# Announcements

- Reading 9.6-9.7

# Handling a page fault

1) Check if the reference is valid
   - if not, terminate the process

2) Find a page frame to allocate for the new process
   - for now we assume there is a free page frame.

3) Schedule a read operation to load the page from disk
   - we can run other processes while waiting for this to complete

4) Modify the page table entry to the page

5) Restart the faulting instruction
   - hardware normally will abort the instruction so we just return from the trap to the correct location.

# What happens when we fault and there are no more physical pages?

- Need to remove a page from main memory
  - if it is “dirty” we must store it to disk first.
    - dirty pages have been modified since they were last stored on disk.
- How to we pick a page?
  - Need to choose an appropriate algorithm
    - should it be global?
    - should it be local (one owned by the faulting process)

# Page Replacement Algorithms

- FIFO
  - Replace the page that was brought in longest ago
  - However
    - old pages may be great pages (frequently used)
    - number of page faults may increase when one increases number of page frames (discouraging!)
      - called belady's anomaly
      - 1,2,3,4,1,2,5,1,2,3,4,5 (consider 3 vs. 4 frames)
- Optimal
  - Replace the page that will be used furthest in the future
  - Good algorithm(!) but requires knowledge of the future
  - With good compiler assistance, knowledge of the future is sometimes possible

# Page Replacement Algorithms

- LRU
  - Replace the page that was actually used longest ago
  - Implementation of LRU can be a bit expensive
    - e.g. maintain a stack of nodes representing pages and put page on top of stack when the page is accessed
    - maintain a time stamp associated with each page
- Approximate LRU algorithms
  - maintain reference bit(s) which are set whenever a page is used
  - at the end of a given time period, reference bits are cleared

# FIFO Example (3 frames)

- Reference string: 1,2,3,4,1,2,5,1,2,3,4,5
  - access 1 - (1) fault
  - access 2 - (1,2) fault
  - access 3- (1,2,3) fault
  - access 4 - (2,3,4) fault, replacement
  - access 1 - (3,4,1) fault, replacement
  - access 2 - (4,1,2) fault, replacement
  - access 5 - (1,2,5) fault, replacement
  - access 1- (1,2,5)
  - access 2 - (1,2,5)
  - access 3 - (2,5,3) fault, replacement
  - access 4 - (5,3,4) fault, replacement
  - access 5 - (5,3,4)
- 9 page faults

# LRU Example (3 frames)

- Reference string: 1,2,3,4,1,2,5,1,2,3,4,5
  - access 1 - (1) fault
  - access 2 - (1,2) fault
  - access 3- (1,2,3) fault
  - access 4 - (2,3,4) fault, replacement
  - access 1 - (3,4,1) fault, replacement
  - access 2 - (4,1,2) fault, replacement
  - access 5 - (1,2,5) fault, replacement
  - access 1- (2,5,1)
  - access 2 - (5,1,2)
  - access 3 - (1,2,3) fault, replacement
  - access 4 - (2,3,4) fault, replacement
  - access 5 - (3,4,5) fault, replacement
- 10 page faults

# LRU Example (4 frames)

- Reference string: 1,2,3,4,1,2,5,1,2,3,4,5
  - access 1 - (1)   fault
  - access 2 - (1,2)  fault
  - access 3-  (1,2,3)  fault
  - access 4 - (1,2,3,4)  fault, replacement
  - access 1 - (2,3,4,1)
  - access 2 - (3,4,1,2)
  - access 5 - (4,1,2,5)  fault, replacement
  - access 1-  (4,2,5,1)
  - access 2 - (4,5,1,2)
  - access 3 - (5,1,2,3)  fault, replacement
  - access 4 - (1,2,3,4)  fault, replacement
  - access 5 - (2,3,4,5)  fault, replacement
- 8 faults

