sending at the same time
- persistence: when to send after detecting channel in use
 - 1-persistent
 - as soon as the channel is free, starting sending
 - nonpersistent CSMA
 - if channel is sensed busy, wait a random time and try again
 - p-persistent CSMA
 - if slot is idle send with probability p , else wait for next idle slot

Impact of Carrier Sense



From: *Computer Networks*, 3rd Ed. by Andrew S. Tanenbaum, (c)1996 Prentice Hall.

Collision Detection

- If a sender senses a collision
 - stop sending at once
 - apply random backoff
- “contention” period
 - after contention period, there will be no collision
 - send for for 2τ (max propagation delay)
 - need 2τ since might be a collision at far end at $\tau - \epsilon$

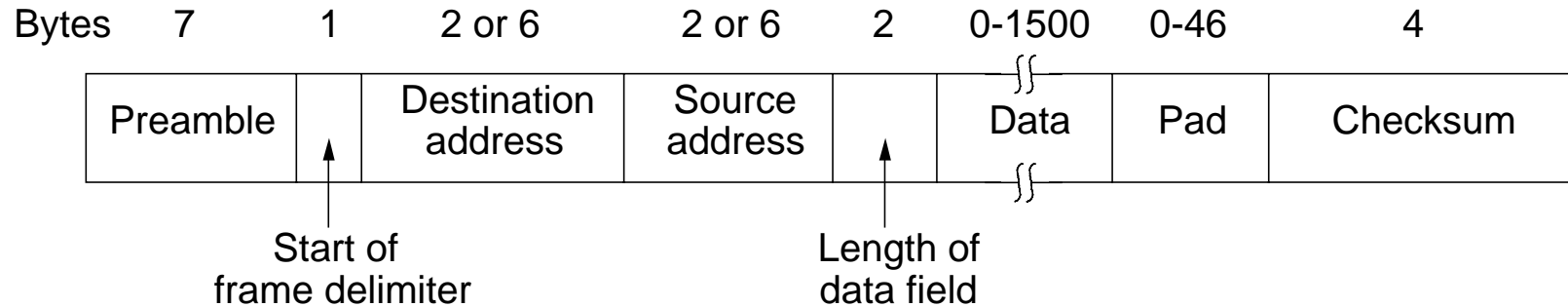
Ethernet Cable Options

- **10base5: Thicknet - first Ethernet**
 - Thick cable, doesn't bend well
 - vampire taps used to "tap" the network
 - max run is 500 meters
- **10Base2: Thin coax (cheaper net),**
 - uses "T" connectors
 - max run is 200 meters
- **10baseT: twisted pair**
 - uses a central hub
 - easier to find faults and problems
 - max run is 100 meters to hub

Manchester Encoding

- **Problem: How to send zero/ones?**
 - need to know timing information
 - when does on bit end?
- **Answer: Force many transitions**
 - every bit is half low and half high
 - 1 is high then low
 - 0 is low then high
 - but this doubles bandwidth
- **Differential Manchester Encoding**
 - better noise immunity
 - 0 is a transition at the start, 1 none
 - both transition during the middle

Ethernet Frame Format



From: *Computer Networks*, 3rd Ed. by Andrew S. Tanenbaum, (c)1996 Prentice Hall.

- Preamble used to sync clock
- Addresses
 - 48 bits
 - if it starts with a 0 it is globally unique (assigned by IEEE)
 - if it starts with a 1 it is locally unique
- Length
 - 0 to 1500 bytes
 - **min** length is 46 bytes
 - ensures frame reaches end of cable before end of frame is sent
- Checksum
 - 32 bit CRC to detect garbled data at link level

Collision Management

- Binary Exponential Backoff

- after collision, divide into slot times
- after first collision, wait either 0 or 1 slot times
- after second collision, wait either 0, 1, 2, or 3 slot times
- limited to 1023 slots
- after 16 collisions, link layer gives up

- Performance

- each station wants to transmit with probability p , then
 - $A = k [p^1(1-p)^{k-1}]$
 - $A \rightarrow 1/e$ as $k \rightarrow \text{infinity}$
- probability a contention interval has j slots is $A(1-A)^{j-1}$
- mean number of slots per contention is:

$$\sum_{j=0}^{\infty} jA(1-A)^{j-1} = \frac{1}{A} \quad \text{mean contention interval is then } 2\tau/A$$

Ethernet Performance (cont.)

- Ethernet Channel efficiency is then:

$$\frac{P}{P + 2\tau / A} = \frac{1}{1 + 2BLE / cF}$$

B = bandwidth

L = cable length

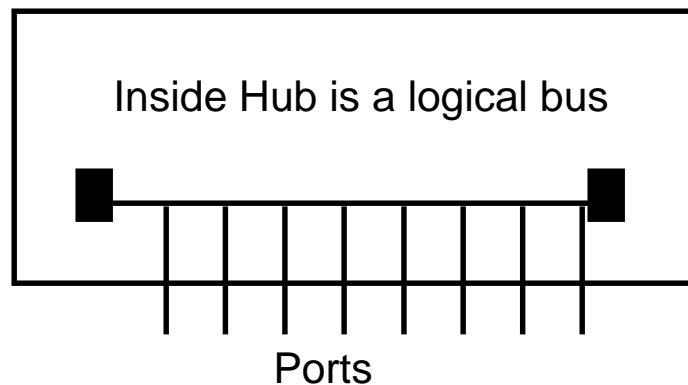
c = speed of light

F = frame length

- Traffic models
 - traditional analysis assume Poisson arrival
 - recent studies have demonstrated self similar properties
 - traffic variance does not decrease with wider samples

