Announcements

- Lecture provided by Bobby Bhattacharjee

Page Replacement Algorithms

- FIFO
  - Replace the page that was brought in longest ago
  - However
    - old pages may be great pages (frequently used)
    - number of page faults may increase when one increases number of page frames (discouraging!)
      - called belady's anomaly
      - 1,2,3,4,1,2,5,1,2,3,4,5 (consider 3 vs. 4 frames)
- Optimal
  - Replace the page that will be used furthest in the future
  - Good algorithm(!) but requires knowledge of the future
  - With good compiler assistance, knowledge of the future is sometimes possible

Page Replacement Algorithms

- LRU
  - Replace the page that was actually used longest ago
  - Implementation of LRU can be a bit expensive
    - e.g. maintain a stack of nodes representing pages and put page on top of stack when the page is accessed
    - maintain a time stamp associated with each page
- Approximate LRU algorithms
  - maintain reference bit(s) which are set whenever a page is used
  - at the end of a given time period, reference bits are cleared

FIFO Example (3 frames)

- Reference string: 1,2,3,4,1,2,5,1,2,3,4,5
  - access 1 - (1) fault
  - access 2 - (1,2) fault
  - access 3- (1,2,3) fault
  - access 4 - (2,3,4) fault, replacement
  - access 1 - (3,4,1) fault, replacement
  - access 2 - (4,1,2) fault, replacement
  - access 5 - (1,2,5) fault, replacement
  - access 1- (1,2,5)
  - access 2 - (1,2,5)
  - access 3 - (2,5,3) fault, replacement
  - access 4 - (5,3,4) fault, replacement
  - access 5 - (5,3,4)
- 9 page faults

LRU Example (3 frames)

- Reference string: 1,2,3,4,1,2,5,1,2,3,4,5
  - access 1 - (1) fault
  - access 2 - (1,2) fault
  - access 3- (1,2,3) fault
  - access 4 - (2,3,4) fault, replacement
  - access 1 - (3,4,1) fault, replacement
  - access 2 - (4,1,2) fault, replacement
  - access 5 - (1,2,5) fault, replacement
  - access 1- (2,5,1)
  - access 2 - (5,1,2)
  - access 3 - (1,2,3) fault, replacement
  - access 4 - (2,3,4) fault, replacement
  - access 5 - (3,4,5) fault, replacement
- 10 page faults

LRU Example (4 frames)

- Reference string: 1,2,3,4,1,2,5,1,2,3,4,5
  - access 1 - (1) fault
  - access 2 - (1,2) fault
  - access 3- (1,2,3) fault
  - access 4 - (1,2,3,4) fault, replacement
  - access 1 - (2,3,4,1)
  - access 2 - (3,4,1,2)
  - access 5 - (4,1,2,5) fault, replacement
  - access 1- (4,2,5,1)
  - access 2 - (4,5,1,2)
  - access 3 - (5,1,2,3) fault, replacement
  - access 4 - (1,2,3,4) fault, replacement
  - access 5 - (2,3,4,5) fault, replacement
- 8 faults

FIFO Example (4 frames)

- Reference string: 1,2,3,4,1,2,5,1,2,3,4,5
  - access 1 - (1) fault
  - access 2 - (1,2) fault
  - access 3- (1,2,3) fault
  - access 4 - (1,2,3,4) fault, replacement
  - access 1 - (1,2,3,4)
  - access 2 - (1,2,3,4)
  - access 5 - (2,3,4,5) fault, replacement
  - access 1- (3,4,5,1) fault, replacement
  - access 2 - (4,5,1,2) fault, replacement
  - access 3 - (5,1,2,3) fault, replacement
  - access 4 - (1,2,3,4) fault, replacement
  - access 5 - (2,3,4,5) fault, replacement
- 10 Page faults

Thrashing

- Virtual memory is not “free”
  - can allocate so much virtual memory that the system spends all its time getting pages
  - the situation is called thrashing
  - need to select one or more processes to swap out
- Swapping
  - write all of the memory of a process out to disk
  - don’t run the process for a period of time
  - part of medium term scheduling
- How do we know when we are thrashing?
  - check CPU utilization?
  - check paging rate?
  - Answer: need to look at both
    - low CPU utilization plus high paging rate --> thrashing