

Loop Invariants

Annotating Programs

- General intuition behind annotations: label points in program with assertions that should hold when control is at that point!
 - You can do this using your intuition
 - Strong postconditions / weakest preconditions give you a systematic way to generate these assertions
 - In many cases (e.g. assignment, statement blocks, if-then-else) strongest postconditions / weakest preconditions can be computed automatically!
- When is an annotation of a piece of code complete and correct?
 - An annotation is complete if every statement in the code has both a precondition and a postcondition (these will be shared: the postcondition of one statement will be a precondition of the following statement)
 - An annotation is correct if every embedded Hoare triple is valid
- If an annotation is complete and correct, then the Hoare triple consisting of the precondition of the code, the code itself, and the postcondition is valid!

$$\{Q \ [X \mapsto a]\} \ X:=a \ \{Q\}$$

(hoare_asgn)

$$\{P\} \ c_1 \ \{Q\}$$
$$\{Q\} \ c_2 \ \{R\}$$

$$\{P\} \ c_1; c_2 \ \{R\}$$

(hoare_seq)

$$\{P \wedge b\} \ c_1 \ \{Q\}$$
$$\{P \wedge \sim b\} \ c_2 \ \{Q\}$$

$$\{P\} \ \text{if } b \text{ then } c_1 \ \text{else } c_2 \ \text{end } \{Q\}$$

(hoare_if)

Recall: Three Key Concepts in Systematic Annotation Construction

- Strongest postconditions
- Weakest preconditions
- *Loop invariants*

Annotations and Loops

- Strongest postconditions / weakest preconditions still exist for loops!
- However, they cannot generally be computed automatically
- *Loop invariants* fill this gap
 - They are propositions
 - They must be added manually in Dafny
 - Once added, Dafny can check that they really are invariants!

Defining “Loop Invariant”

- Let code S be `while B { S' }` ($\{ S' \}$ is the loop body)
- Then a proposition I is a *loop invariant* for S if and only if $\{ I \wedge B \} S' \{ I \}$ is valid
 - If you start S' in a state satisfying I and loop condition B ...
 - ... then whenever S' terminates the result state satisfies I !
- This means that as the loop “loops”, I is being kept true
- Also: if I is a loop invariant for S then $\{ I \} S \{ I \wedge \neg B \}$ is valid
 - If loop terminates then B must be false (so $\neg B$ must be true)
 - Since loop body keeps I true, when loop exists $I \wedge \neg B$ must hold!

$$\{Q \ [X \mapsto a]\} \ X := a \ \{Q\}$$

(hoare_asgn)

$$\begin{array}{l} \{P\} \ c_1 \ \{Q\} \\ \{Q\} \ c_2 \ \{R\} \end{array}$$

$$\{P\} \ c_1; c_2 \ \{R\}$$

(hoare_seq)

$$\begin{array}{l} \{P \wedge b\} \ c_1 \ \{Q\} \\ \{P \wedge \sim b\} \ c_2 \ \{Q\} \end{array}$$

$$\{P\} \ \text{if } b \ \text{then } c_1 \ \text{else } c_2 \ \text{end } \{Q\}$$

(hoare_if)

$$\{P \wedge b\} \ c \ \{P\}$$

$$\{P\} \ \text{while } b \ \text{do } c \ \text{end } \{P \wedge \sim b\}$$

(hoare_while)

Loop Invariants in Dafny

```
method FindMinVal (a : array<int>) returns (min : int)
  requires a.Length > 0 // Precondition
  ensures forall i : int :: 0 <= i < a.Length ==> min <= a[i] // Postcondition
{
  min := a[0];
  var i := 1;
  while (i < a.Length)
    invariant
    {
      if a[i] < min {
        min := a[i];
      }
      i := i+1;
    }
}
```

- Declared with keyword “invariant” after loop invocation, before body
- You can have as many invariant declarations as you like; multiple invariants are interpreted as being conjoined

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  min := a[0];
  var i := 1;
  while (i < a.Length)
    invariant forall j : int :: 0 <= j < i ==> min <= a[j]
    {
      if a[i] < min {
        min := a[i];
      }
      i := i+1;
    }
}
```

- Declared with keyword “invariant” after loop invocation, before body
- You can have as many invariant declarations as you like; multiple invariants are interpreted as being conjoined

Strengthening Invariants

- Sometimes Dafny complains that it cannot complete the verification of a given invariant
- Often you can add extra invariants to give facts to Dafny that it needs

Adding Invariants

```
method FindMinVal (a : array<int>) returns (min : int)
  requires a.Length > 0 // Precondition
  ensures forall i : int :: 0 <= i < a.Length ==> min <= a[i] // Postcondition
{
  min := a[0];
  var i := 1;
  while (i < a.Length)
    invariant 0 <= i <= a.Length // Extra invariant to constrain i
    invariant forall j : int :: 0 <= j < i ==> min <= a[j]
    {
      if a[i] < min {
        min := a[i];
      }
      i := i+1;
    }
}
```

- Dafny could not complete the previous proof because it did not know that $i \leq a.Length$ is preserved by the loop
- Adding this enables completion of verification

Another Example

```
method Search (key : int, a : array<int>) returns (found : bool)
  ensures found <==> exists i : int :: 0 <= i < a.Length && key == a[i]
{
  var i : int := 0;
  found := false;
  while (i < a.Length)
    invariant i <= a.Length;
    invariant found <==> exists j : int :: 0 <= j < i && key == a[j]
    {
      if (key == a[i])
      {
        found := true;
      }
      i := i+1;
    }
}
```

Yet Another Example

```
method Locate (key : int, a : array<int>) returns (found : bool, index : int)
  ensures -1 <= index < a.Length
  ensures found ==> index >= 0 && key == a[index]
  ensures !found ==> index == -1
{
  var i : int := 0;
  found := false;
  index := -1;
  while (i < a.Length)
    invariant i <= a.Length
    invariant found ==> key == a[index]
    invariant (!found) ==> index == -1
    {
      if (key == a[i])
      {
        return true, i;
      }
      i := i+1;
    }
}
```

Verifying Methods in Dafny

- Add requires, ensures clauses
- Add invariants to all loops
- If it verifies, you are done!
- Otherwise
 - Strengthen / weaken invariants
 - Strengthen requires, ensures
 - Start constructing the annotation on your own to see if that helps Dafny