Variations on Ethernet

- Traditional Ethernet is a bus
 - limited to one host at a time
 - Requires long wires
- Hub based Ethernet
 - Cables form a star
 - Basic bus is still used (one large collision domain)



Bridges

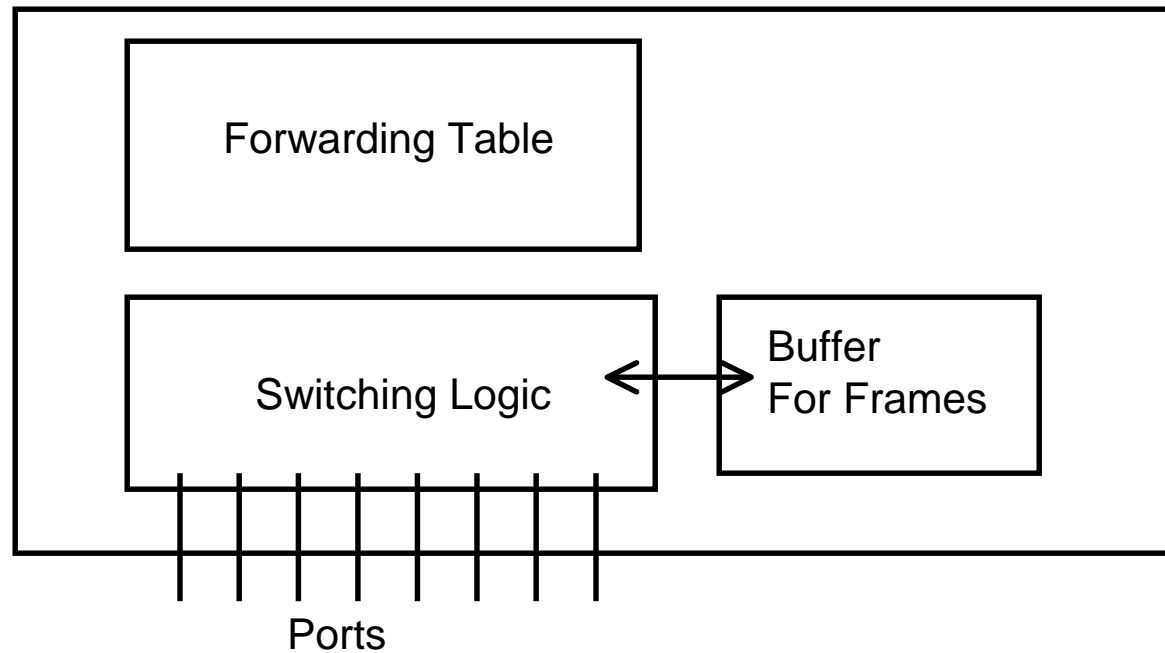
- Split one logical LAN into multiple physical LANs
 - permit mixing types of 802.X networks
 - 100 Megabit Ethernet with 10Mbps
 - token ring with Ethernet
 - extend the physical network
 - limits on cable length
 - improve security
 - reduce traffic
- Forward traffic between the physical layers
 - regenerate the signal
 - convert between 802.X formats
 - this is non-trivial

Learning Bridges

- **Transparent to users**
 - traffic just gets to the correct location
 - no software configuration required
- **Selectively forward traffic among segments**
 - used 48bit Ethernet addresses
 - at first, forward all traffic via flooding
 - use **source** address to learn where a host is located
 - do not forward a packet if the destination is known to be on the local network
- **need to have a spanning tree to prevent loops**
 - use lowest serial number to elect root
 - compute shortest path to root as the spanning tree
 - some bridge may be disabled to ensure a tree

Switched Ethernet

- Make Bridge with multiple ports
 - different ports can each form their own Ethernet segment
 - frames for other segment travel over backplane
 - individual stations retain the same card and cabling



Source Routing Bridge (skipped Fall 2001)

- Each host knows how to reach other hosts
 - it builds a full path to that host
- Every LAN and bridge has a number
 - a LAN has a 12 bit identifier
 - a bridge a 4 bit id
- To discover a route
 - broadcast a discovery packet
 - destination responds
 - bridges fill in their information in the response
 - results in a full path to the remote destination

Source vs. Transparent Bridges (skipped Fall 2001)

- Source Bridges

- always use optimal routes
- could exploit multiple paths between two LANs for load sharing

- Transparent Bridges

- require no changes to nodes
 - nodes are now more complex
- no need to configure the bridges
 - source bridges need LAN and Bridge Ids

Faster Ethernet

- Based on hubs
 - advantages of hubs rendered bus cables useless
 - limits cable length to 100 meters for copper
 - can be switched or use a single collision domain
- 100 Mbps
 - 100Base-T4
 - uses 4 pair cat 3 wiring
 - 33Mbps in each direction and two reversible channels
 - 25Mhz with trinary signaling and 4 bits per baud
 - 100Base-TX
 - two pairs of cat 5 wiring
 - 125Mhz with 4bits out of 5 for data
- 1000 Mbps
 - Uses 4 pair of cat 5e wiring (at 250Mhz)
 - Uses fiber