# FIFO Example (4 frames)

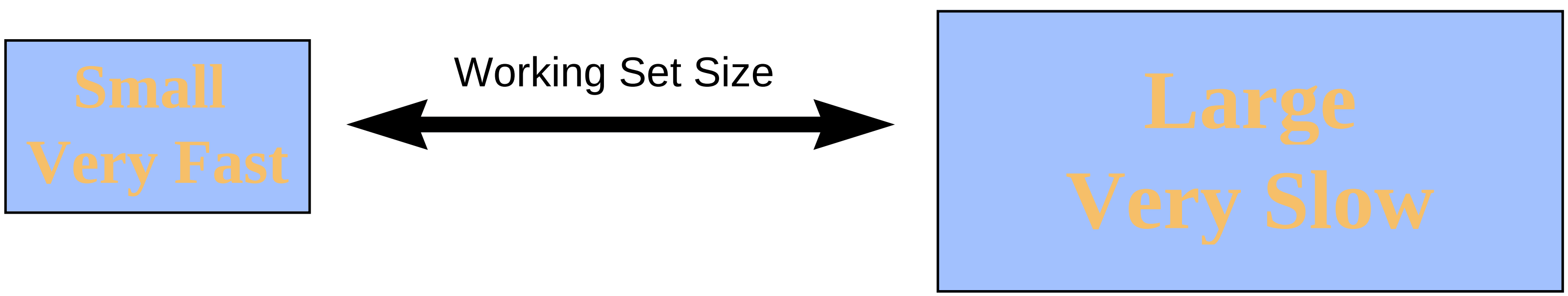
- Reference string: 1,2,3,4,1,2,5,1,2,3,4,5
  - access 1 - (1) fault
  - access 2 - (1,2) fault
  - access 3- (1,2,3) fault
  - access 4 - (1,2,3,4) fault, replacement
  - access 1 - (1,2,3,4)
  - access 2 - (1,2,3,4)
  - access 5 - (2,3,4,5) fault, replacement
  - access 1- (3,4,5,1) fault, replacement
  - access 2 - (4,5,1,2) fault, replacement
  - access 3 - (5,1,2,3) fault, replacement
  - access 4 - (1,2,3,4) fault, replacement
  - access 5 - (2,3,4,5) fault, replacement
- 10 Page faults

# Thrashing

- Virtual memory is not "free"
  - can allocate so much virtual memory that the system spends all its time getting pages
  - the situation is called thrashing
  - need to select one or more processes to swap out
- Swapping
  - write all of the memory of a process out to disk
  - don't run the process for a period of time
  - part of medium term scheduling
- How do we know when we are thrashing?
  - check CPU utilization?
  - check paging rate?
  - Answer: need to look at both
    - low CPU utilization plus high paging rate --> thrashing

# Working Sets and Page Replacement

- Programs usually display reference locality
  - temporal locality
    - repeated access to the same memory location
  - spatial locality
    - consecutive memory locations access nearby memory locations
  - memory hierarchy design relies heavily on locality reference
    - sequence of  nested storage media
- Working set
  - set of pages referenced in the last delta references



# Preventing Threashing

- Need to ensure that we can keep the working set in memory
  - if the working sets of the processes in memory exceed total page frames, then we need to swap a process out
- How do we compute the working set?
  - can approximate it using a reference bit

# Implementation Issues

- How big should a page be?
  - want to trade cost of fault vs. fragmentation
    - cost of fault is: trap + seek + latency + transfer
  - Does the OS page size have to equal the HW page size?
    - no, just needs to be a multiple of it
- How does I/O relate to paging
  - if we request I/O for a process, need to lock the page
    - if not, the I/O device can overwrite the page
- Can the kernel be paged?
  - most of it can be.
  - what about the code for the page fault handler?

# Segmentation

- Segmentation is used to give each program several independent protected address spaces
  - each segment is an independent protected address space
  - access to segments is controlled by data which describes size, privilege level required to access, protection (whether segment is read-only etc)
  - segments may or may not overlap
    - disjoint segments can be used to protect against programming errors
    - separate code, data stack segments

- Disjoint Segments can be used to exploit expanded address space
  - In 16 bit architectures e.g. (8086 and 80x86 in V86 mode) each segment has only 16 bits of address space
  - In distributed networks consisting of multiple 32 bit machines, segmentation can be used to support single huge address space
- Segments can span identical regions of address space - *flat model*
  - Windows NT and Windows '95 use 4 Gbyte code segments, stack segments, data segments