

Announcements

- Reading
 - Chapter 6 (6.4 & 6.6)
- Midterm #1
 - Re-grades were returned
 - Changes ranged from -7 to $+10$

Multiplexing in the Transport Layer

- Upward multiplexing
 - putting multiple transport connections onto one network connection
 - used to accommodate pricing strategies that charge for connections
- Downward multiplexing
 - using several network connections per transport connection
 - permits use of multiple copies of network resources
 - if the network layer uses sliding windows
 - a high latency network may under utilize the link
 - multiple connections each get a window
 - per connection buffer allocation
 - get more buffers
 - round-robin scheduling
 - get a larger share of link bandwidth

Crash Recovery

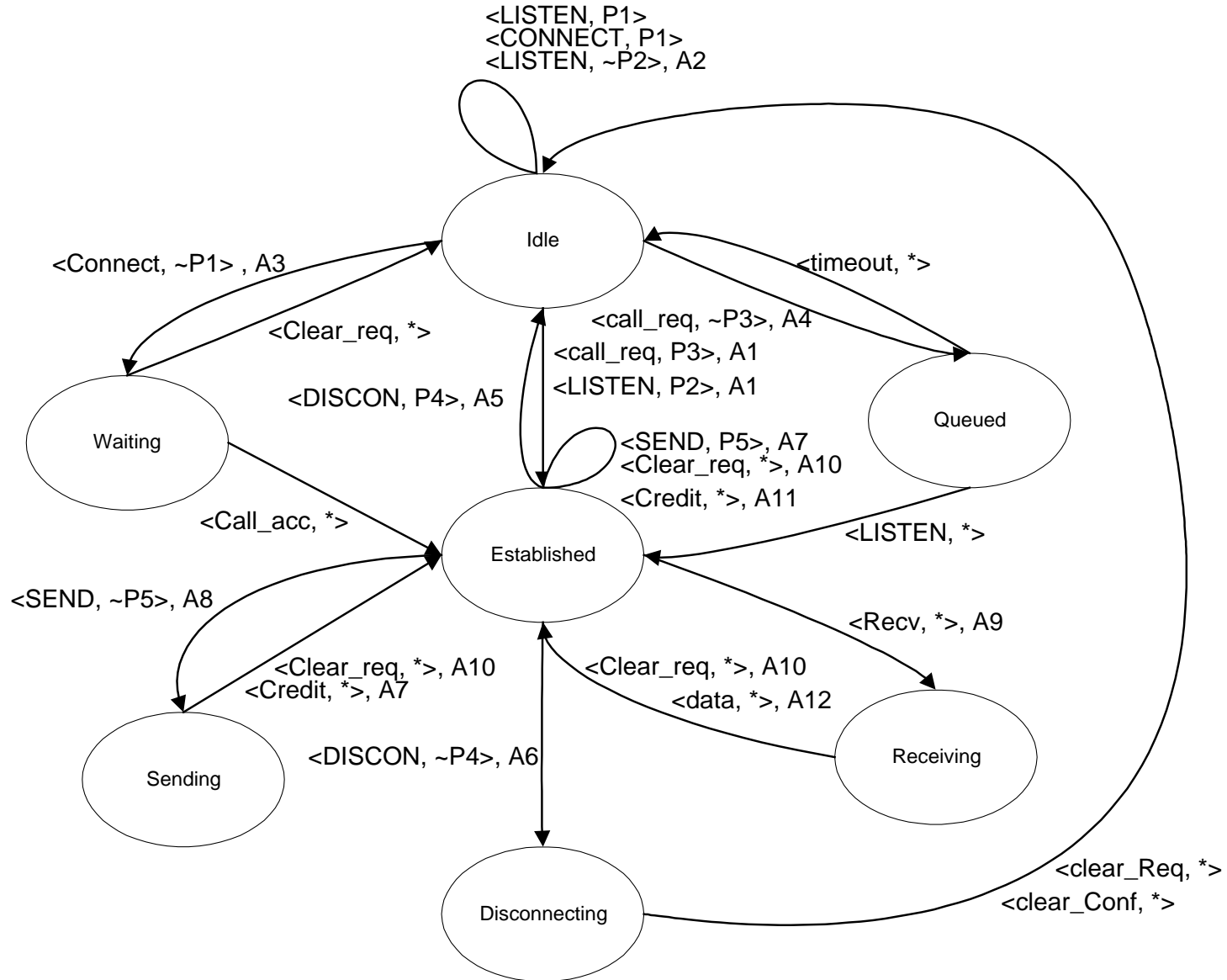
- Router or Link Crashes

- Data in transit can be lost.
- End nodes have sufficient state to recover lost data.
- Transport protocol can hide network failures from the application.

- Host Crashes

- Transport level state will be lost at one end.
- Does the transport layer have sufficient info to recover?, **No!**
 - Information must flow down to network and up to transport user
 - ACKs go down, and data goes up.
 - It is not possible to make these two operations atomic.
 - lack of stable storage causes this problem
- Result, higher up layer must deal with host crashes

Protocol State Machines



Predicates And State Transitions

Pred	Meaning
P1	Connection table full
P2	Call_req pending
P3	LISTEN Pending
P4	Clear_req Pending
P5	Credit Available

Act	Meaning
A1	Send Call_acc
A2	Wait for Call_req
A3	Send Call_req
A4	Start Timer
A5	Send Clear_conf
A6	Send Clear_req
A7	Send message
A8	Wait for credit
A9	Send Credit
A10	Set Clr_req_rcv flag
A11	Record credit
A12	Accept message

TCP Timer Management

- Problem: How to pick timeout value?
 - need to estimate round-trip latency
 - need low variance in round trip latency
- Solution: dynamic estimates of RTT
 - $RTT = \alpha RTT + (1 - \alpha) M$
M time of an ACK
 $\alpha = 7/8$
 - Need to pick retransmission time
 - old policy, use Timeout = RTT β , with $\beta = 2$
 - estimate standard deviation of RTT using mean deviation
$$D = \alpha D + (1 - \alpha) | RTT - M |$$
$$\text{Timeout} = RTT + 4 * D$$
 - How to update RTT on retransmission's
 - double Timeout on a retransmission

Other TCP Timers

- Persistence Timer
 - Prevents deadlock due to dropped window packets
 - This is a problem if the window is set to 0
- Keepalive Timer
 - Prevents half dead connections
 - may consume bandwidth
 - may kill live connections when net hiccups
- TIMED Wait
 - prevents re-use of a connection before max packet life is over
 - set to twice max packet lifetime