

CMSC330 - Organization of Programming Languages Spring 2023 - Final

CMSC330 Course Staff
University of Maryland
Department of Computer Science

Name: _____

UID: _____

I pledge on my honor that I have not given or received any unauthorized assistance on this assignment/examination

Signature: _____

Ground Rules

- You may use anything on the accompanying reference sheet anywhere on this exam
- Please write legibly. **If we cannot read your answer you will not receive credit**
- You may not leave the room or hand in your exam within the last 10 minutes of the exam
- If anything is unclear, ask a proctor. If you are still confused, write down your assumptions in the margin

Question	Points
Q1	10
Q2	7
Q3	15
Q4	15
Q5	12
Q6	15
Q7	18
Q8	8
EC	5
Total	100 + 5

Problem 1: Language Concepts

[Total 10 pts]

$(\lambda x.abx)$ is alpha-equivalent to $(\lambda c.xyc)$	True	False
For statically typed languages, type checking occurs during the parsing phase	True	False
Dangling Pointers are prevented in Rust	True	False
Lifetimes are part of a variable's type in Rust	True	False
"Missing semicolon on line 12" is an error that would be raised during evaluation	True	False
$S \rightarrow S - S n$ is an ambiguous grammar	True	False
Grammar is a subset of Syntax	True	False
Mark and Sweep is faster than Reference Counting on average	True	False
A rust function with the following header will compile: <code>fn myst(a:&str, b:&u32, c:&u32) -> &str</code>	True	False
OCaml's <code>let x = x + 1 in x</code> is operationally the same as Ruby's <code>x = x + 1</code>	True	False

Problem 2: Regex

[Total 7 pts]

(a) Which of the following strings are accepted by the regular expression below?

 $/[\lambda\delta\sigma]^+\omega|\beta/$

Circle NONE if none of the first five (5) options match.

[3 pts]

 $\lambda\lambda\beta$ δ $\delta\omega\lambda$ $\sigma\lambda\beta\beta$ $\omega\beta$

NONE

(b) Write a regular expression that describes a comma separated integer list of odd length.

[4 pts]

	Valid	Invalid
Examples:	1	1,2
	1,2,3	1.3
	-6,-1,-3	

Problem 3: Higher Order Functions

Given the following type, write an expression that matches that type. You may not use type annotations and all pattern matching must be exhaustive. **You must use map or fold in your answer**

(a) `string list -> string`

(b) `'a list -> 'b list -> ('a list -> 'b -> 'a list) -> ('a -> 'c) -> 'c list`

Given the expression, write down its type. **You will need to evaluate it first**

(c) `fun a b c -> if a b then [b+1] else c`

(d) `(fun x -> fun y -> y x) ((fun y -> y + 1) 5)`

(e) `let c = if true then false else true in fun a -> fun b c -> b c > a c`

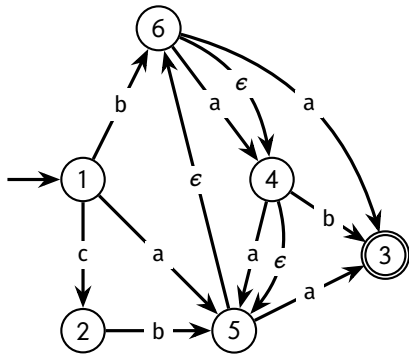
Problem 4: Finite State Machines

[Total 15 pts]

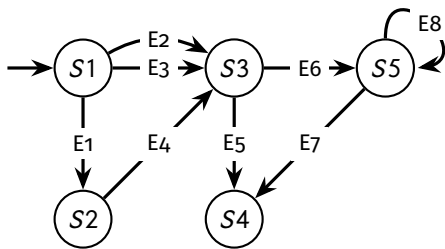
Using the subset algorithm, convert the following NFA to a DFA, and fill in the blanks appropriately matching the DFA provided with the right nodes and transitions. Only the blanks will be graded.

NFA:

Scratch Space (if needed)



DFA:



S1: S2: S3: S4: S5:

E1: E2: E3: E4:

E5: E6: E7: E8:

Final States:

S1 S2 S3 S4 S5

Problem 5: Operational Semantics

Consider the following rules for 2 Languages, using Ruby as the Metalanguage:

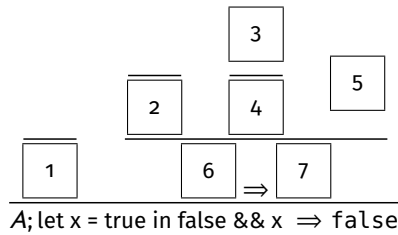
<p style="text-align: center;">Language 1:</p> $\frac{}{\text{true} \rightarrow \text{true}}$ $\frac{}{\text{false} \rightarrow \text{false}}$ $\frac{A(x) = v}{A; x \Rightarrow v}$ $\frac{A; e_1 \Rightarrow v_1 \quad A; e_2 \Rightarrow v_2 \quad v_3 = v_1 \text{ and } v_2}{A; e_1 \ \&\& \ e_2 \Rightarrow v_3}$ $\frac{A; e_1 \Rightarrow v_1 \quad A, x : v_1; e_2 \Rightarrow v_2}{A; \text{let } x = e_1 \text{ in } e_2 \Rightarrow v_2}$	<p style="text-align: center;">Language 2</p> $\frac{}{\text{true} \rightarrow \text{true}}$ $\frac{}{\text{false} \rightarrow \text{false}}$ $\frac{A(x) = v}{A; x \Rightarrow v}$ $\frac{A; e_1 \Rightarrow v_1 \quad A; e_2 \Rightarrow v_2 \quad v_3 = v_1 \text{ and } v_2}{A; (\lambda x. \lambda y. x y x) e_1 e_2 \Rightarrow v_3}$ $\frac{A; e_2 \Rightarrow v_1 \quad A, x : v_1; e_1 \Rightarrow v_2}{A; (\lambda x : e_1) e_2 \Rightarrow v_2}$
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(a) Convert the following Language 1 sentence to it's language 2 counterpart

A; let x = true in false && x

(b) Complete the opsem proof for the following program using Language 1:

let x = true in false && x



Blank 1: Blank 2: Blank 3:

Blank 4: Blank 5: Blank 6:

Blank 7:

Problem 6: Lambda Calculus

[Total 15 pts]

Perform a single β -reduction using lazy (call by name) evaluation on the outermost expression. If you cannot reduce it, write **Beta Normal Form**. Do **not** α -convert your final answer.

(a) $(a \lambda x. x a)(\lambda y. y y)$

[3 pts]

Perform a single β -reduction using Eager (call by value) evaluation on the outermost expression. If you cannot reduce it, write **Beta Normal Form**. Do **not** α -convert your final answer.

(b) $(\lambda x. a b c)((\lambda x. (x x)) x)$

[3 pts]

Convert the following expressions to Beta Normal Form. If it is already in Beta Normal Form, circle BNF. If the answer is not given, circle None.

(c) $(\lambda x. \lambda y. x y)((\lambda b. b b) y)$

[3 pts]

 $\lambda y. y y y$ $\lambda y. x x y$ $\lambda a. y y a$ $y y y$

BNF

infinite recursion

None

(d) $(\lambda x. x x x) (\lambda x. x x x)$

[3 pts]

 $(\lambda x. x x x)$ $x x x$ $(\lambda x. x x x)(\lambda x. x x x)$ x

BNF

infinite recursion

None

(e) $\lambda x. (\lambda b. a b) (\lambda b. a b)$

[3 pts]

 $\lambda x. (\lambda b. a b)$ $(\lambda b. a b)$ $a b$ $\lambda x. a \lambda b. a b$

BNF

infinite recursion

None

Problem 7: Coding

Consider the following Grammar, where n is any integer:

$$\begin{aligned} S &\rightarrow N + S|(N) \\ N &\rightarrow n \end{aligned}$$

(a) Ruby Lexer

Write a lexer for this grammar in **Ruby**, you may use the following as tokens

```
# tokens: n, "Plus", "RParen", "LParen"
# example input-output
lex("2 * -5 + 6") = IOError
lex("2 -7 9 -10") = ["2", "-7", "9", "-10"]
lex("(-2) + (3)") = ["LParen", "-2", "RParen", "Plus", "LParen", "3", "RParen"]
#If an error occurs, you may raise an error
raise IOError.new("Error")
```

```
def lex(str)
```

(b) Ocaml Parser

[10 pts]

Using the same grammar as before, where n is any integer:

$$\begin{aligned} S &\rightarrow N + S | (N) \\ N &\rightarrow n \end{aligned}$$

Write a parser for the S non-terminal in **OCaml**. You may use the following types and functions:

```
type tok = Int of int | Plus | RParen | LParen
```

```
type tree = Add of tree * tree | Leaf of int
```

```
let lookahead toks = match toks with [] -> None | h::t -> Some h
```

```
let match_tok toks tok = match toks [] -> raise Error | h::t when h = tok -> t | _ -> raise Error
```

```
(* You may assume raise Error is valid and compiles *)
```

You may assume there is a parse_n function of type tok list \rightarrow (tree * tok list) and that it is correct.

The type of parse_s is tok list \rightarrow (tree * tok list)

```
let rec parse_s toks =
```


Problem 8: Rust

```

1 fn main(){
2     let m = String::from("Hello");
3     let t = String::from("World");
4     { let y = m;
5       { let c = myfunc(y,t);
6         let d = &c;
7       }
8     }
9 }
10
11 fn myfunc<'a>(a:String, b: String) -> String{
12     if a.len() > b.len() {a} else {b}
13 }

```

Problem 9: Extra Credit

What is your favorite pun?

Problem 10: Extra Credit

Who is your discussion TA and what is your section number?

Ownership

If there is no owner, write "NONE".

Who is the owner of "Hello" immediately after line 11 is run?

Who is the owner of "World" immediately after line 5 is run?

Lifetimes

What is the last line executed before "Hello" dropped?

What is the last line executed before "World" dropped?

You may use this area as scratch space